

Center for
Quantum &
Topological
Systems

Quantum Data Types via Linear HoTT

presentation at:

Workshop on Quantum Software @ QTML 2022

Urs Schreiber (NYU Abu Dhabi)

on joint work at CQTS with

D. J. Myers, M. Riley,

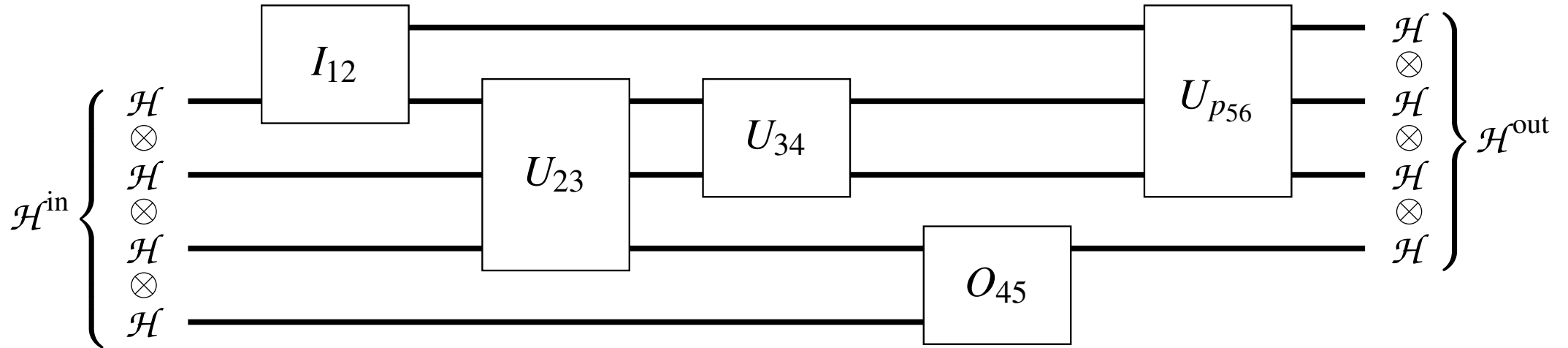
and Hisham Sati

slides and further pointers at: ncatlab.org/schreiber/show/QDataInLHoTT#QTML2022

The Problem

Pure quantum circuits are easy...

Linear operator composed & tensored from given *quantum logic gates*



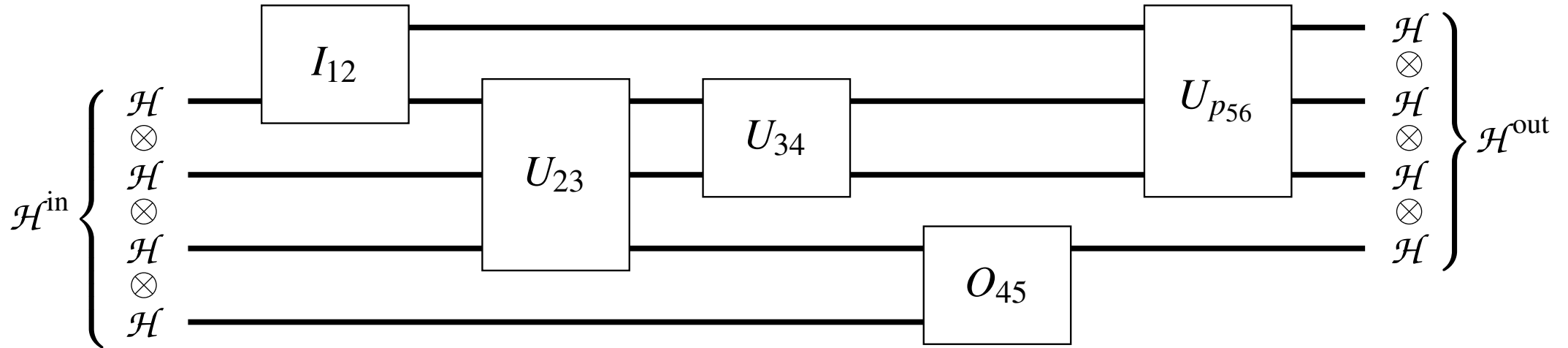
Hilbert space of possible **input** quantum states

linear transformation
upon execution

Hilbert space of possible **output** quantum states

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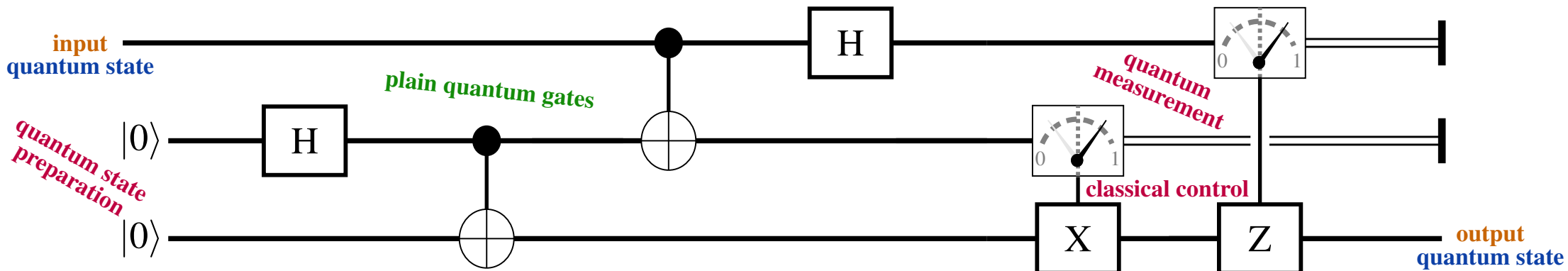
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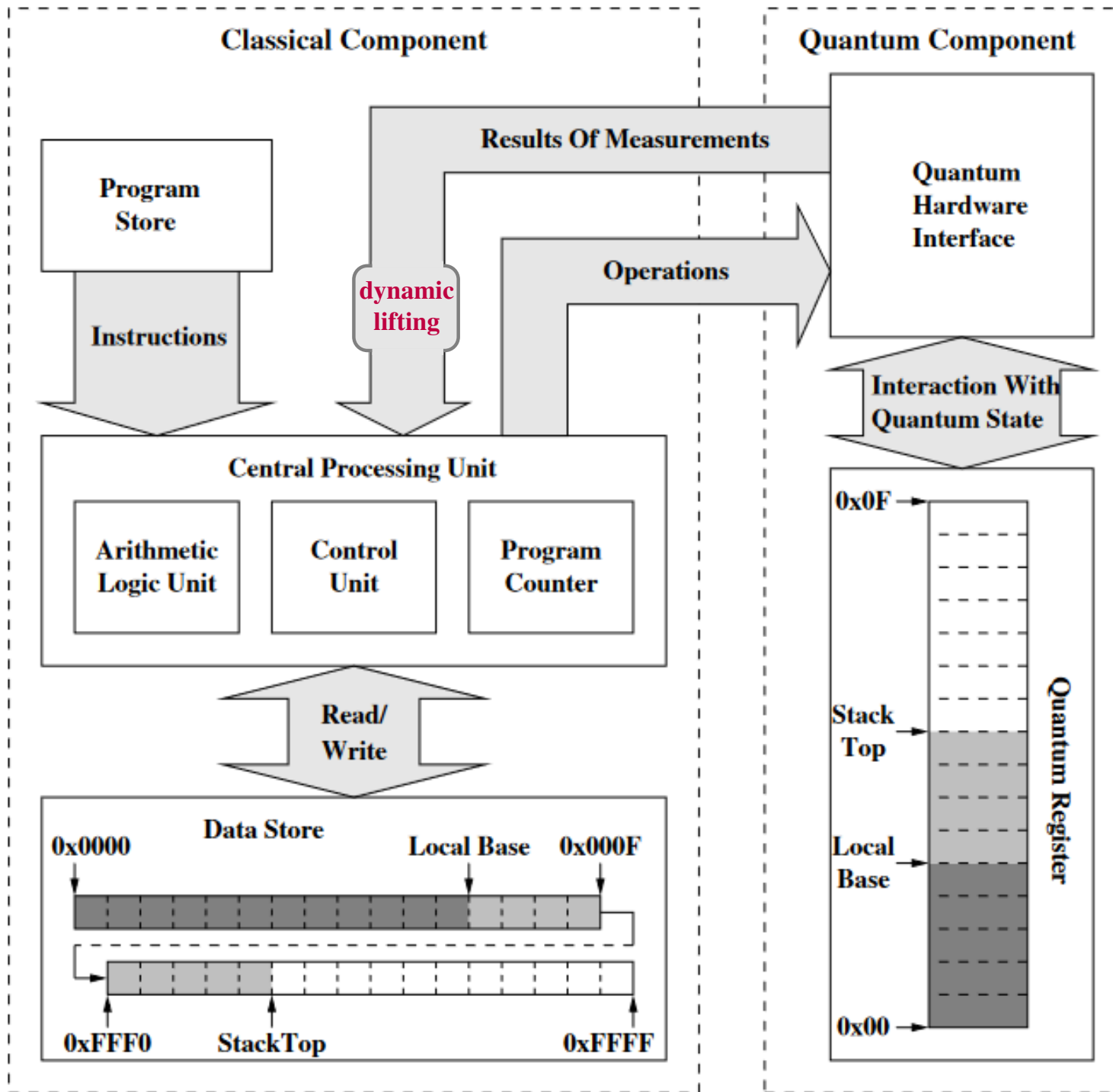
Hilbert space of possible **output** quantum states

but real quantum circuits have **classical control & effects**

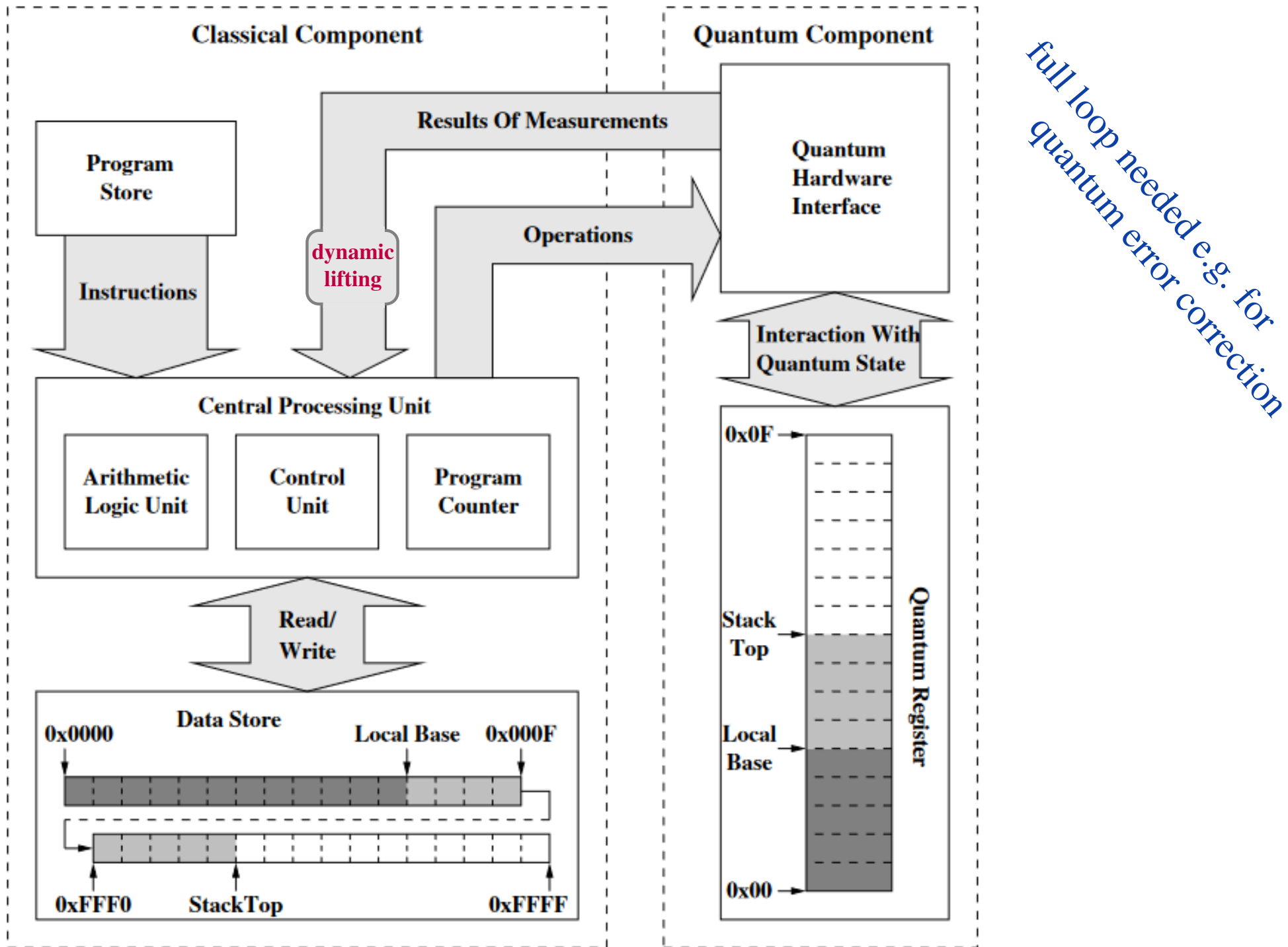
(Example: QBit Teleportation protocol)



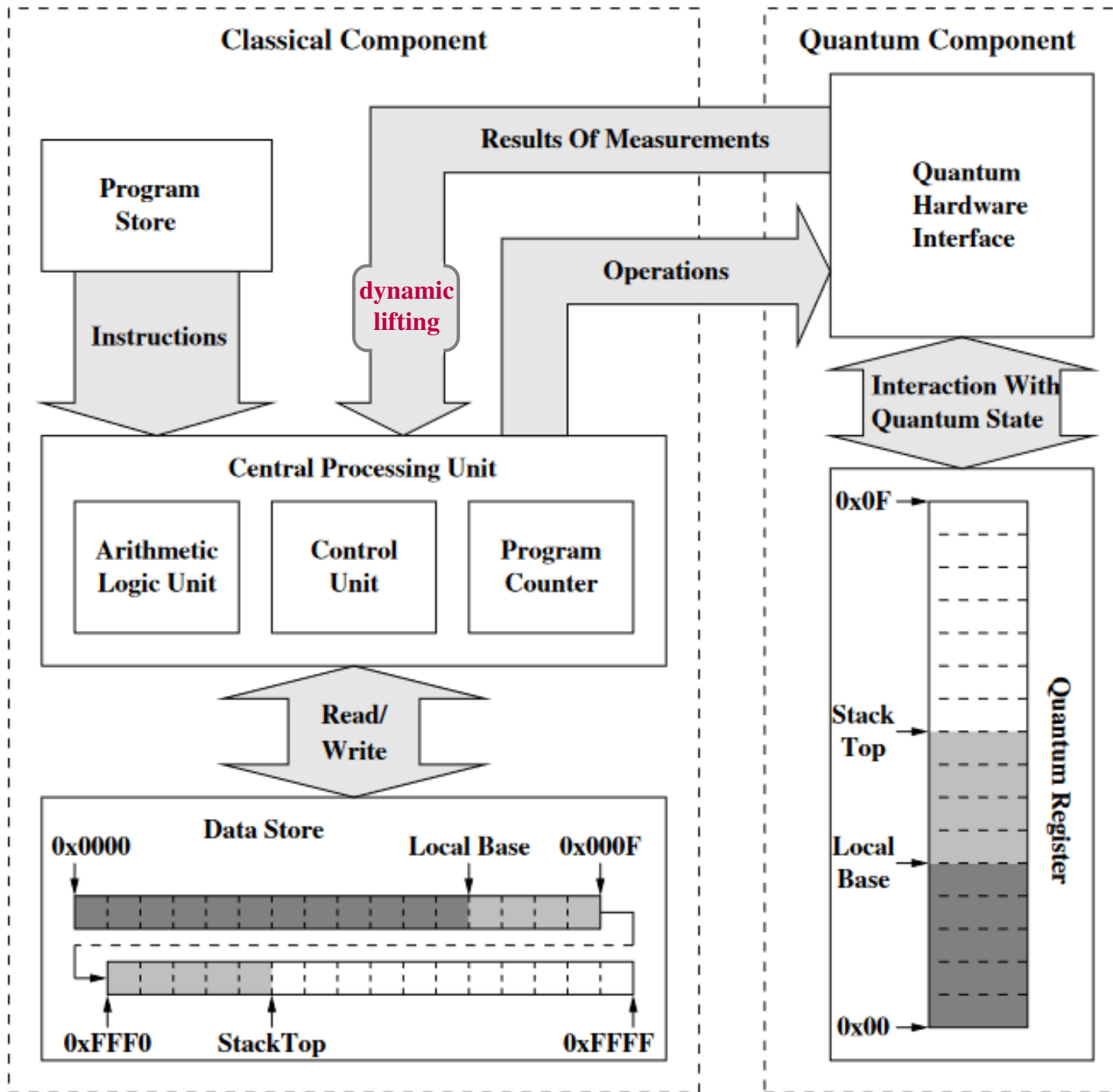
full reality is a loop: Classical $\xleftarrow{\text{measure}}$ Quantum $\xrightarrow{\text{prepare}}$



full reality is a loop: Classical $\begin{matrix} \leftarrow \text{measure} \\ \rightarrow \text{prepare} \end{matrix}$ Quantum

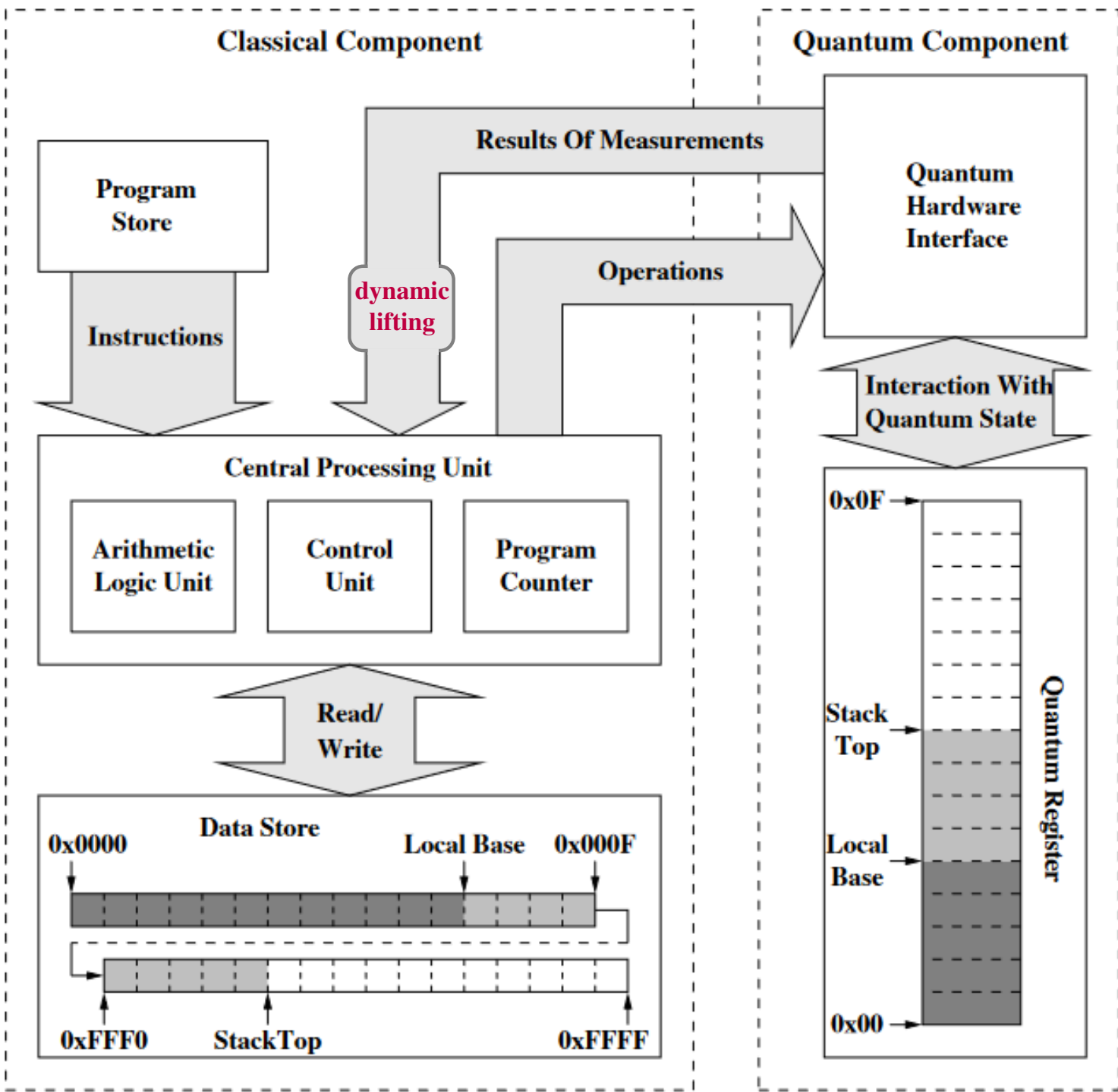


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existing models for dynamic lifting are ad hoc & unverified

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are embedded inside *classical* type theories:

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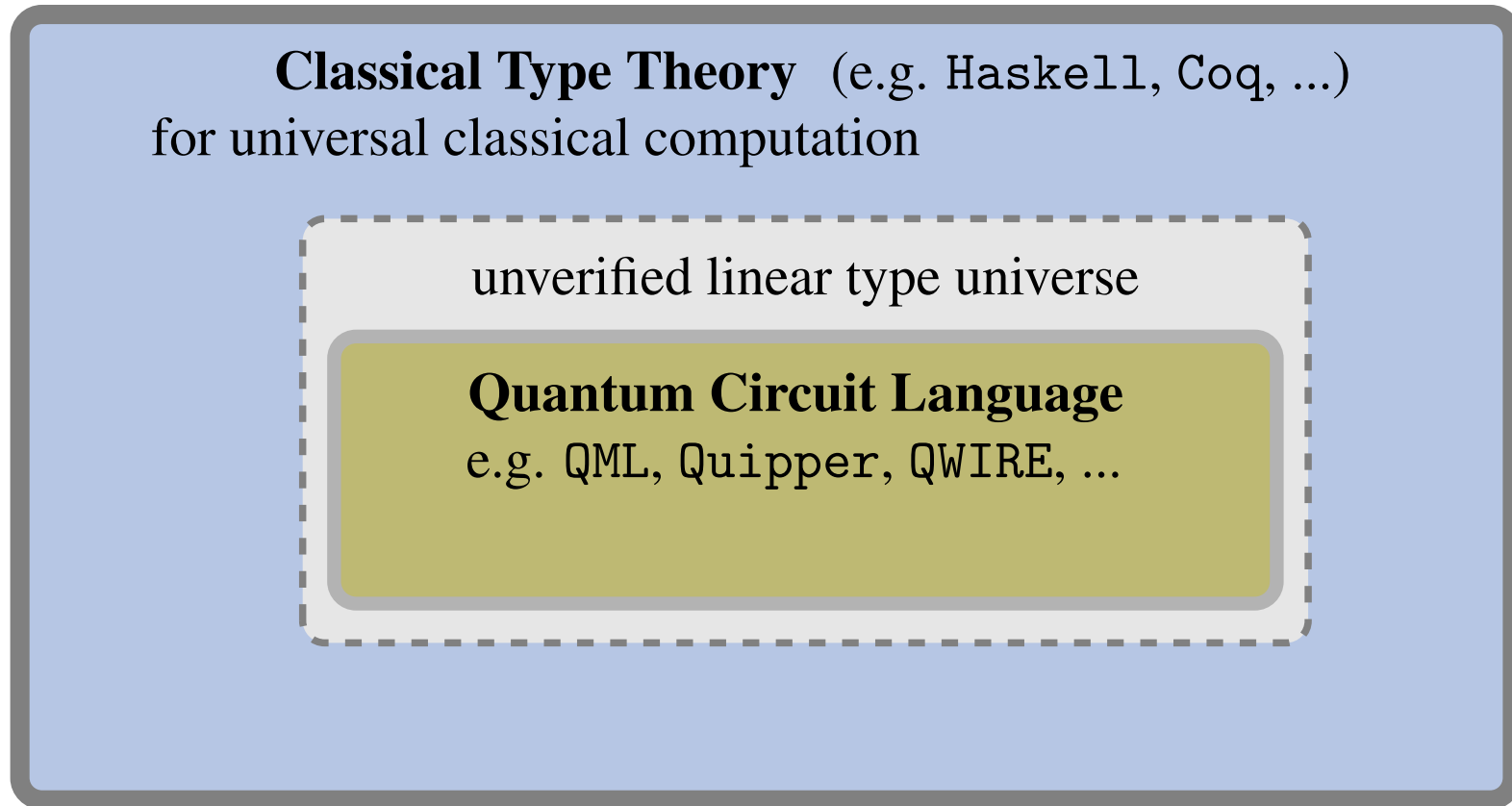
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Quantum Circuit Language

e.g. QML, Quipper, QWIRE, ...

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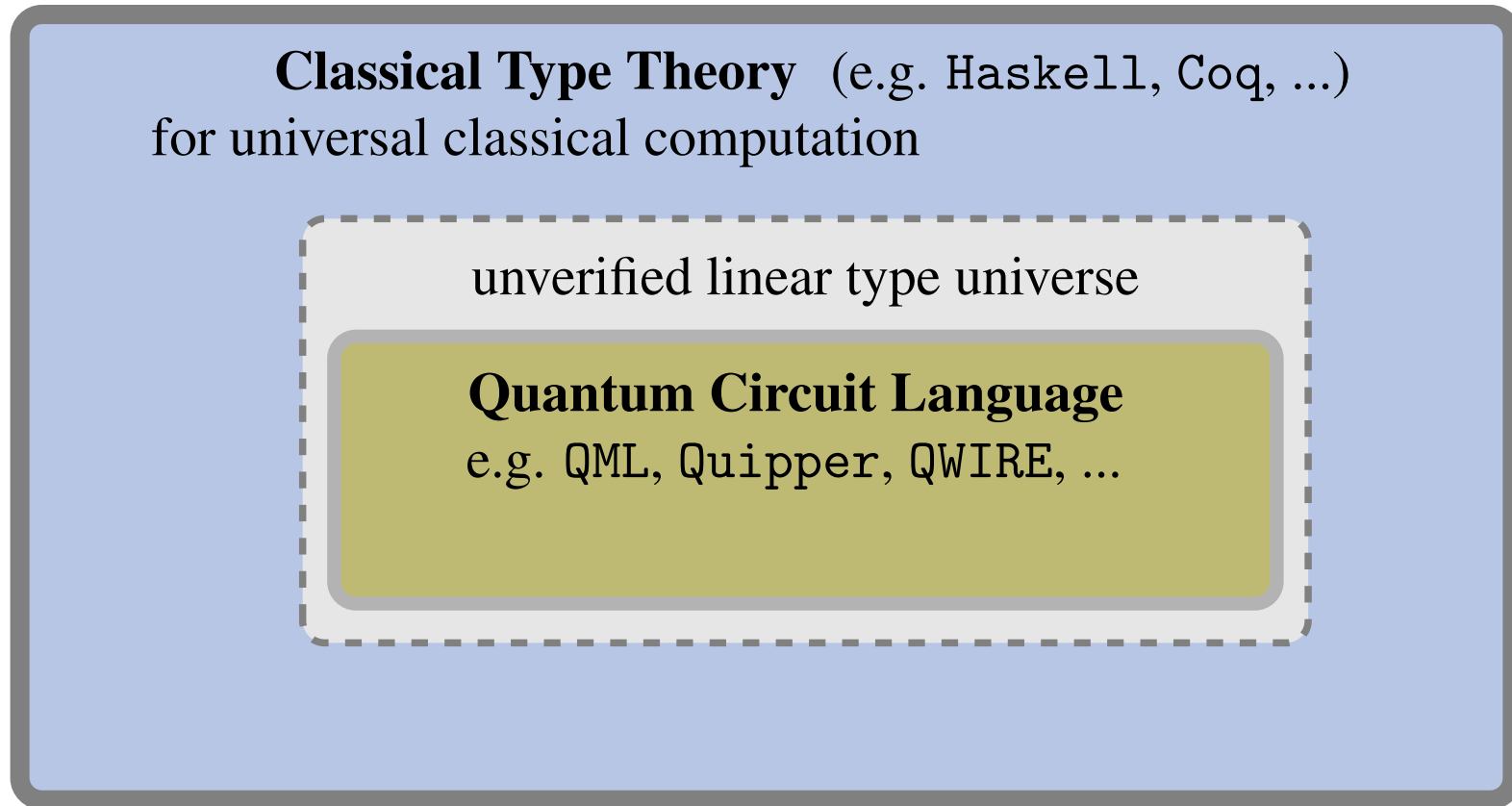
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Until now...

Our Solution

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dLHoTT is like a quantum microscope for Classical Data Types

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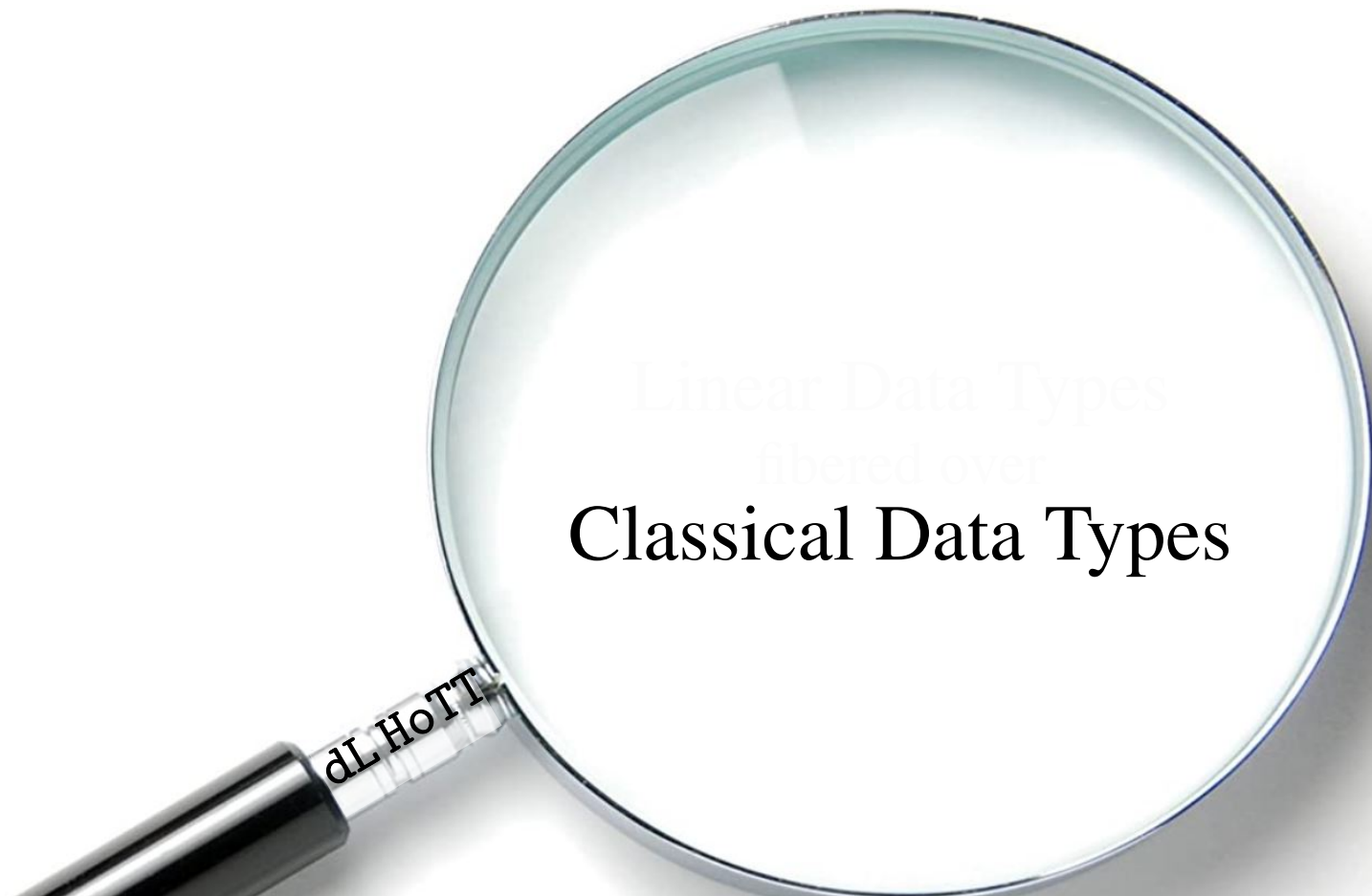
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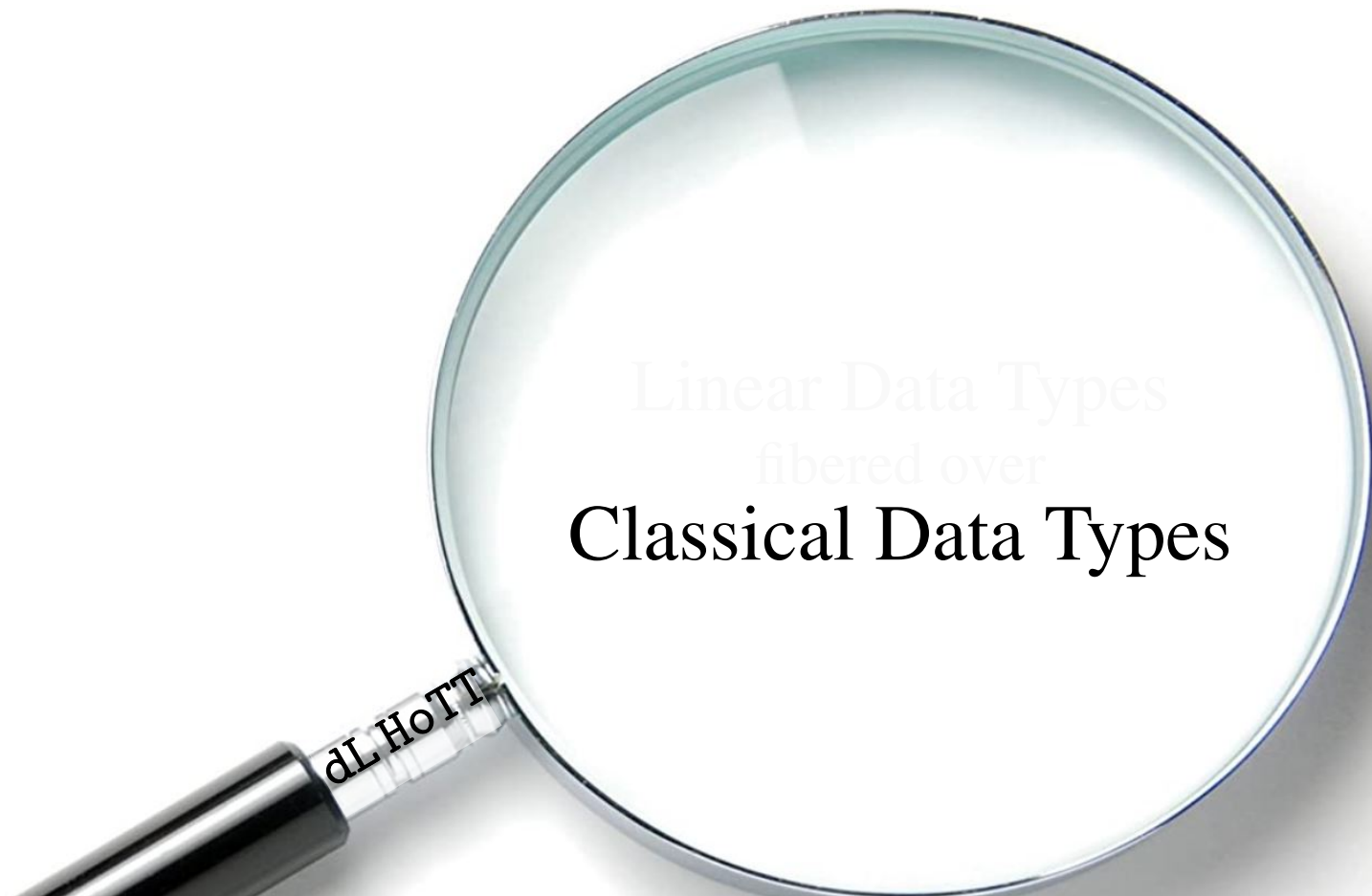
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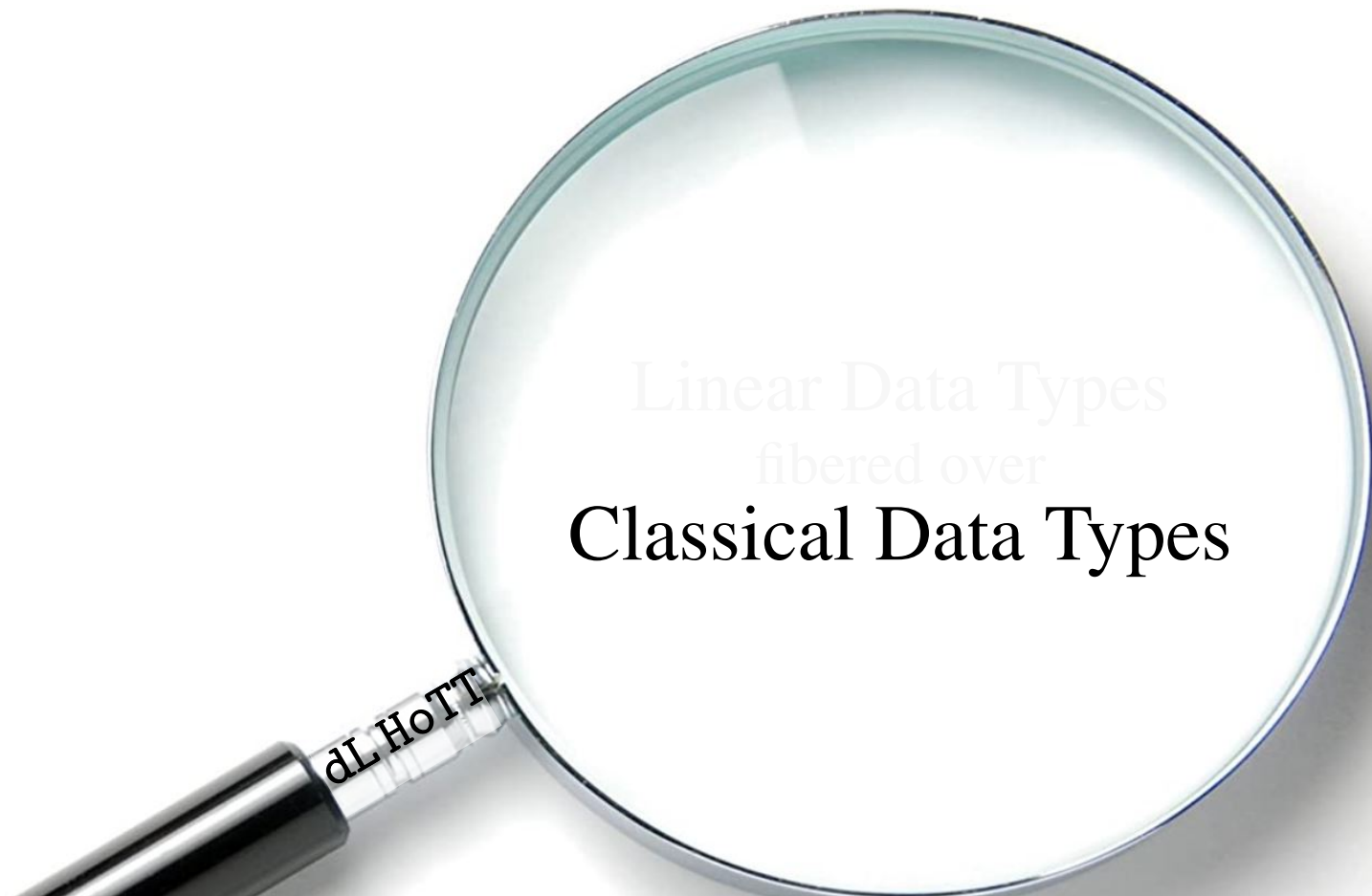
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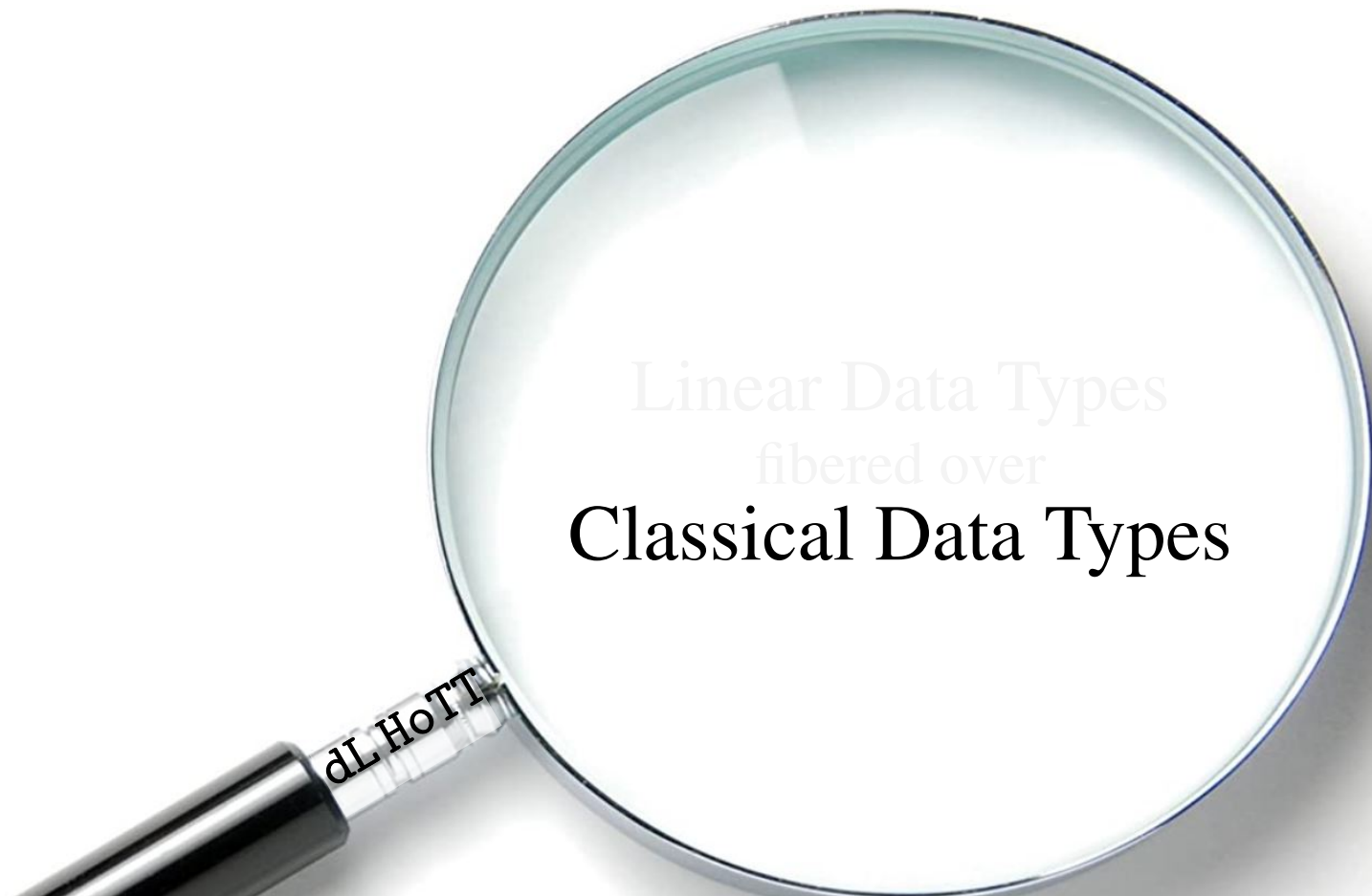
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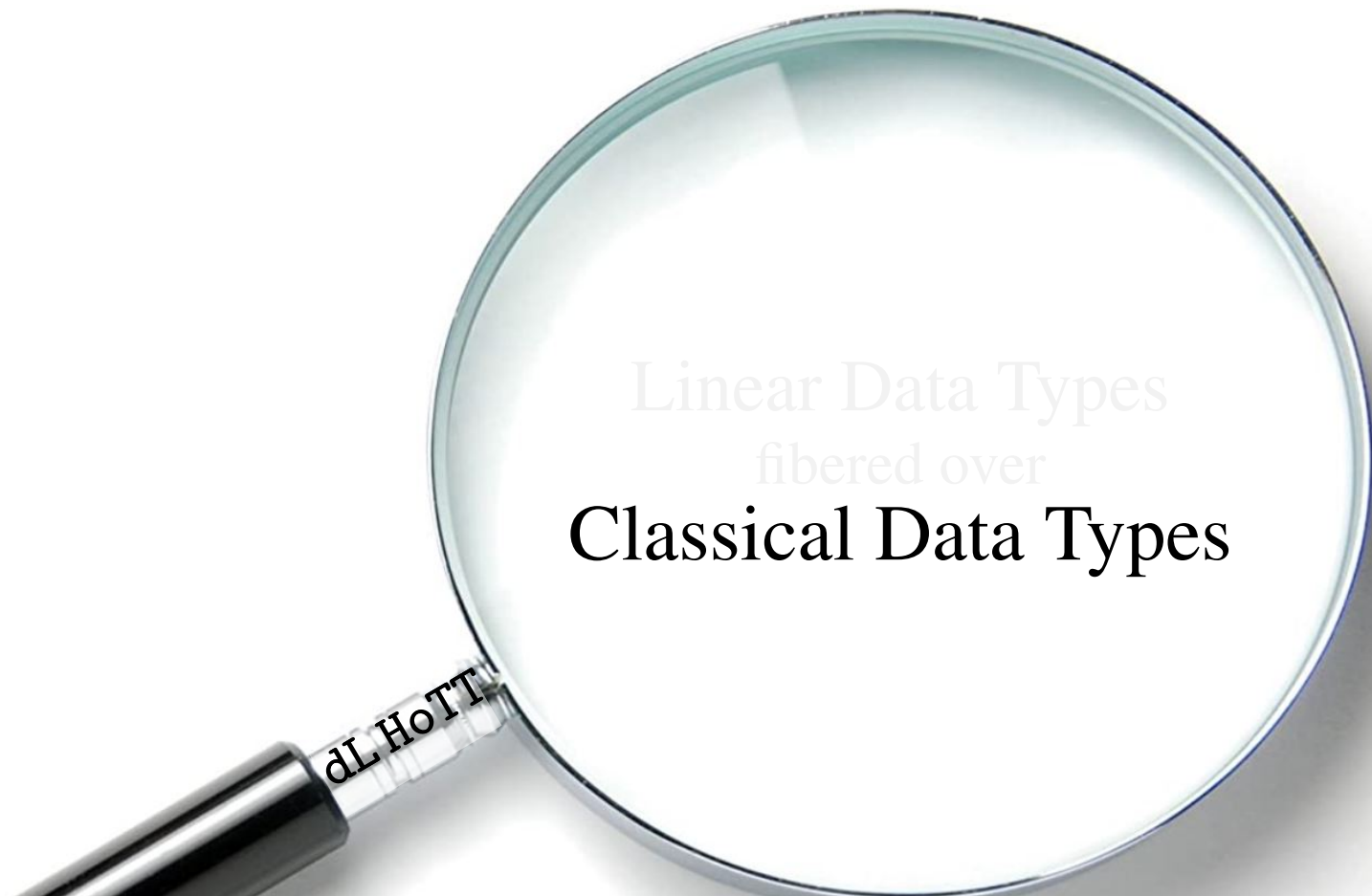
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constituting

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quantum measurement is handling of linear indefiniteness effects

quantum state preparation is handling of linear randomness co-effects

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\exists *universal* quantum+classical specification language

conservative over classical *Homotopy Type Theory* (HoTT)

and

verifying axiom scheme “**Motivic Yoga**” [Riley, §2.4, anticipated in S. (2014), §3.2]

(i.e. Grothendieck’s six operations *à la* Wirthmüller — more on all this below)

Theorem [CQTS (2022)]:

Motivic Yoga induces a system of monadic computational effects

constituting

linear modalities of actuality and potentiality

which happen to

know all about quantum information theory:

quantum measurement is handling of linear indefiniteness effects

quantum state preparation is handling of linear randomness co-effects

quantum+classical circuits are the effectful string diagrams

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ambient dLHoTT

verifies

classically dependent quantum linear types

ambient HoTT

provides

specification of topological quantum gates

ambient dTT

provides

full verified classical control

Quantum Data Types

Linear/Quantum Data Types

Characteristic Property			
Symbol			
Formula (for $B : \text{FinType}$)			
AlgTop Jargon			
Linear Logic			
Physics Meaning			

Linear/Quantum Data Types

Characteristic Property	1. their cartesian product blends into the co-product:		
Symbol			
Formula (for $B : \text{FinType}$)			
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Linear Logic			
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Linear/Quantum Data Types

Characteristic Property	1. their cartesian product blends into the co-product:		
Symbol	\oplus direct sum		
Formula (for $B : \text{FinType}$)			
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Physics Meaning			

Linear/Quantum Data Types

Characteristic Property	1. their cartesian product blends into the co-product:	2. a tensor product appears & distributes over direct sum	
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AlgTop Jargon	biproduct, stability, ambidexterity	Frobenius reciprocity	mapping spectrum
Linear Logic			
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Linear Logic	additive disjunction	multiplicative conjunction	linear implication
Physics Meaning	superselection sectors / quantum parallelism	compound quantum systems / quantum entanglement	QRAM systems

Linear/Quantum Data Types

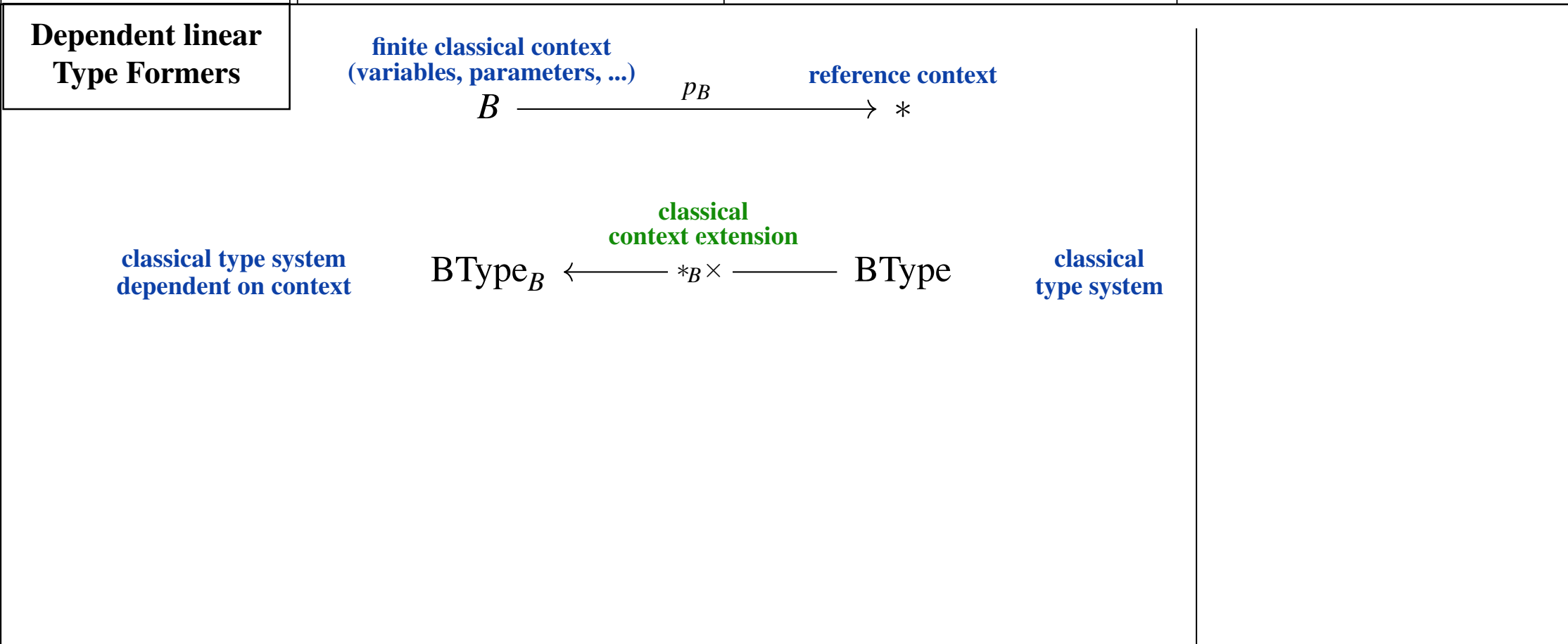
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Dependent linear Type Formers	$ \begin{array}{ccc} \text{finite classical context} & & \text{reference context} \\ \text{(variables, parameters, ...)} & & \\ B & \xrightarrow{p_B} & * \end{array} $		
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		<small>classical type system</small>	

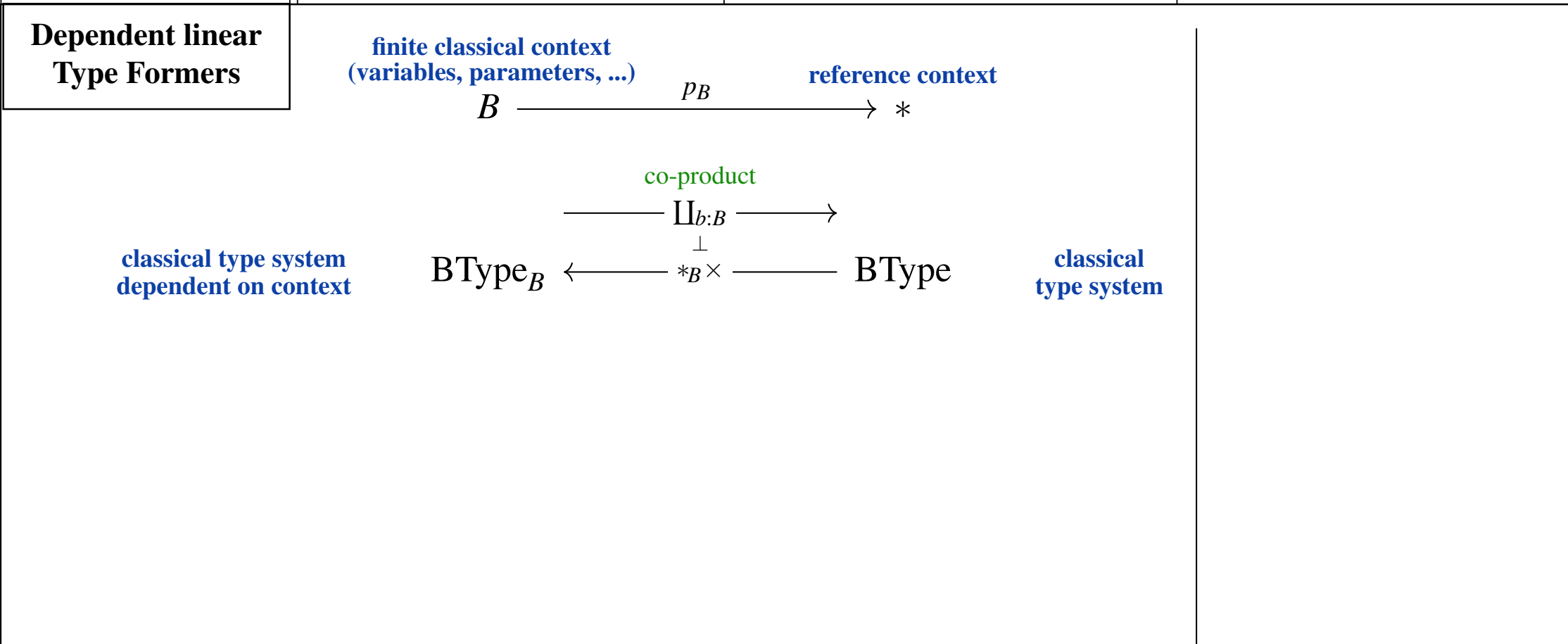
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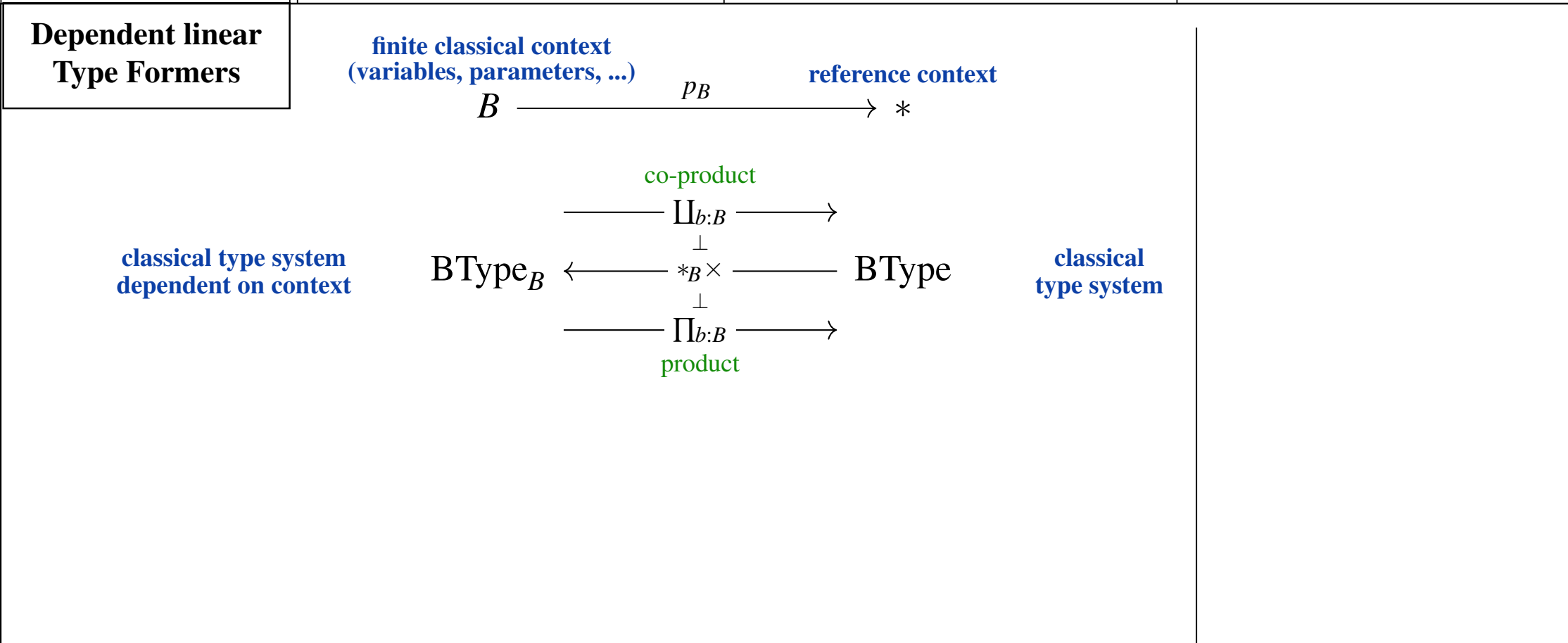
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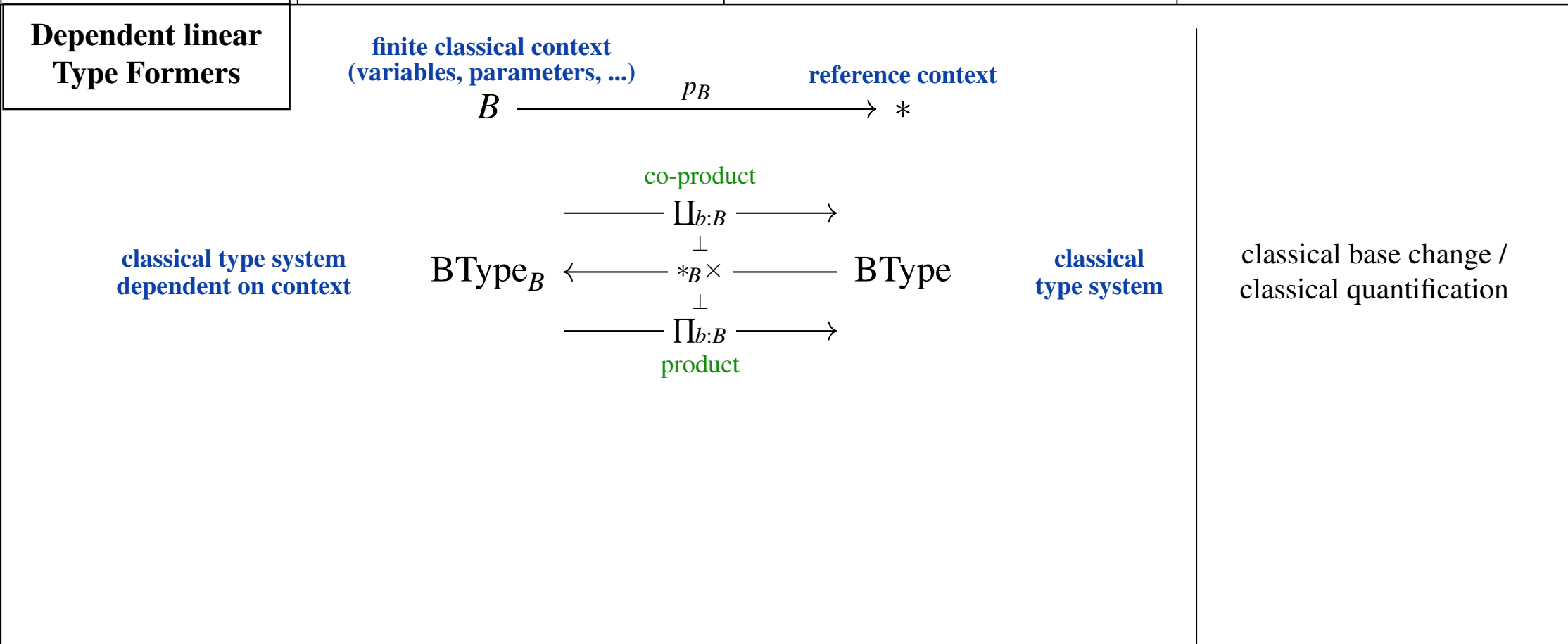
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Dependent linear Type Formers	<p>finite classical context (variables, parameters, ...)</p> $B \xrightarrow{p_B} *$ <p>reference context</p>		
classical type system dependent on context	$\text{BType}_B \xleftarrow{\quad} * \times \xrightarrow{\quad} \text{BType}$ <p style="text-align: center;"> \perp \perp </p>	classical type system	classical base change / classical quantification
	$\xrightarrow{\quad} \coprod_{b:B} \xrightarrow{\quad}$ <p style="text-align: center;">co-product</p>		
	$\xrightarrow{\quad} \prod_{b:B} \xrightarrow{\quad}$ <p style="text-align: center;">product</p>		

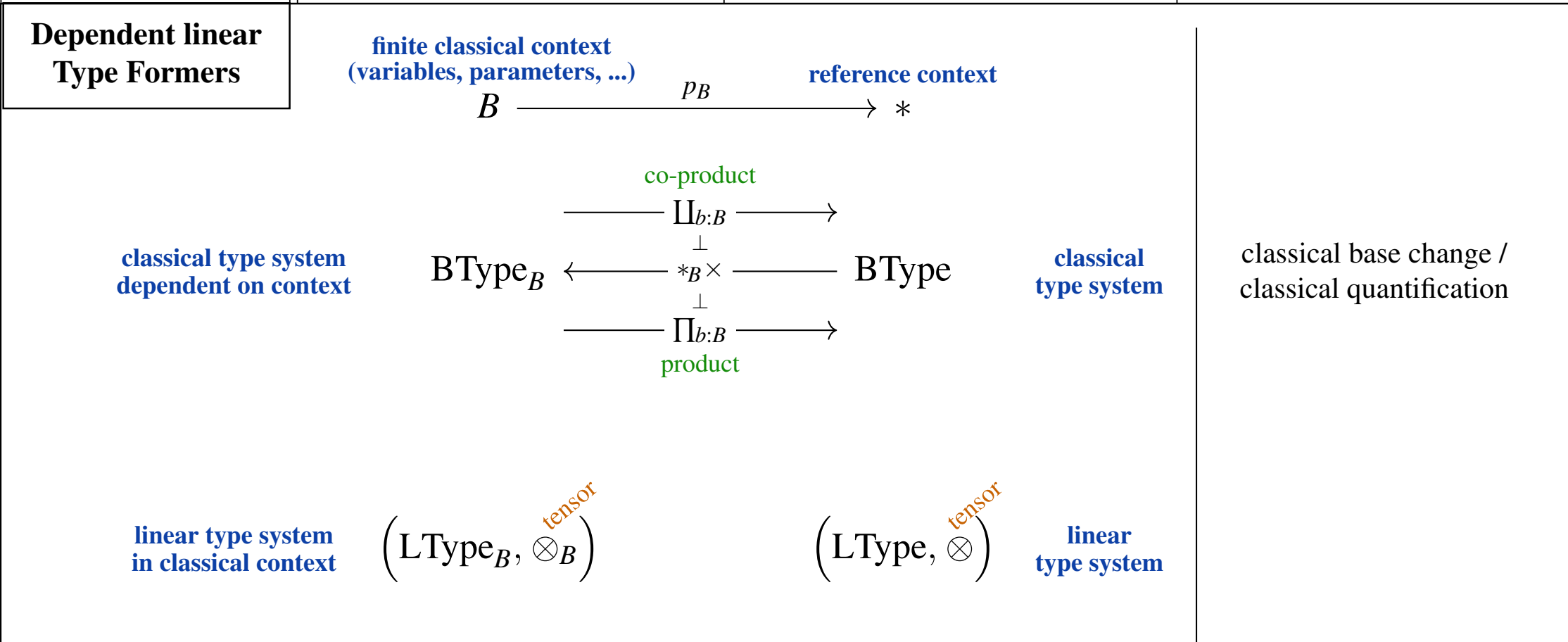
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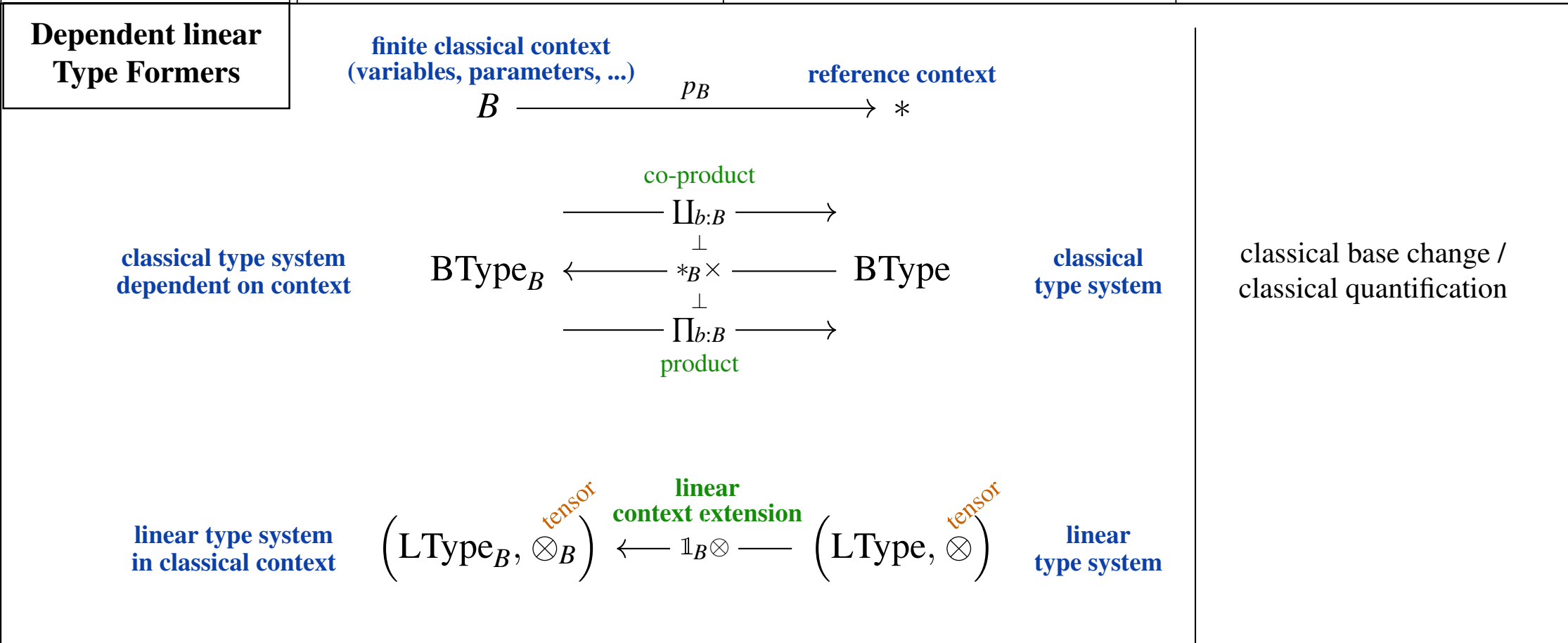
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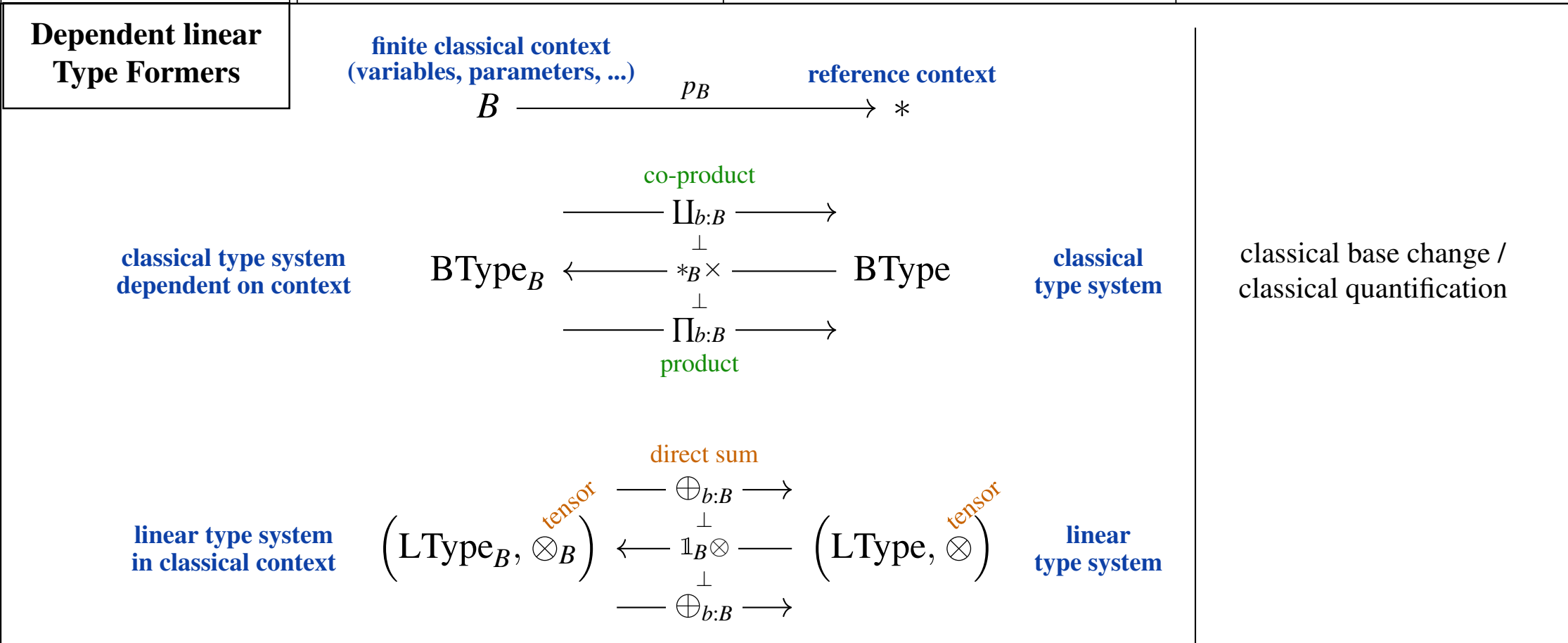
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Dependent linear Type Formers	<p>finite classical context (variables, parameters, ...)</p> $B \xrightarrow{p_B} *$ <p>reference context</p>		
classical type system dependent on context	$\text{BType}_B \xleftarrow{\quad} *_B \times \xrightarrow{\quad} \text{BType}$ <p style="text-align: center;"> \perp \perp </p> $\xrightarrow{\quad} \prod_{b:B} \xrightarrow{\quad}$ <p style="text-align: center;"> <small>co-product</small> <small>product</small> </p>	classical type system	classical base change / classical quantification
linear type system in classical context	$\left(\text{LType}_B, \otimes_B \right) \xleftarrow{\quad} \mathbb{1}_B \otimes \xrightarrow{\quad} \left(\text{LType}, \otimes \right)$ <p style="text-align: center;"> <small>direct sum</small> <small>tensor</small> \perp <small>tensor</small> </p> $\xrightarrow{\quad} \bigoplus_{b:B} \xrightarrow{\quad}$	linear type system	quantum base change / Motivic Yoga

Linear/Quantum Data Types

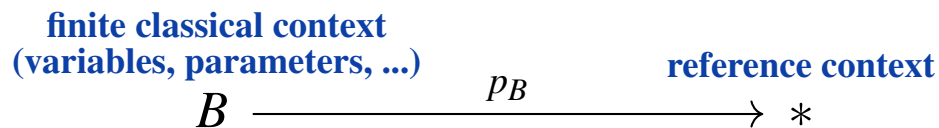
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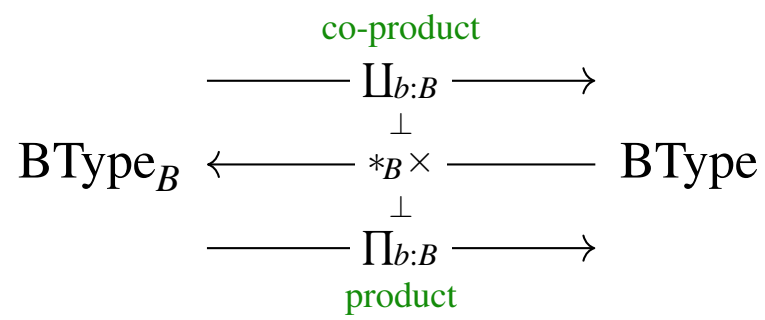
Characteristic Property	1. their cartesian product blends into the co-product:	2. a tensor product appears & distributes over direct sum	3. a linear function type appears adjoint to tensor
Symbol	\oplus direct sum	\otimes tensor product	\multimap linear function type
Formula (for $B : \text{FinType}$)	cart. product co-product $\prod_B \mathcal{H}_b \simeq \bigoplus_B \mathcal{H}_b \simeq \coprod_B \mathcal{H}_b$ direct sum	$\mathcal{V} \otimes \left(\bigoplus_{b:B} \mathcal{H}_b \right) \simeq \bigoplus_{b:B} (\mathcal{V} \otimes \mathcal{H}_b)$	$(\mathcal{V} \otimes \mathcal{H}) \multimap \mathcal{K}$ $\simeq \mathcal{V} \multimap (\mathcal{H} \multimap \mathcal{K})$

Dependent linear Type Formers



all available
in dLHoTT

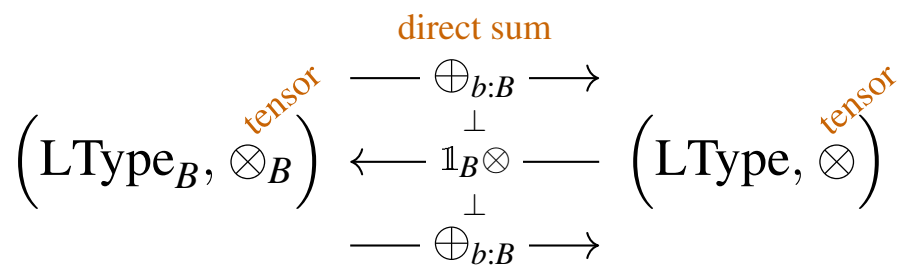
classical type system
dependent on context



classical
type system

classical base change /
classical quantification

linear type system
in classical context



linear
type system

quantum base change
/ Motivic Yoga

Quantum Effects

Recall: **Monadic computational effects.**

A monad $\mathcal{E}(-)$ on a data type system encodes *computational effects*:

effectful program

$$D_1 \xrightarrow{\text{prog}_{12}} \mathcal{E}(D_2)$$

**output data of nominal type D_2
causing effects of type $\mathcal{E}(-)$**

Recall: **Monadic computational effects.**

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first program

$$D_1 \xrightarrow{\text{prog}_{12}} \mathcal{E}(D_2)$$

**output data of nominal type D_2
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second program

$$D_2 \xrightarrow{\text{prog}_{23}} \mathcal{E}(D_3)$$

**input data of type D_2
causing effects of type $\mathcal{E}(-)$**

Recall: Monadic computational effects.

A monad $\mathcal{E}(-)$ on a data type system encodes *computational effects*:

first program

$$D_1 \xrightarrow{\text{prog}_{12}} \mathcal{E}(D_2)$$

output data of nominal type D_2
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second program

$$D_2 \xrightarrow{\text{prog}_{23}} \mathcal{E}(D_3)$$


input data of type D_2
causing effects of type $\mathcal{E}(-)$

$$D_1 \xrightarrow{\text{prog}_{12}} \mathcal{E}(D_2)$$

$$\mathcal{E}(D_2) \xrightarrow{\text{bind}^{\mathcal{E}} \text{prog}_{23}} \mathcal{E}(D_3)$$

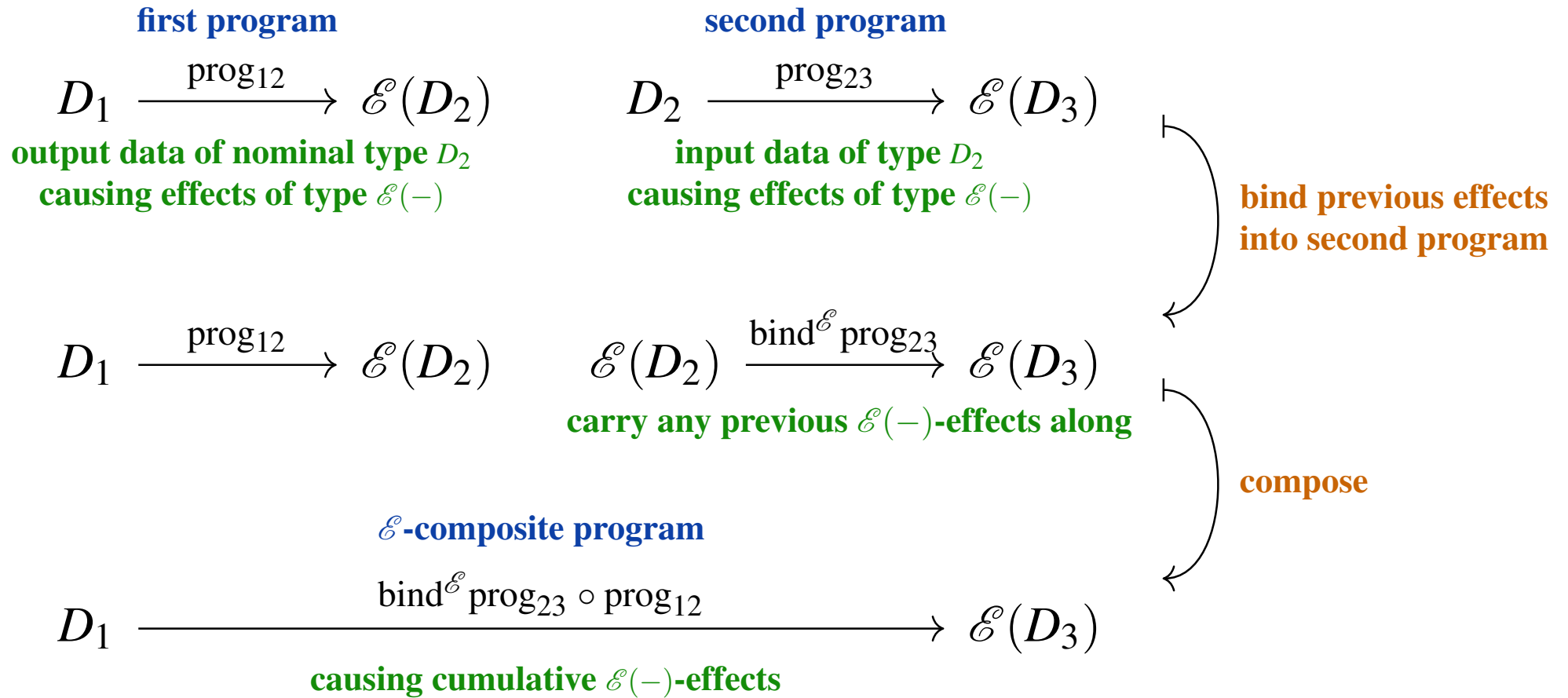
carry any previous $\mathcal{E}(-)$ -effects along

bind previous effects
into second program



Recall: Monadic computational effects.

A monad $\mathcal{E}(-)$ on a data type system encodes *computational effects*:



Recall: **Monadic effect handlers.**

$D_1 \xrightarrow{\text{prog}_{12}} D_2$ **data type to absorb \mathcal{E} -effects**
in-effectful program

Recall: Monadic effect handlers.


$$D_1 \xrightarrow{\text{prog}_{12}} D_2$$

in-effectful program

$$\mathcal{E}(D_1) \xrightarrow{\text{hdl}_{D_2}^{\mathcal{E}} \text{prog}_{12}} D_2$$

**in-effectful program
handling effects of type $\mathcal{E}(-)$**

**incorporate handling
of $\mathcal{E}(-)$ -effects**



Recall: Monadic effect handlers.

$$D_1 \xrightarrow{\text{prog}_{12}} D_2$$

in-effectful program

$$\mathcal{E}(D_1) \xrightarrow{\text{hdl}_{D_2}^{\mathcal{E}} \text{prog}_{12}} D_2$$

in-effectful program
handling effects of type $\mathcal{E}(-)$

incorporate handling
of $\mathcal{E}(-)$ -effects

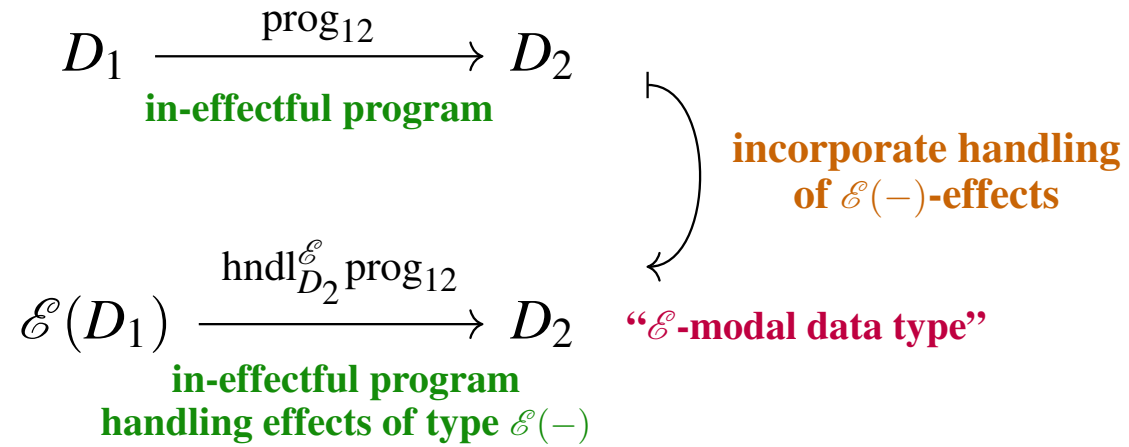
$$D_1 \xrightarrow{\text{ret}_{D_1}^{\mathcal{E}}} \mathcal{E}(D_1) \xrightarrow{\text{hdl}_{D_2}^{\mathcal{E}} \text{prog}_{12}} D_2$$

produce trivial effect **handle effects running program**

prog₁₂
no effect

consistency conditions

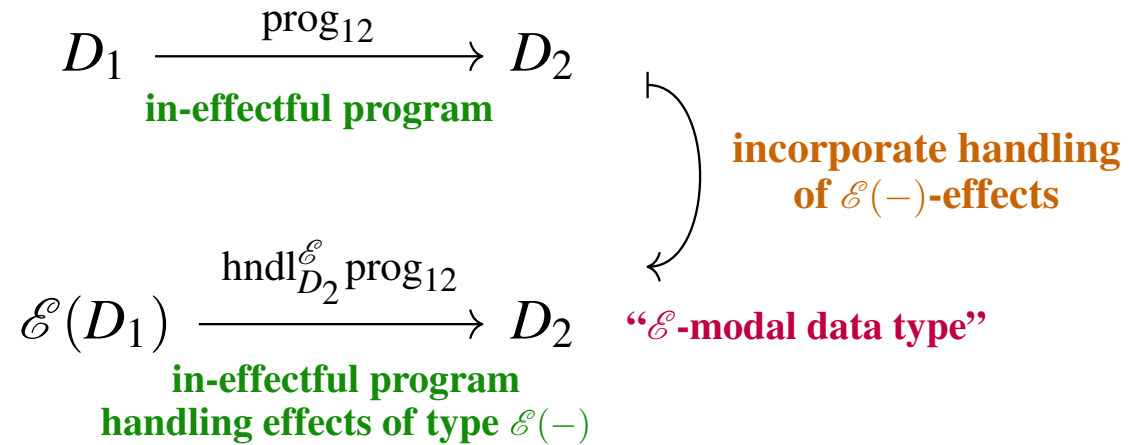
Recall: Data type system of Monadic effect handlers.



Monadicity:

\mathcal{E} -modales in Type
("EM-category") $\text{Type}^{\mathcal{E}}$

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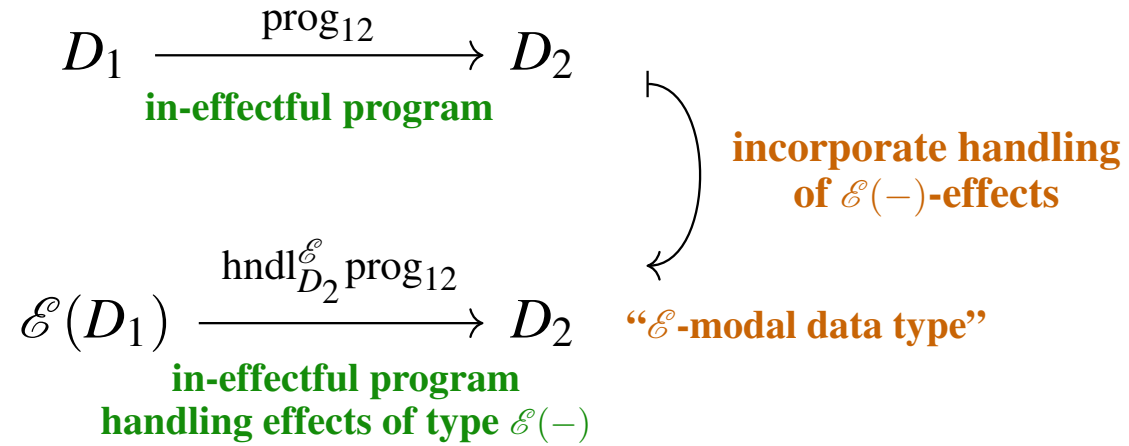


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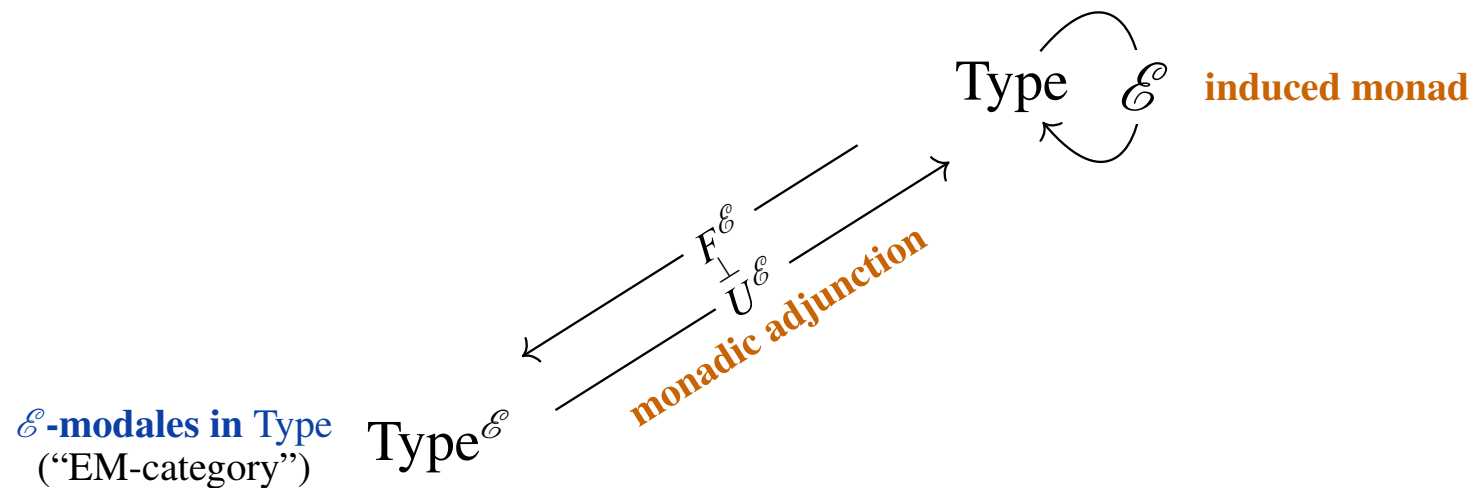


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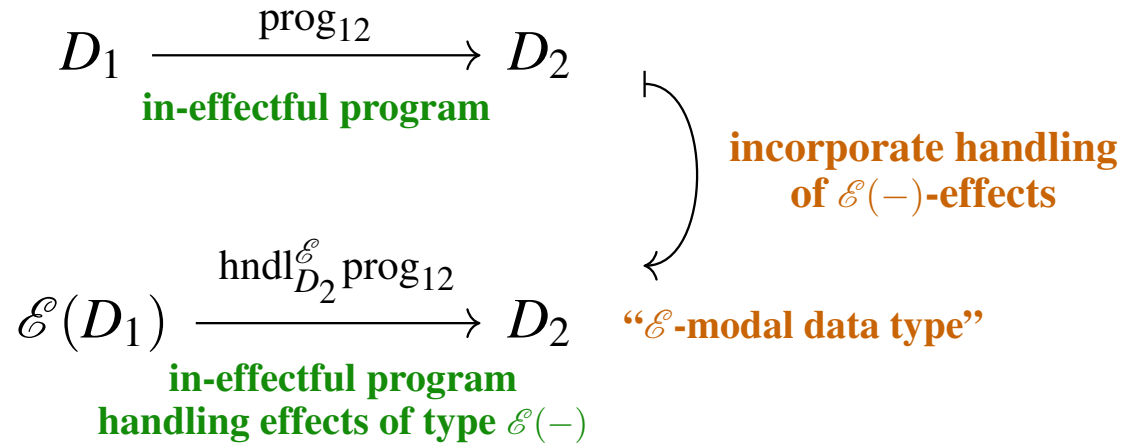
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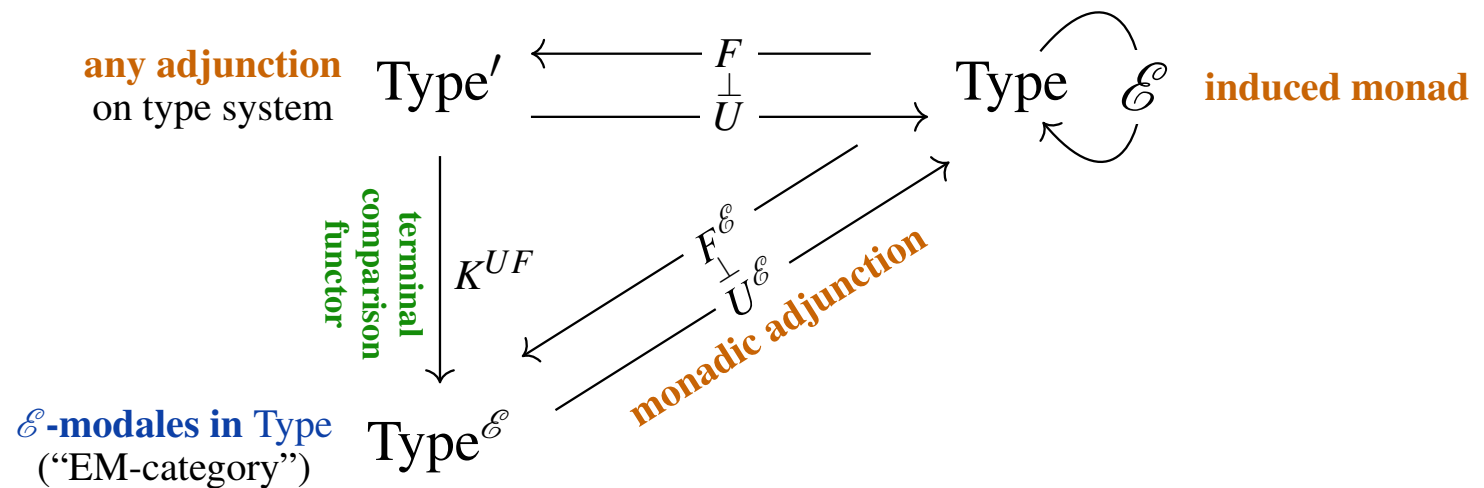
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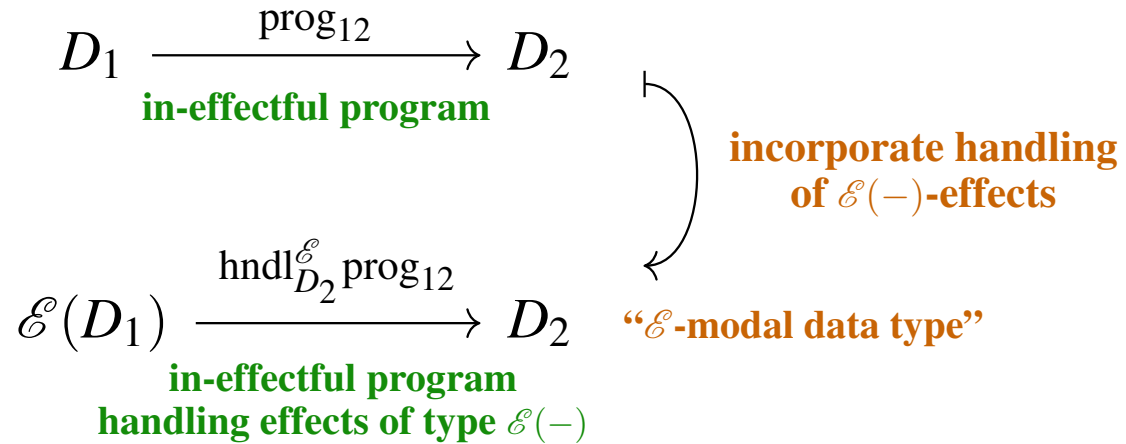
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Recall: Data type system of Monadic effect handlers.



Monadicity:

free \mathcal{E} -modales in Type
 (“Kleisli category”)

Type $_{\mathcal{E}}$

initial
comparison
functor

K_{UF}

any adjunction
on type system

Type'

F
 \perp
 U

Type \mathcal{E}

induced monad

terminal
comparison
functor

K_{UF}

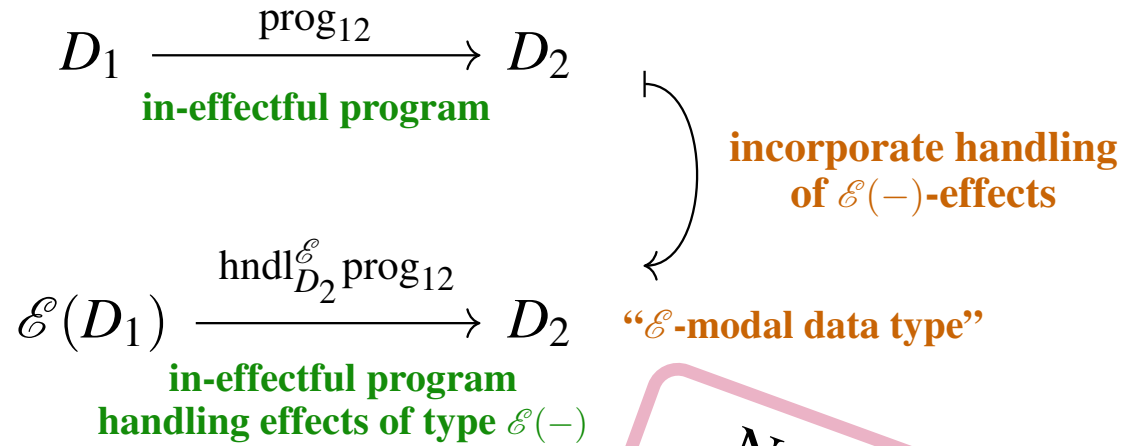
\mathcal{E} -modales in Type
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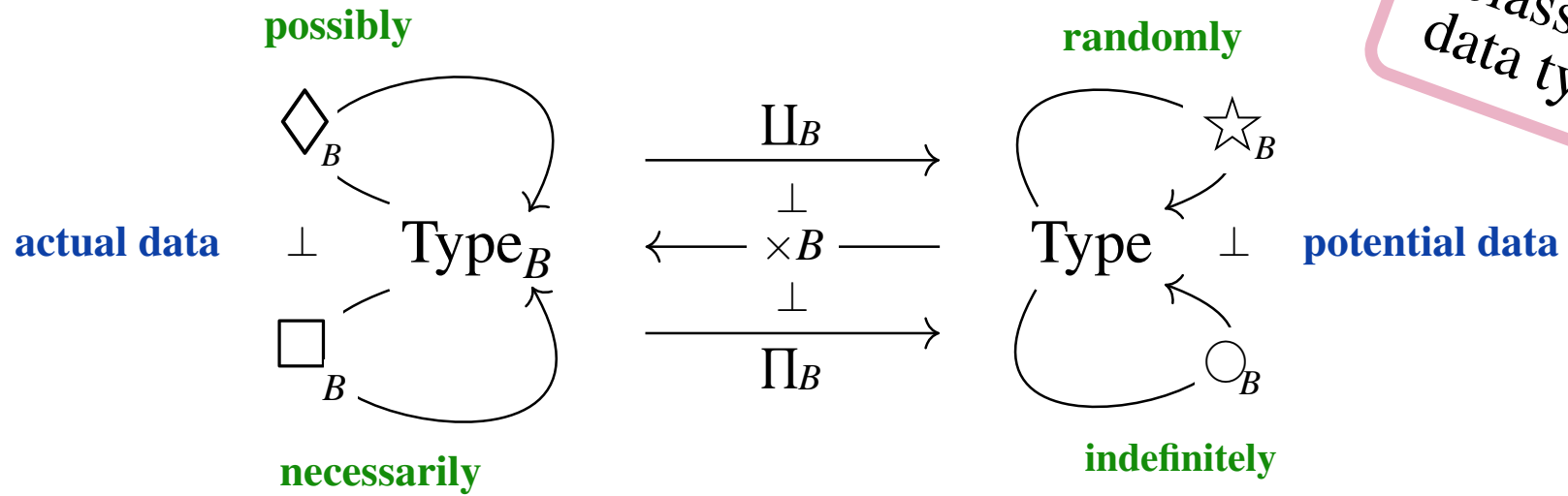
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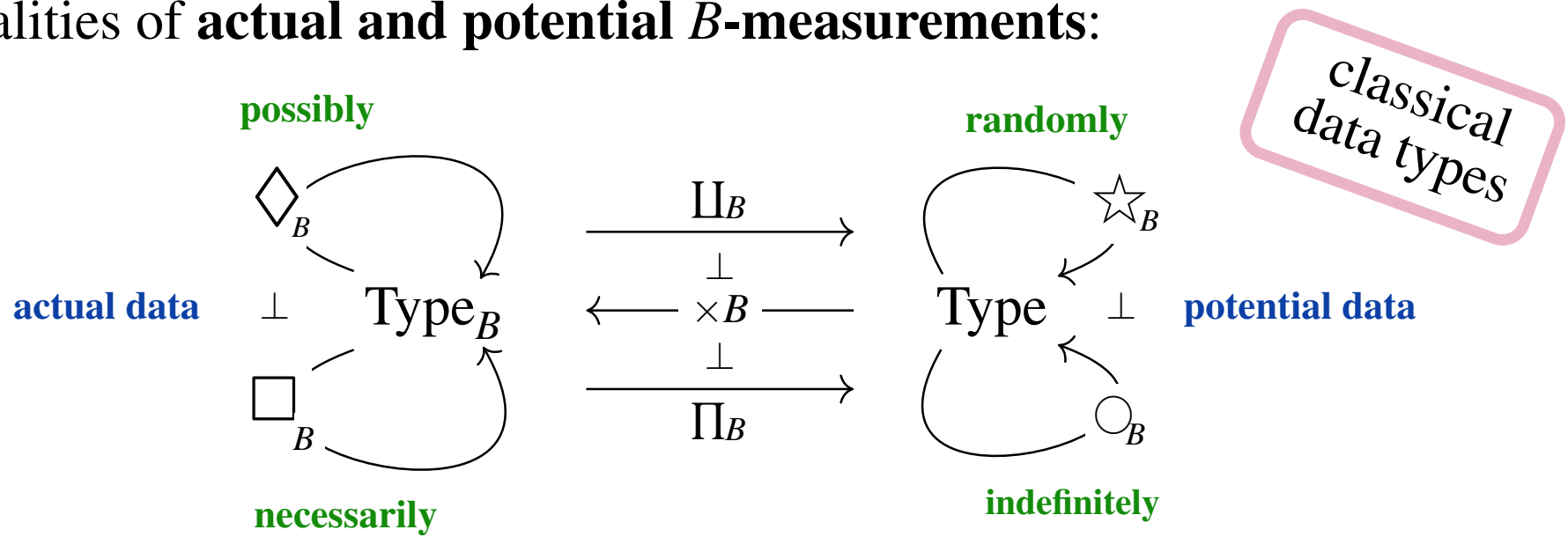
$F^{\mathcal{E}}$
 \perp
 $U^{\mathcal{E}}$
monadic adjunction

Now just to work this out
 for the effects induced by
 dependent data type formers
 in dLHoTT

Given $B : B\text{Type}$ of possible measurement outcomes (“possible worlds”) **the monadic effects of B -dependent data type formers** constitute modalities of **actual and potential B -measurements**:



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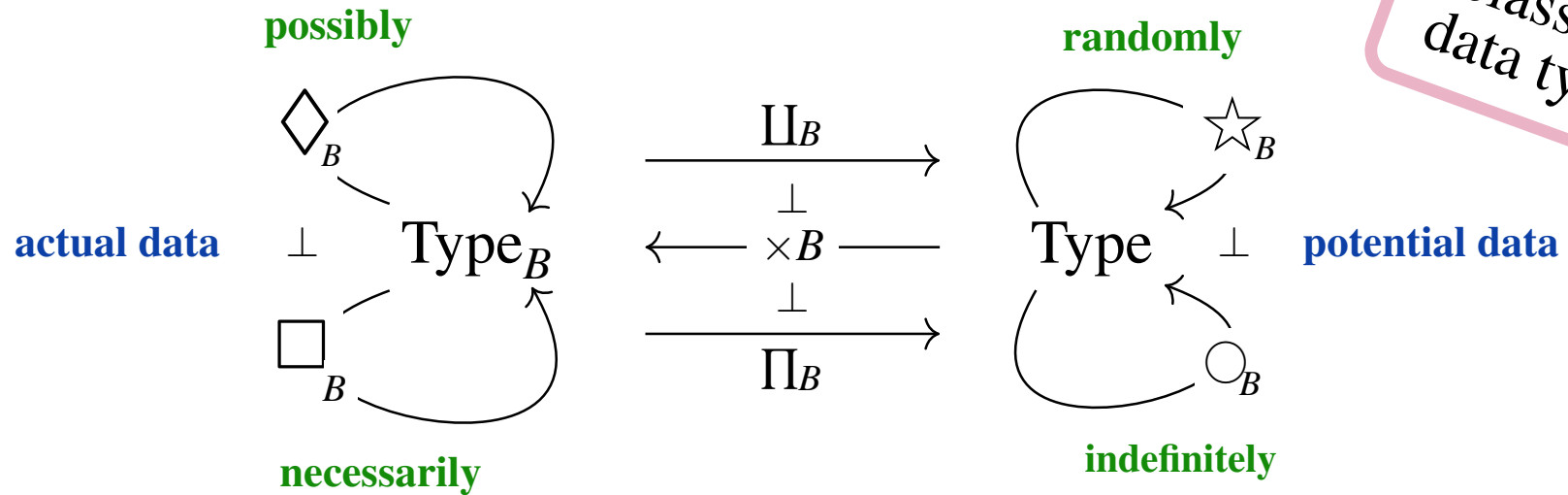


necessarily P .

$\square_B P.$

$b : B \vdash \prod_{b':B} P_{b'}$

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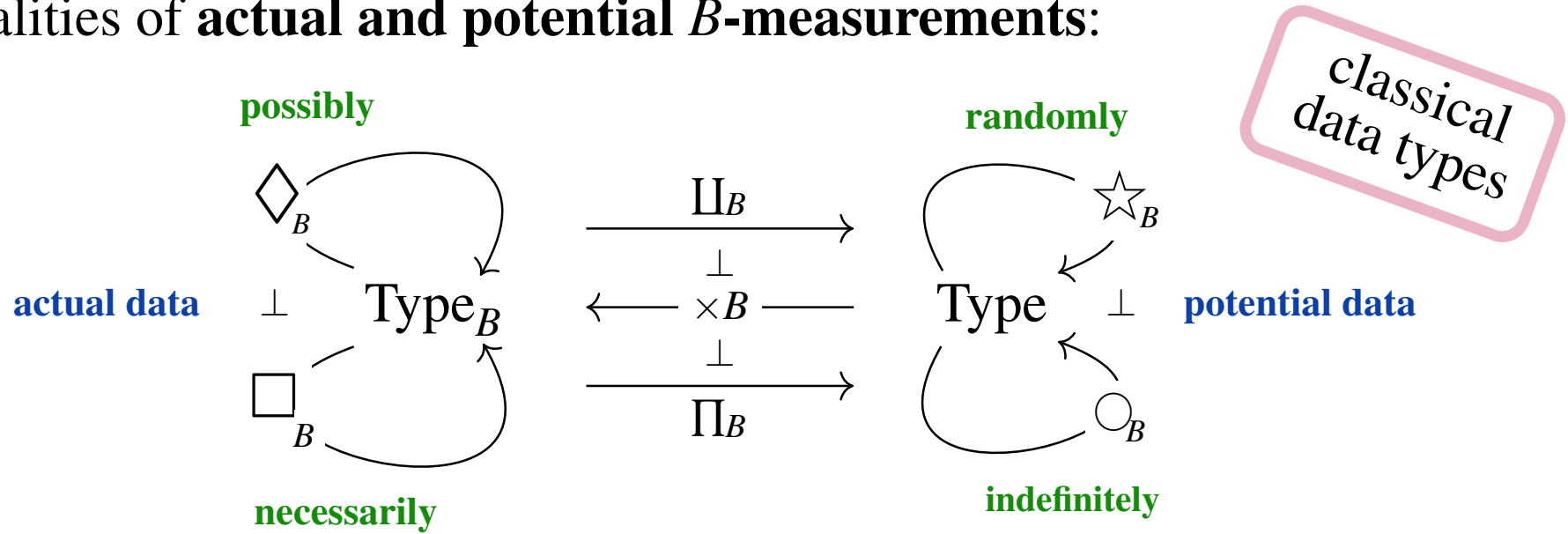


necessarily P_\bullet entails actually P_\bullet .

$$\square_B P_\bullet \longrightarrow \varepsilon_{P_\bullet}^{\square_B} \longrightarrow P_\bullet$$

$$b : B \vdash \prod_{b' : B} P_{b'} \xrightarrow{(P_{b'})_{b' : B} \mapsto P_b} P_b$$

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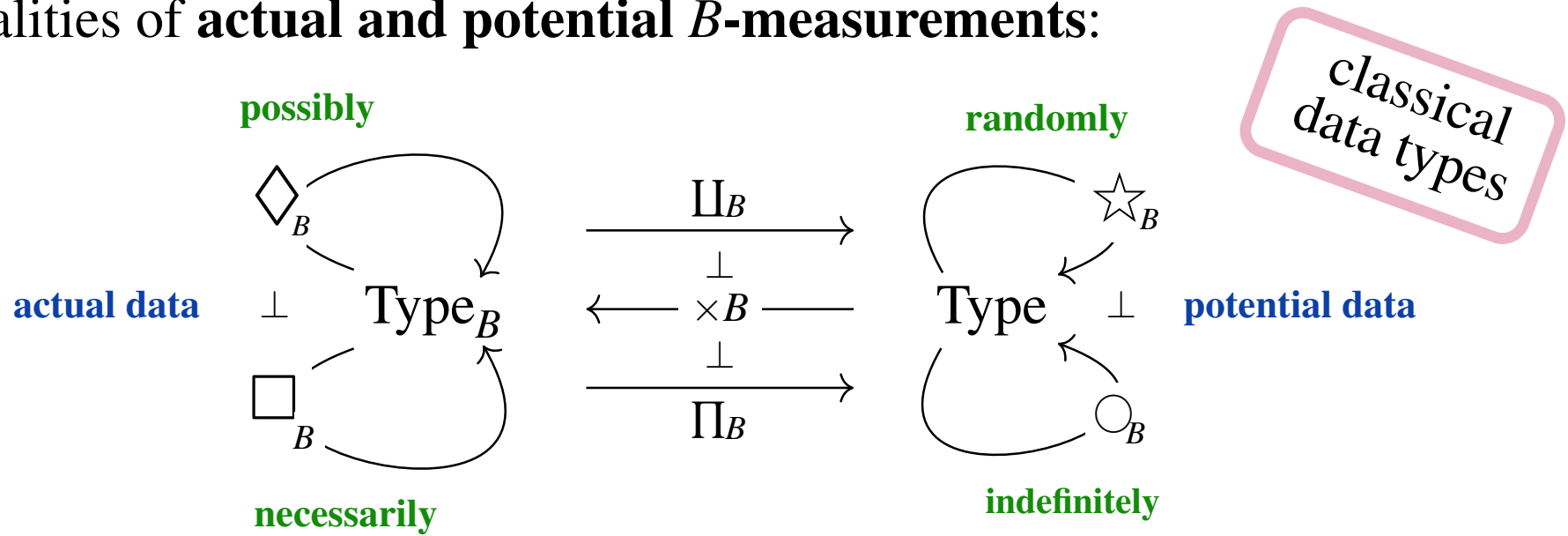


necessarily P_\bullet entails actually P_\bullet entails possibly P_\bullet

$$\square_B P_\bullet \xrightarrow{\varepsilon_{P_\bullet}^{\square_B}} P_\bullet \xrightarrow{\eta_{P_\bullet}^{\diamond_B}} \diamond_B P_\bullet$$

$$b : B \vdash \prod_{b':B} P_{b'} \xrightarrow{(p_{b'})_{b':B} \mapsto p_b} P_b \xrightarrow{p_b \mapsto (p_b)_b} \coprod_{b':B} P_{b'}$$

Given $B : BType$ of possible measurement outcomes (“possible worlds”) the monadic effects of B -dependent data type formers constitute modalities of actual and potential B -measurements:



necessarily P_\bullet entails actually P_\bullet entails possibly P_\bullet

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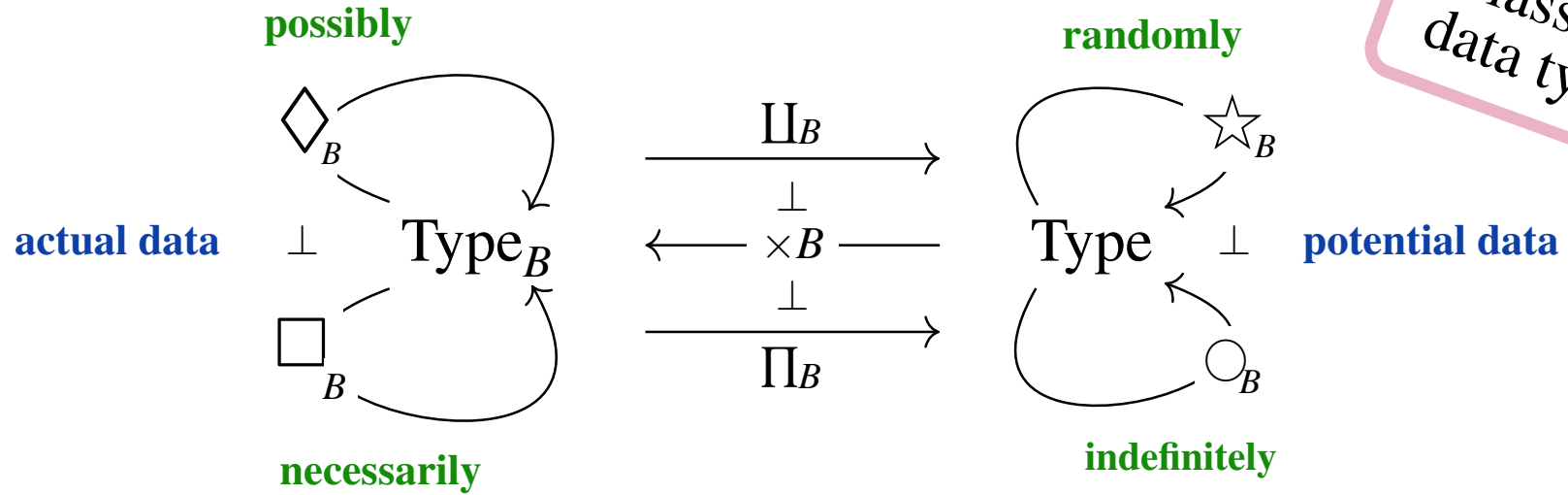
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randomly P

$$\star_B P$$

$$\coprod_{b:B} P$$

Given $B : B\text{Type}$ of possible measurement outcomes (“possible worlds”) the monadic effects of B -dependent data type formers constitute modalities of actual and potential B -measurements:



necessarily P_\bullet entails actually P_\bullet entails possibly P_\bullet

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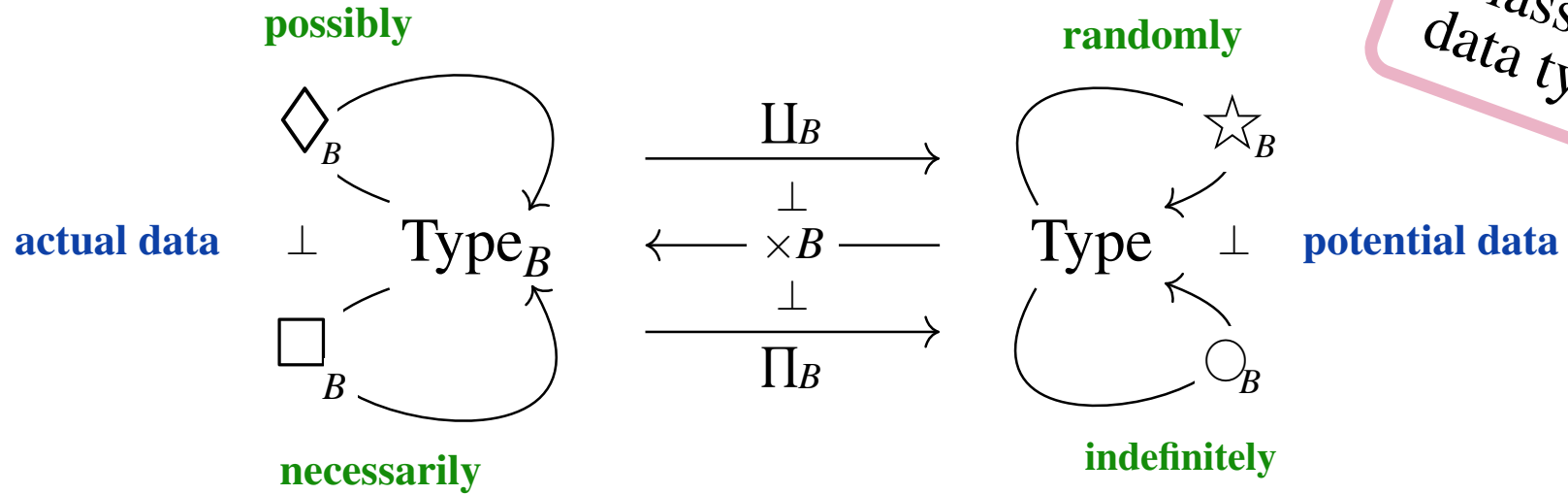
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randomly P entails potentially P

$$\star_B P \xrightarrow{\varepsilon_P^{\star_B}} P$$

$$\coprod_{b:B} P \xrightarrow{(p)_b \mapsto p} P$$

Given $B : B\text{Type}$ of possible measurement outcomes (“possible worlds”) the monadic effects of B -dependent data type formers constitute modalities of actual and potential B -measurements:



necessarily P_\bullet entails actually P_\bullet entails possibly P_\bullet

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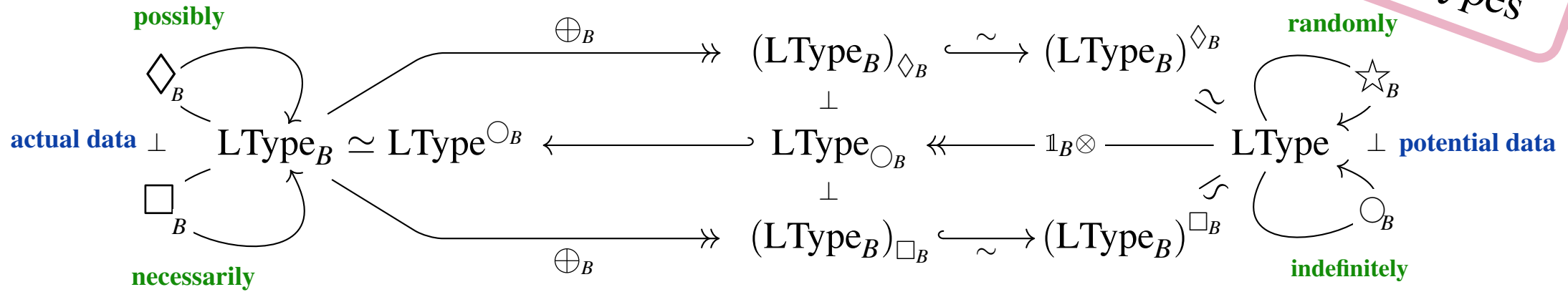
randomly P entails potentially P entails indefinitely P

$$\star_B P \xrightarrow{\varepsilon_P^{\star_B}} P \xrightarrow{\eta_P^{\circ_B}} \circ_B P$$

$$\coprod_{b:B} P \xrightarrow{(p)_b \mapsto p} P \xrightarrow{p \mapsto (p)_{b:B}} \prod_{b:B} P$$

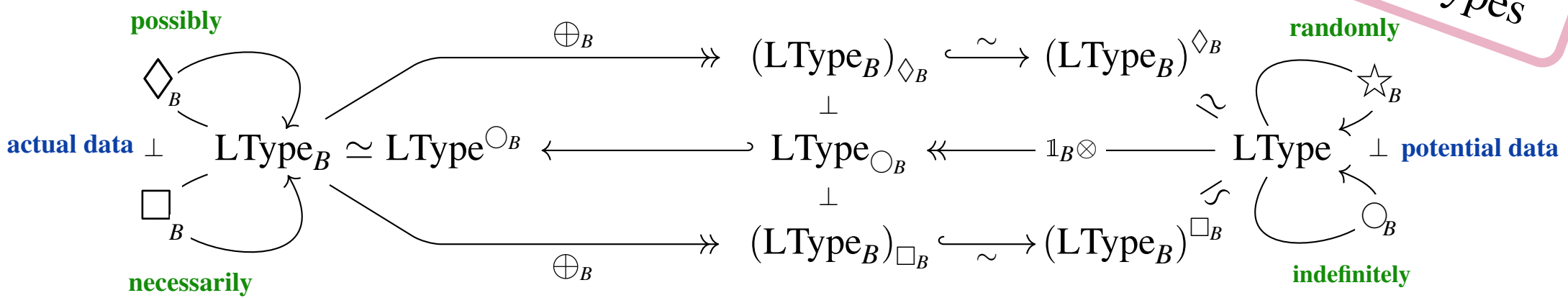
Given $B : B\text{Type}$ of possible measurement outcomes (“possible worlds”) the monadic effects of B -dependent **linear** data type formers constitute modalities of **actual** and **potential quantum** B -measurements.

quantum data types



Given $B : \text{BType}$ of possible measurement outcomes (“possible worlds”) the monadic effects of B -dependent **linear** data type formers constitute modalities of **actual** and **potential quantum** B -measurements.

quantum data types



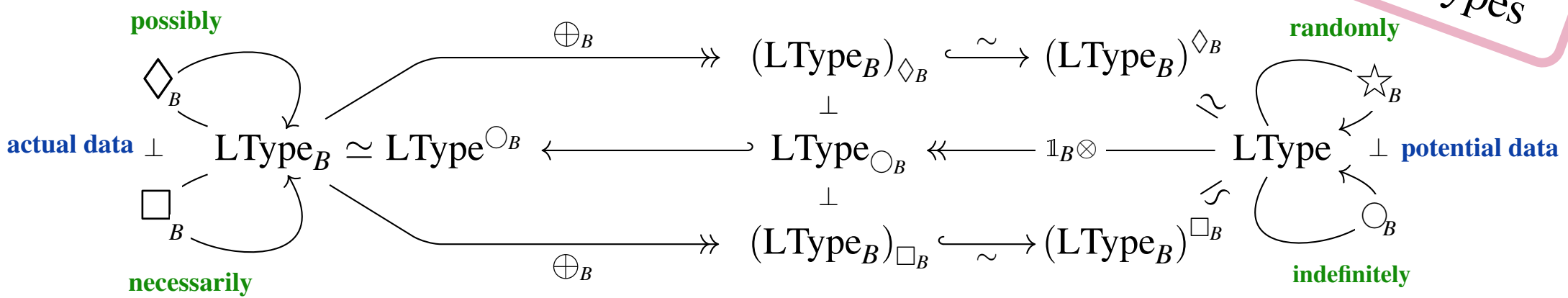
necessarily $\mathcal{H} \bullet$
 $\square_B \mathcal{H} \bullet$

Given... obtain...
 $b : B \vdash \mathcal{H}$
 measurement result

where $\mathcal{H} := \bigoplus_{b' : B} \mathcal{H}_{b'}$

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quantum data types



necessarily \mathcal{H}_\bullet entails actually \mathcal{H}_\bullet

$$\square_B \mathcal{H}_\bullet \xrightarrow{\varepsilon_{\mathcal{H}_\bullet}^{\square B}} \mathcal{H}_\bullet$$

Given... obtain...

$b : B \vdash$ $\mathcal{H} \xrightarrow[\text{measurement collapse}]{\Sigma_{b'} |\psi_{b'}\rangle \mapsto |\psi_b\rangle} \mathcal{H}_b$

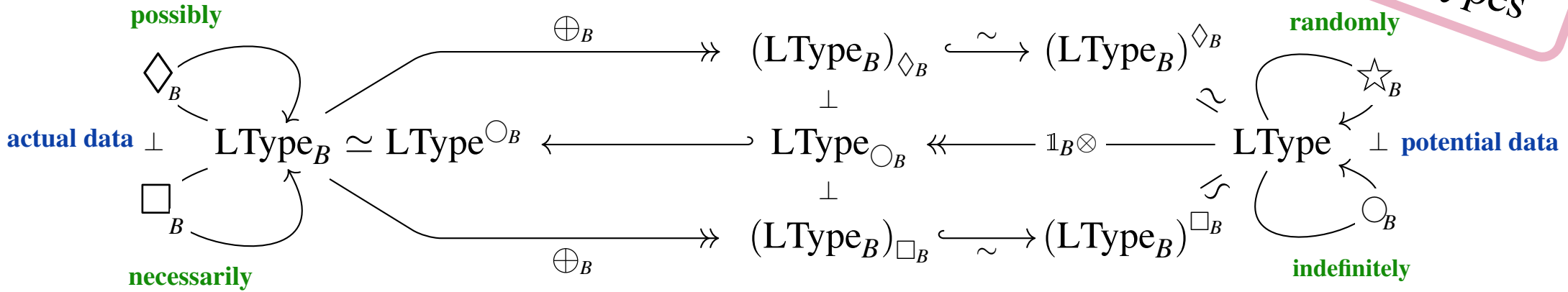
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quantum data types



$$\text{necessarily } \mathcal{H}_\bullet \quad \text{entails} \quad \text{actually } \mathcal{H}_\bullet \quad \text{entails} \quad \text{possibly } \mathcal{H}_\bullet$$

$$\square_B \mathcal{H}_\bullet \xrightarrow{\varepsilon_{\mathcal{H}_\bullet}^{\square B}} \mathcal{H}_\bullet \xrightarrow{\eta_{\mathcal{H}_\bullet}^{\diamond B}} \diamond_B \mathcal{H}_\bullet$$

Given... $b : B$
obtain... \vdash
measurement result

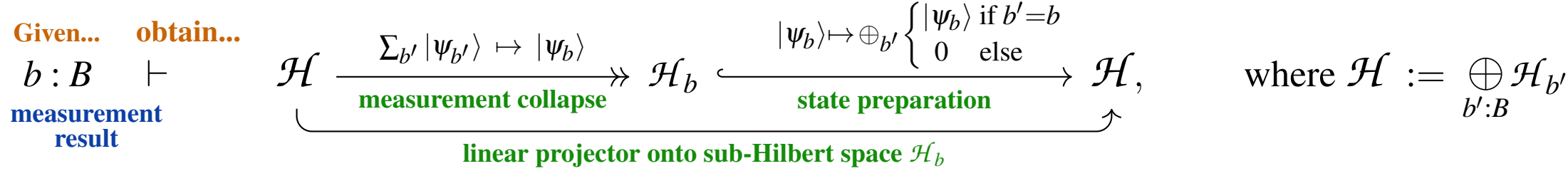
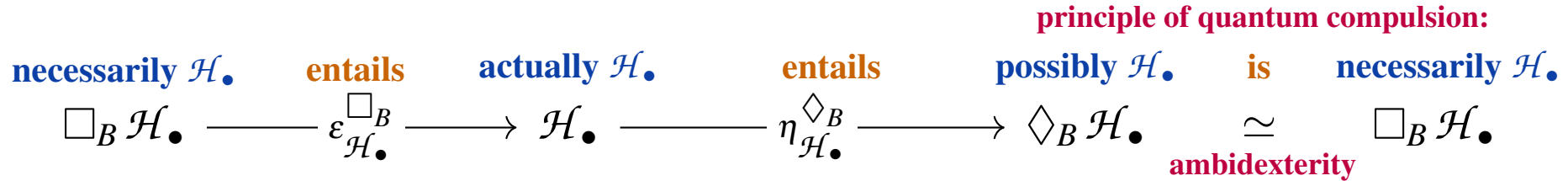
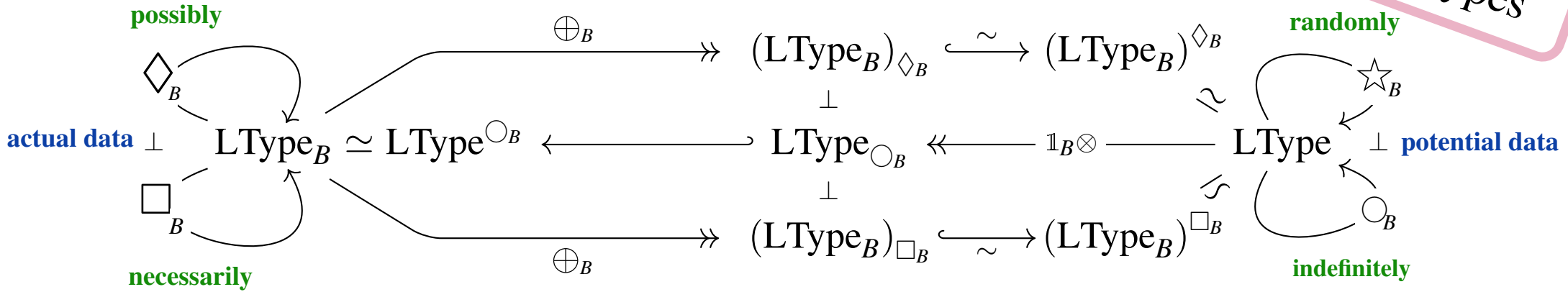
$$\mathcal{H} \xrightarrow[\text{measurement collapse}]{\sum_{b'} |\psi_{b'}\rangle \mapsto |\psi_b\rangle} \mathcal{H}_b \xrightarrow[\text{state preparation}]{|\psi_b\rangle \mapsto \oplus_{b'} \begin{cases} |\psi_b\rangle & \text{if } b'=b \\ 0 & \text{else} \end{cases}} \mathcal{H}$$

linear projector onto sub-Hilbert space \mathcal{H}_b

where $\mathcal{H} := \bigoplus_{b':B} \mathcal{H}_{b'}$

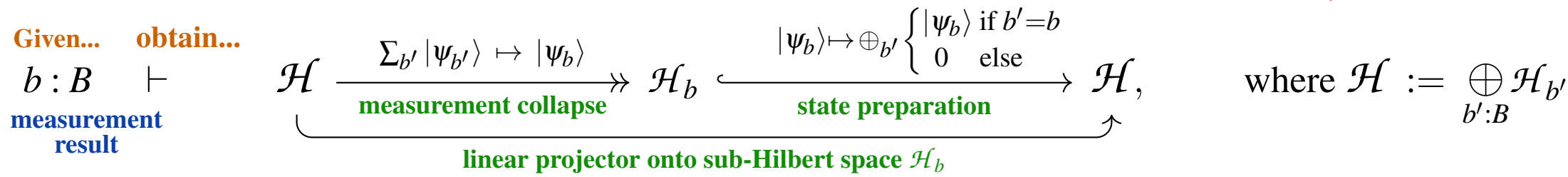
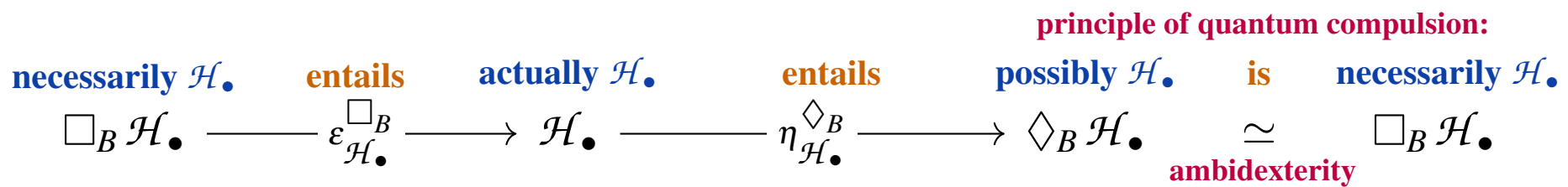
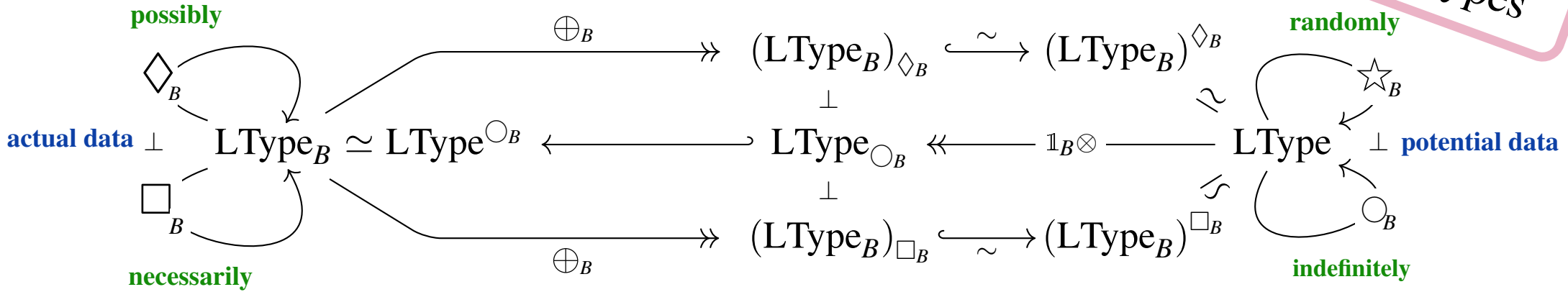
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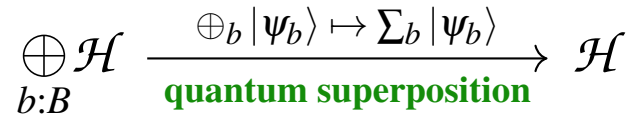
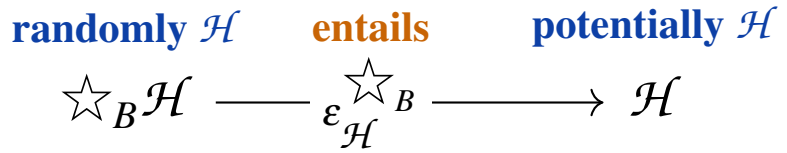
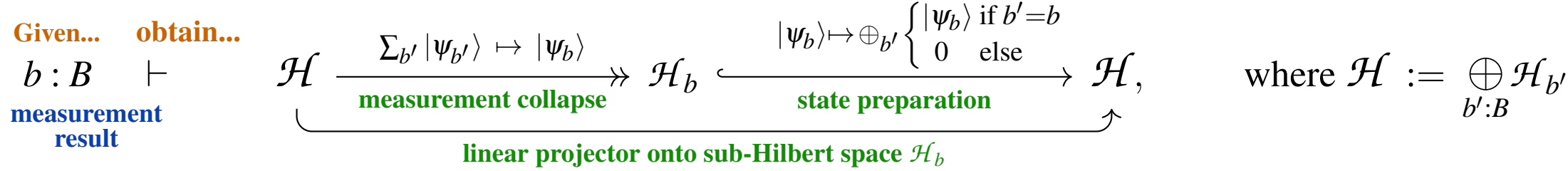
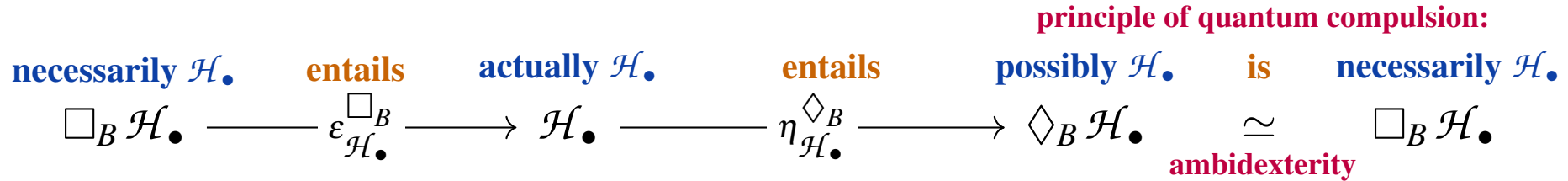
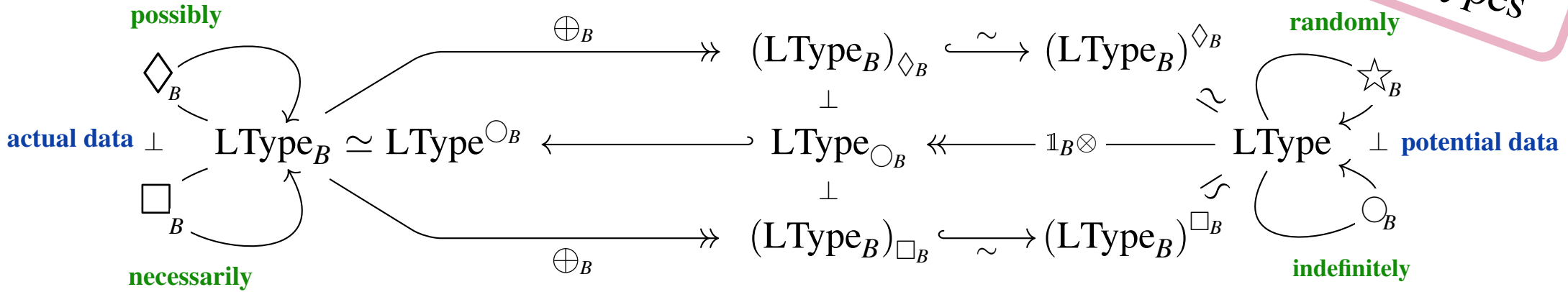
randomly \mathcal{H}

$\star_B \mathcal{H}$

$\bigoplus_{b:B} \mathcal{H}$

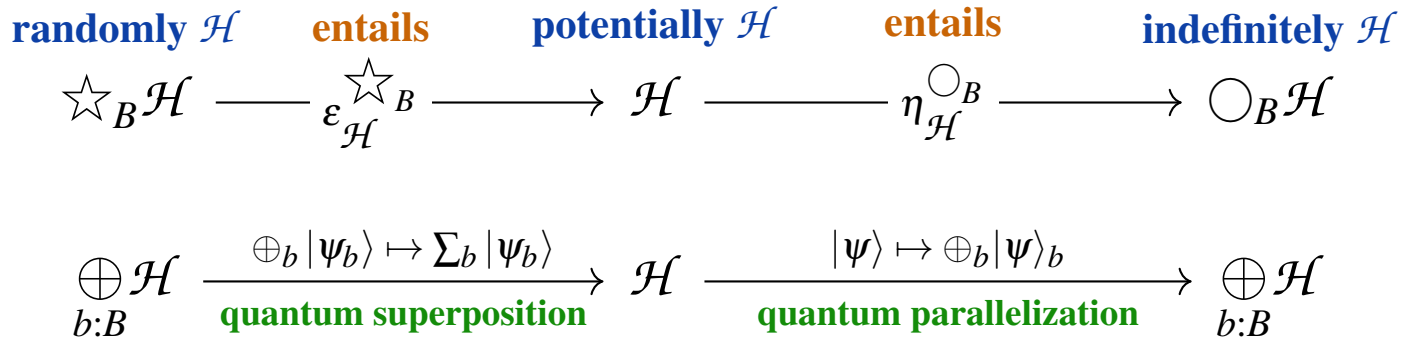
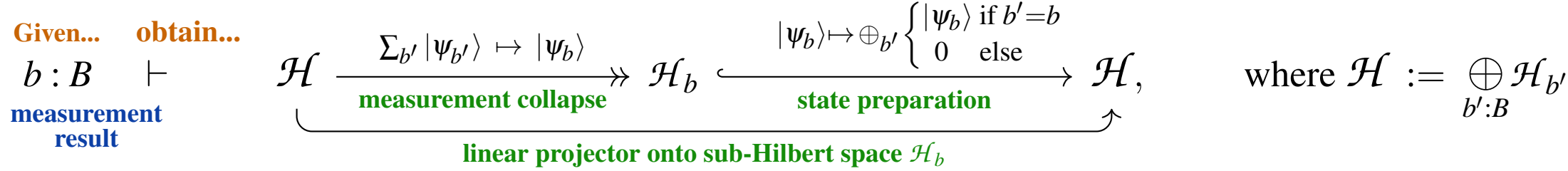
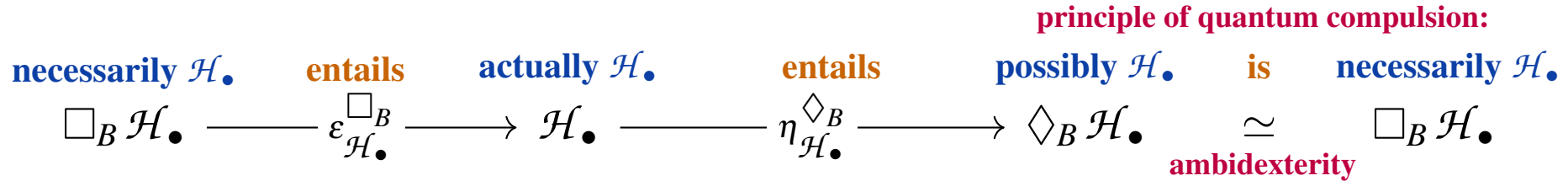
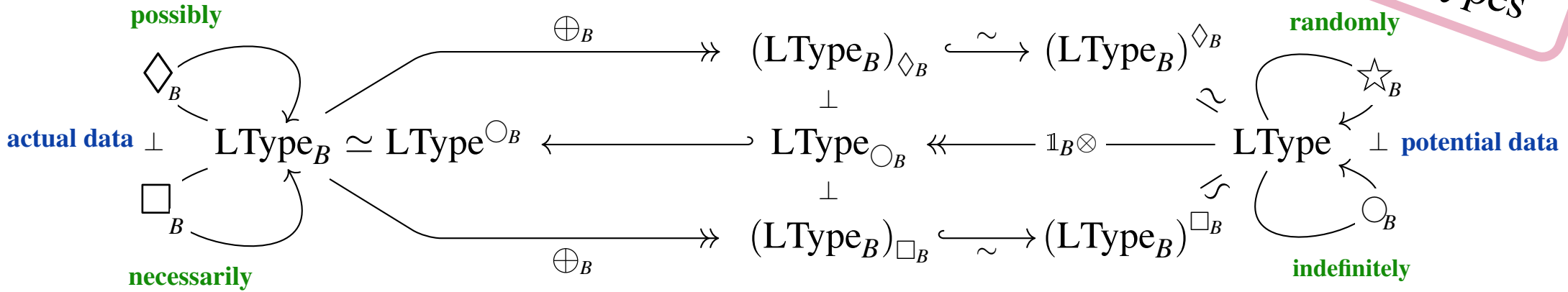
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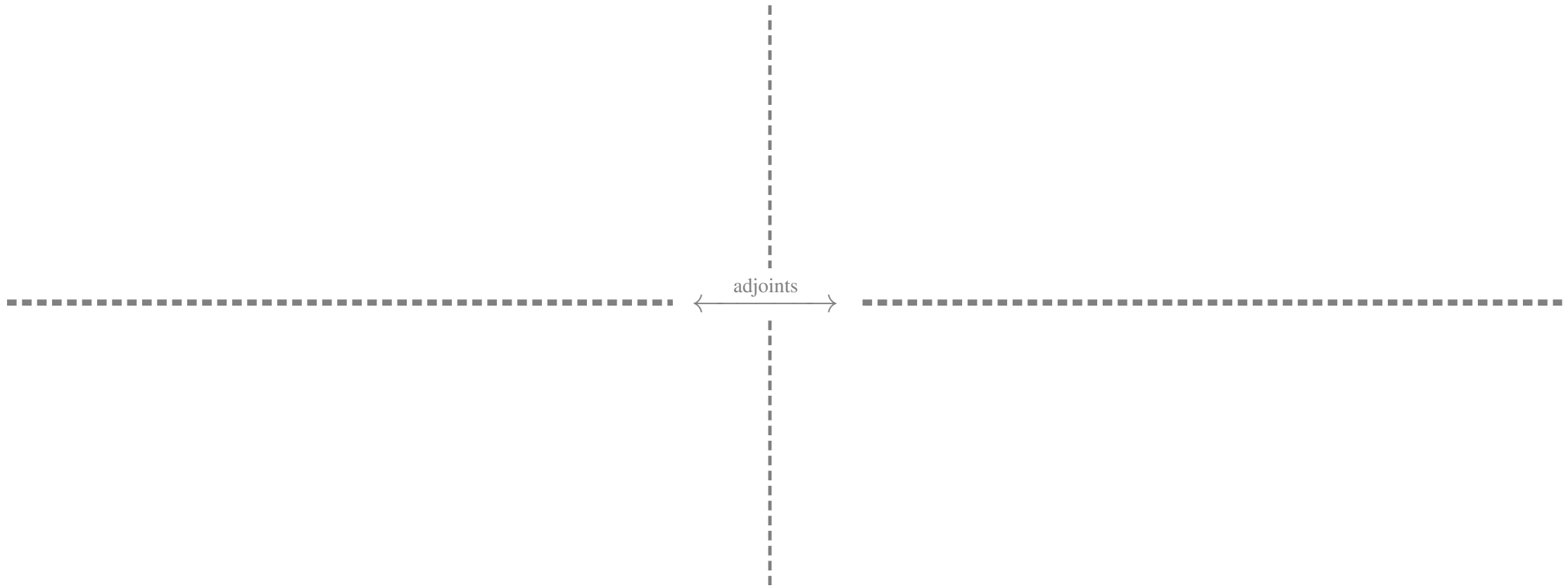
quantum data types



The pure effects of these modalities of dependent linear data type formation

are remarkable in their sheer quantum information-theoretic content.

To repeat:



The pure effects of these modalities of dependent linear data type formation

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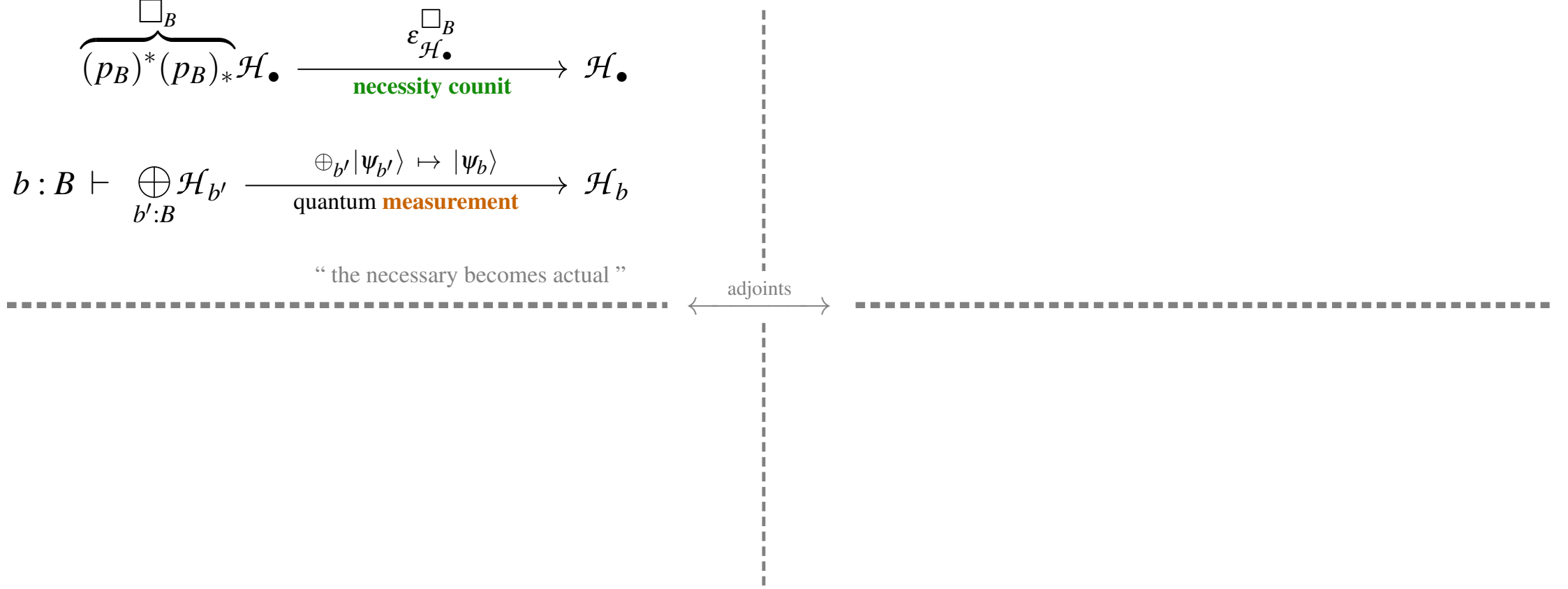
To repeat:

$$\overbrace{(p_B)^*(p_B)_* \mathcal{H}_\bullet}^{\square_B} \xrightarrow[\text{necessity counit}]{\varepsilon_{\mathcal{H}_\bullet}^{\square_B}} \mathcal{H}_\bullet$$

$$b : B \vdash \bigoplus_{b' : B} \mathcal{H}_{b'} \xrightarrow[\text{quantum measurement}]{\bigoplus_{b'} |\psi_{b'}\rangle \mapsto |\psi_b\rangle} \mathcal{H}_b$$

“ the necessary becomes actual ”

adjoints



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To repeat:

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“ the necessary becomes actual ”

adjoints

$$\mathcal{H}_\bullet \xrightarrow[\text{possibility unit } \eta_{\mathcal{H}_\bullet}^{\diamond_B}]{} \overbrace{(p_B)^*(p_B)! \mathcal{H}_\bullet}^{\diamond_B}$$

$$b : B \vdash \mathcal{H}_b \xrightarrow[\text{quantum state preparation } |\psi_b\rangle \mapsto \oplus_{b'} \begin{cases} |\psi_b\rangle & \text{if } b'=b \\ 0 & \text{else} \end{cases}]{} \bigoplus_{b':B} \mathcal{H}_{b'}$$

“ the actual is possible ”

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To repeat:

$$\overbrace{(p_B)^*(p_B)_* \mathcal{H}_\bullet}^{\square_B} \xrightarrow[\text{necessity counit } \varepsilon_{\mathcal{H}_\bullet}^{\square_B}]{} \mathcal{H}_\bullet$$

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“ the necessary becomes actual ”

“ the random becomes potential ”

$$\overbrace{(p_B)!(p_B)^* \mathcal{H}}^{\star_B} \xleftarrow[\text{randomness counit } \varepsilon_{\mathcal{H}}^{\star_B}]{} \mathcal{H}$$

$$\bigoplus_{b:B} \mathcal{H} \xrightarrow[\text{quantum superposition } \oplus_b |\psi_b\rangle \mapsto \sum_b |\psi_b\rangle]{} \mathcal{H}$$

$$\mathcal{H}_\bullet \xrightarrow[\text{possibility unit } \eta_{\mathcal{H}_\bullet}^{\diamond_B}]{} \overbrace{(p_B)^*(p_B)! \mathcal{H}_\bullet}^{\diamond_B}$$

$$b : B \vdash \mathcal{H}_b \xrightarrow[\text{quantum state preparation } |\psi_b\rangle \mapsto \oplus_{b'} \begin{cases} |\psi_b\rangle & \text{if } b'=b \\ 0 & \text{else} \end{cases}]{} \bigoplus_{b':B} \mathcal{H}_{b'}$$

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$$b : B \vdash \bigoplus_{b':B} \mathcal{H}_{b'} \xrightarrow[\text{quantum measurement } \oplus_{b'} |\psi_{b'}\rangle \mapsto |\psi_b\rangle]{} \mathcal{H}_b$$

“the necessary becomes actual”

“the random becomes potential”

$$\overbrace{(p_B)! (p_B)^* \mathcal{H}}^{\star_B} \xleftarrow[\text{randomness counit } \varepsilon_{\mathcal{H}}^{\star_B}]{} \mathcal{H}$$

$$\bigoplus_{b:B} \mathcal{H} \xrightarrow[\text{quantum superposition } \oplus_b |\psi_b\rangle \mapsto \sum_b |\psi_b\rangle]{} \mathcal{H}$$

adjoints

$$\mathcal{H}_\bullet \xrightarrow[\text{possibility unit } \eta_{\mathcal{H}_\bullet}^{\diamond_B}]{} \overbrace{(p_B)^*(p_B)! \mathcal{H}_\bullet}^{\diamond_B}$$

$$b : B \vdash \mathcal{H}_b \xrightarrow[\text{quantum state preparation } |\psi_b\rangle \mapsto \oplus_{b'} \begin{cases} |\psi_b\rangle & \text{if } b'=b \\ 0 & \text{else} \end{cases}]{} \bigoplus_{b':B} \mathcal{H}_{b'}$$

“the actual is possible”

“the potential is indeterminate”

$$\mathcal{H} \xrightarrow[\text{indeterminacy unit } \eta_{\mathcal{H}}^{\circ_B}]{} \overbrace{(p_B)_* (p_B)^* \mathcal{H}}^{\circ_B}$$

$$\mathcal{H} \xrightarrow[\text{quantum parallelism } |\psi\rangle \mapsto \oplus_b |\psi\rangle_b]{} \bigoplus_{b:B} \mathcal{H}$$

Q-bits are the free linear indeterminacy-effect handlers over $\text{Bit} = \{0, 1\}$

Coherent q-bits:

$$\begin{array}{c}
 \text{——} \quad \text{QBit} : \text{LType} \xrightarrow{\mathbb{1}_{\text{Bit}} \otimes} \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ B} \\
 \parallel \\
 \circ_{\text{Bit}} \mathbb{1}
 \end{array}$$

Quantum gate with q-bit output:

De-cohered (measured) q-bits:

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 \parallel \\
 \circ_{\text{Bit}} \mathbb{1} = \oplus_{\{0,1\}} \mathbb{C} = \mathbb{C} \cdot |0\rangle \oplus \mathbb{C} \cdot |1\rangle
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———— QBit

\otimes

———— \mathcal{H}

\parallel

$$\bigcirc_{\text{Bit}} \mathcal{H} = \bigoplus_{\{0,1\}} \mathcal{H} = \mathcal{H} \otimes |0\rangle \oplus \mathcal{H} \otimes |1\rangle$$

Quantum gate with q-bit output:

De-cohered (measured) q-bits:

Q-bits are the free linear indeterminacy-effect handlers over $\text{Bit} = \{0, 1\}$

Coherent q-bits:

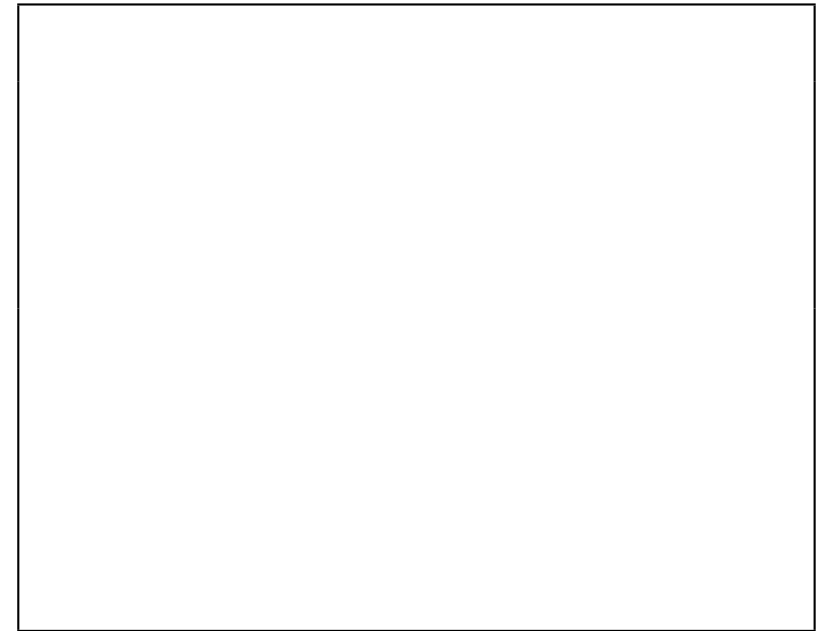
$$\begin{array}{c} \text{————} \quad \text{QBit} : \text{LType} \xrightarrow{\mathbb{1}_{\text{Bit}} \otimes} \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ_B} \\ \parallel \\ \bigcirc_{\text{Bit}} \mathbb{1} = \bigoplus_{\{0,1\}} \mathbb{C} = \mathbb{C} \cdot |0\rangle \oplus \mathbb{C} \cdot |1\rangle \end{array}$$

$$\begin{array}{c} \text{————} \quad \text{QBit} \\ \otimes \\ \text{————} \quad \mathcal{H} \\ \parallel \\ \bigcirc_{\text{Bit}} \mathcal{H} = \bigoplus_{\{0,1\}} \mathcal{H} = \mathcal{H} \otimes |0\rangle \oplus \mathcal{H} \otimes |1\rangle \end{array}$$

De-cohered (measured) q-bits:

$$\begin{array}{c} \text{=====} \quad \mathbb{1}_{\text{Bit}} : \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ_{\text{Bit}}} \\ b : \text{Bit} \quad \vdash \quad \mathbb{C} \cdot |b\rangle : \text{LType} \end{array}$$

Quantum gate with q-bit output:



Q-bits are the free linear indeterminacy-effect handlers over $\text{Bit} = \{0, 1\}$

Coherent q-bits:

$$\begin{array}{c} \text{————} \quad \text{QBit} : \text{LType} \xrightarrow{\mathbb{1}_{\text{Bit}} \otimes} \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ B} \\ \parallel \\ \bigcirc_{\text{Bit}} \mathbb{1} = \bigoplus_{\{0,1\}} \mathbb{C} = \mathbb{C} \cdot |0\rangle \oplus \mathbb{C} \cdot |1\rangle \end{array}$$

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Quantum gate with q-bit output:

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$$\begin{array}{c} \text{====} \quad \mathbb{1}_{\text{Bit}} \\ \otimes \quad b : \text{Bit} \quad \vdash \quad \mathcal{H} \otimes |b\rangle : \text{LType} \\ \text{————} \quad \mathcal{H} \end{array}$$

Q-bits are the free linear indeterminacy-effect handlers over $\text{Bit} = \{0, 1\}$

Coherent q-bits:

$$\begin{array}{c} \text{—————} \\ \text{QBit} : \text{LType} \end{array} \xrightarrow{\mathbb{1}_{\text{Bit}} \otimes} \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ_{\text{B}}}$$

$$\begin{array}{c} \parallel \\ \text{O}_{\text{Bit}} \mathbb{1} = \oplus_{\{0,1\}} \mathbb{C} = \mathbb{C} \cdot |0\rangle \oplus \mathbb{C} \cdot |1\rangle \end{array}$$

$$\begin{array}{c} \text{—————} \\ \text{QBit} \\ \otimes \\ \text{—————} \\ \mathcal{H} \\ \parallel \\ \text{O}_{\text{Bit}} \mathcal{H} = \oplus_{\{0,1\}} \mathcal{H} = \mathcal{H} \otimes |0\rangle \oplus \mathcal{H} \otimes |1\rangle \end{array}$$

De-cohered (measured) q-bits:

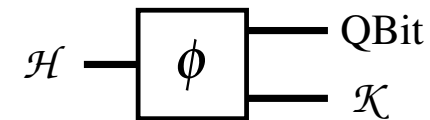
$$\begin{array}{c} \text{=====} \\ \mathbb{1}_{\text{Bit}} : \text{LType}_{\text{Bit}} \end{array} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ_{\text{Bit}}}$$

$$b : \text{Bit} \quad \vdash \quad \mathbb{C} \cdot |b\rangle : \text{LType}$$

$$\begin{array}{c} \text{=====} \\ \mathbb{1}_{\text{Bit}} \\ \otimes \\ \text{—————} \\ \mathcal{H} \end{array} \quad b : \text{Bit} \quad \vdash \quad \mathcal{H} \otimes |b\rangle : \text{LType}$$

Quantum gate with q-bit output:

A quantum gate which may handle O_{Bit} -effects is one with a QBit-output:



$$\mathcal{H} \xrightarrow{\phi} \text{QBit} \otimes \mathcal{K} \simeq \text{O}_{\text{Bit}} \mathcal{K}$$

Q-bits are the free linear indeterminacy-effect handlers over $\text{Bit} = \{0, 1\}$

Coherent q-bits:

$$\begin{array}{c} \text{——} \quad \text{QBit} : \text{LType} \xrightarrow{\mathbb{1}_{\text{Bit}} \otimes} \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ_{\text{B}}} \\ \parallel \\ \circ_{\text{Bit}} \mathbb{1} = \oplus_{\{0,1\}} \mathbb{C} = \mathbb{C} \cdot |0\rangle \oplus \mathbb{C} \cdot |1\rangle \end{array}$$

$$\begin{array}{c} \text{——} \quad \text{QBit} \\ \otimes \\ \text{——} \quad \mathcal{H} \\ \parallel \\ \circ_{\text{Bit}} \mathcal{H} = \oplus_{\{0,1\}} \mathcal{H} = \mathcal{H} \otimes |0\rangle \oplus \mathcal{H} \otimes |1\rangle \end{array}$$

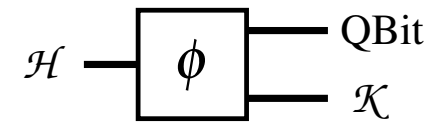
De-cohered (measured) q-bits:

$$\begin{array}{c} \text{====} \quad \mathbb{1}_{\text{Bit}} : \text{LType}_{\text{Bit}} \xrightarrow[\sim]{\oplus_{\text{Bit}}} \text{LType}^{\circ_{\text{Bit}}} \\ b : \text{Bit} \quad \vdash \quad \mathbb{C} \cdot |b\rangle : \text{LType} \end{array}$$

$$\begin{array}{c} \text{====} \quad \mathbb{1}_{\text{Bit}} \\ \otimes \quad b : \text{Bit} \quad \vdash \quad \mathcal{H} \otimes |b\rangle : \text{LType} \\ \text{——} \quad \mathcal{H} \end{array}$$

Quantum gate with q-bit output:

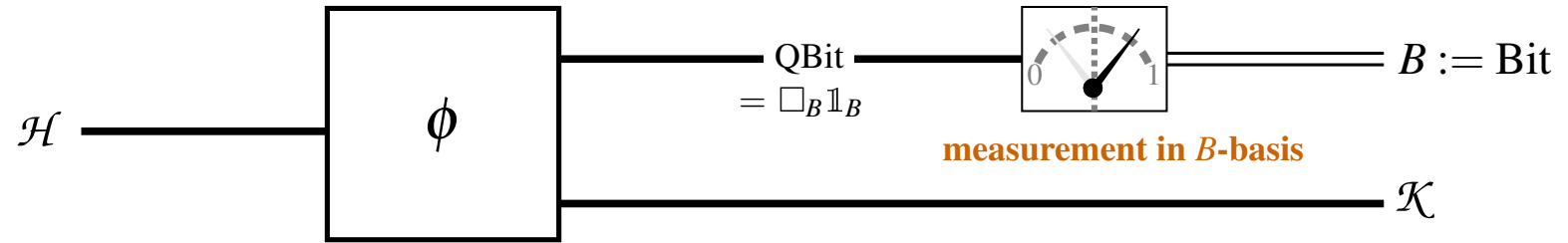
A quantum gate which may handle \circ_{Bit} -effects is one with a QBit-output:



$$\mathcal{H} \xrightarrow{\phi} \text{QBit} \otimes \mathcal{K} \simeq \circ_{\text{Bit}} \mathcal{K}$$

Quantum measurement is Linear indefiniteness-effect handling.

quantum circuit



quantum gate

$$\mathcal{H} \xrightarrow{\phi} \text{QBit} \otimes \mathcal{K} \simeq \circ_B \mathcal{K}$$

\circ_B -effect handling

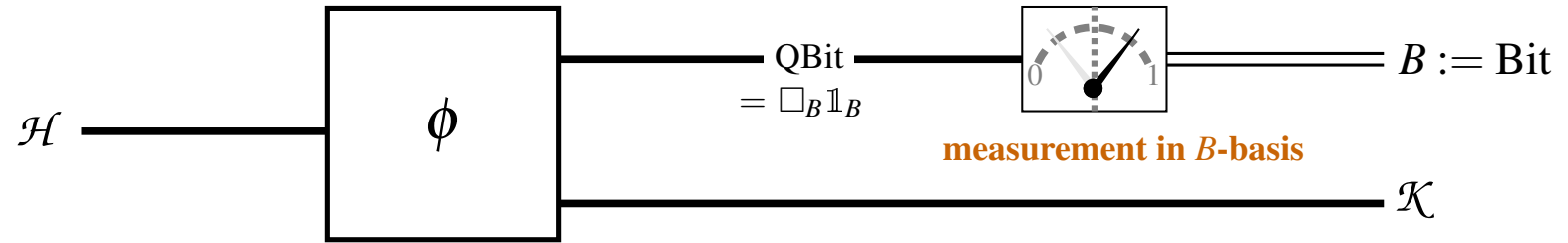
Quantum measurement is Linear indefiniteness-effect handling.

quantum circuit

formalization
↓

\circ_B -modal linear types

LType_{\circ_B}



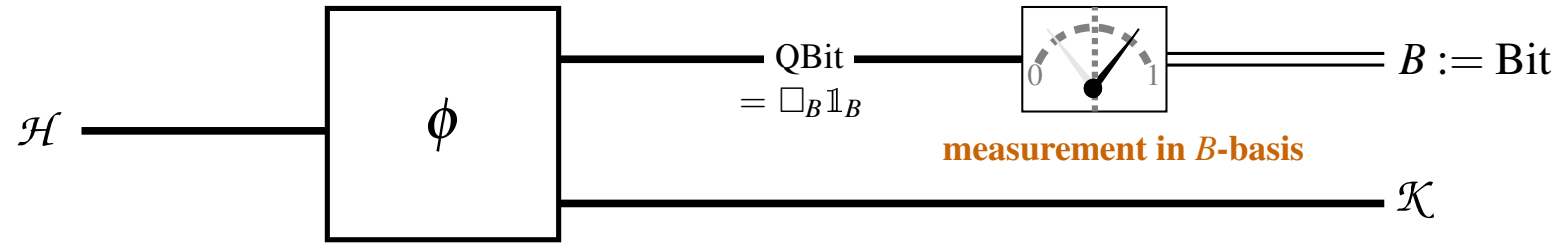
quantum gate

$$\mathcal{H} \xrightarrow{\phi} \text{QBit} \otimes \mathcal{K} \simeq \circ_B \mathcal{K}$$

\circ_B -effect handling

Quantum measurement is Linear indefiniteness-effect handling.

quantum circuit



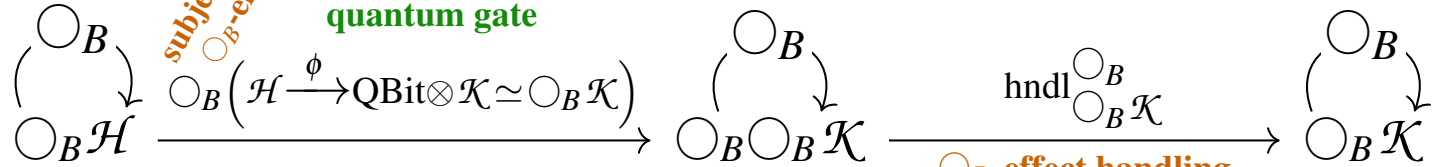
formalization

subjection to \circ_B -effects

quantum gate

measurement in B -basis

\circ_B -effect handling



\circ_B -modal linear types

LType_{\circ_B}

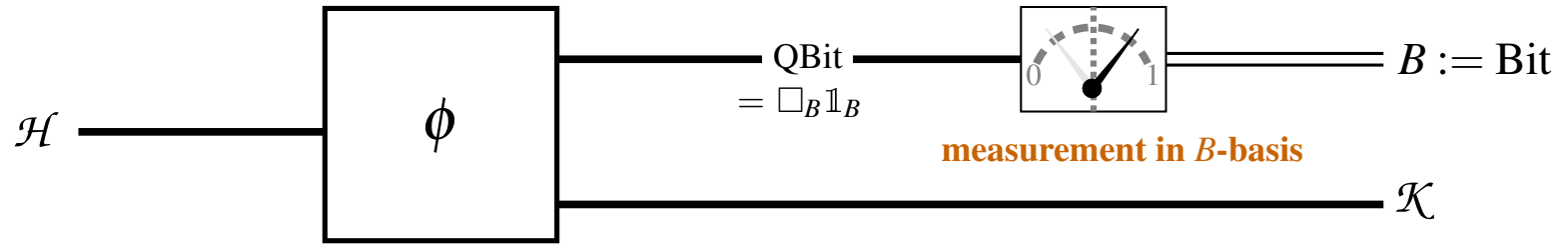
comparison
functor
 $K_{(p_B)^*(p_B)^*}$

LType_B

B -dependent linear types

Quantum measurement is Linear indefiniteness-effect handling.

quantum circuit



formalization

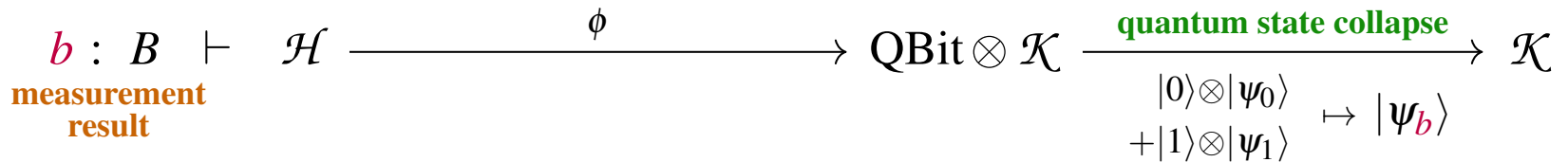
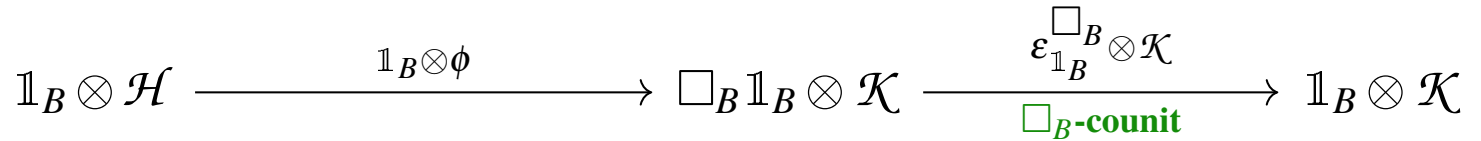
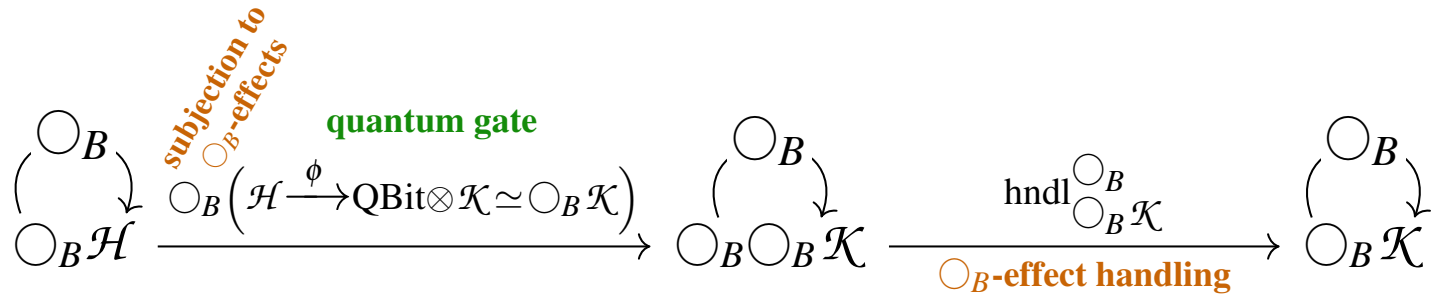
\circ_B -modal linear types

LType \circ_B

comparison
functor
 $K^{(p_B)^* (p_B)^*}$

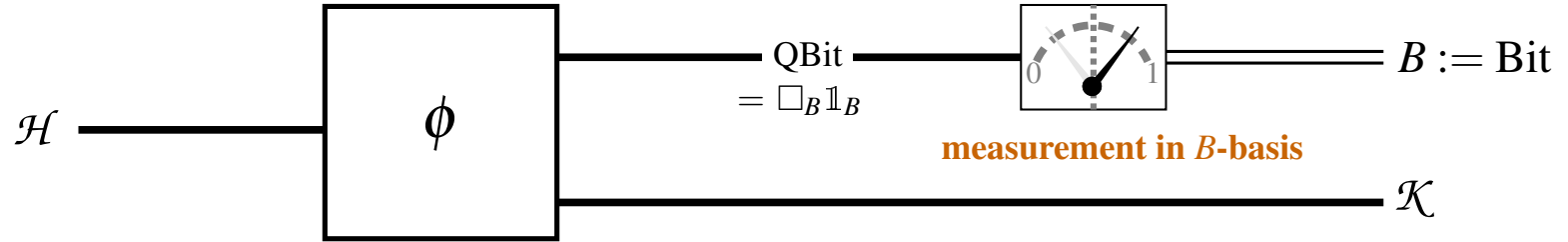
LType $_B$

B -dependent linear types



Quantum measurement is Linear indefiniteness-effect handling.

quantum circuit



formalization

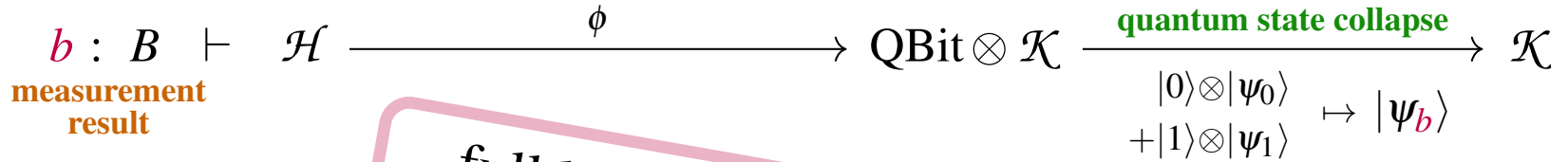
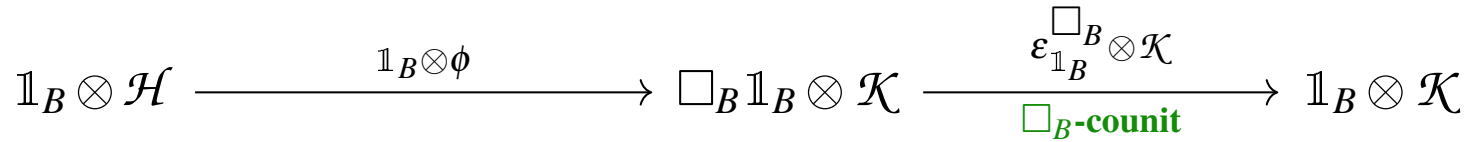
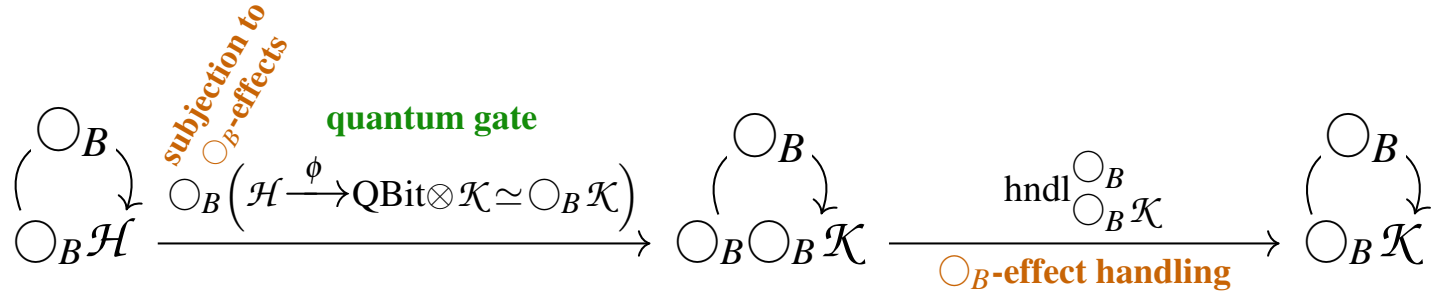
\circ_B -modal linear types

LType \circ_B

comparison
functor
 $K^{(p_B)^*} \rightarrow (p_B)^*$

LType $_B$

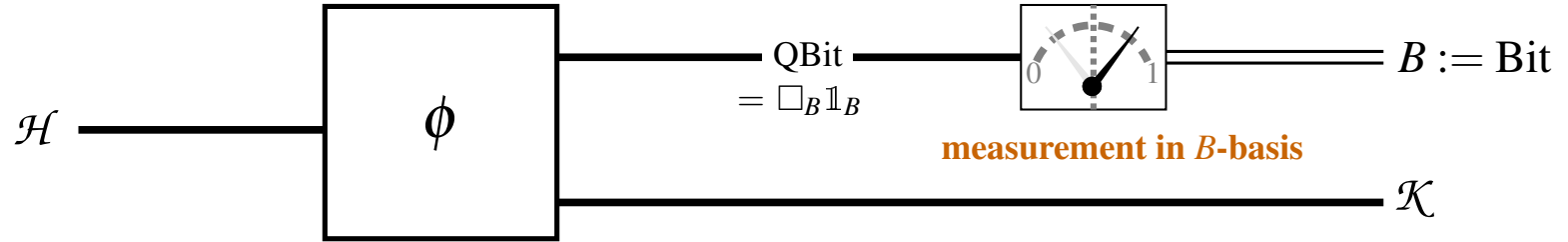
B -dependent linear types



full linearly-typed detail of quantum measurement logic is emergent effect in dLHoTT

Quantum measurement is Linear indefiniteness-effect handling.

quantum circuit



formalization

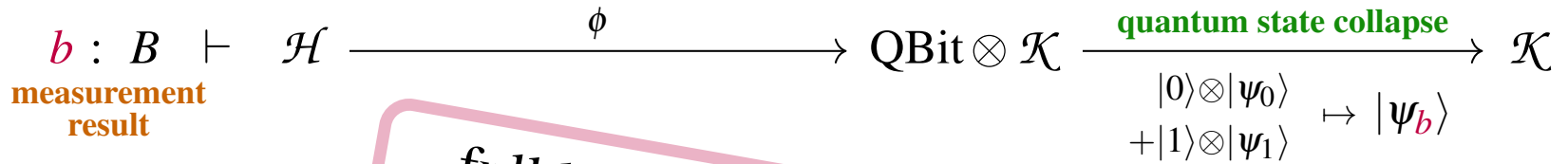
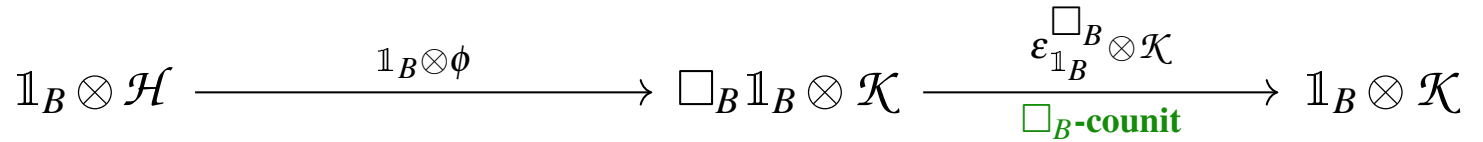
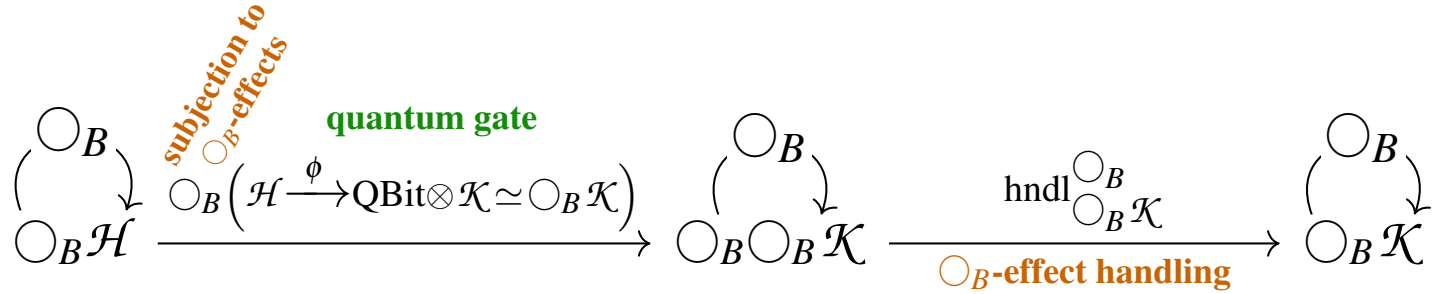
\circ_B -modal linear types

LType_{\circ_B}

comparison functor
 $K_{(p_B)^*(p_B)^*}$

LType_B

B -dependent linear types



full linearly-typed detail of quantum measurement logic is emergent effect in dLHoTT

Aside: **Linear indefiniteness monad recovers Coecke’s “classical structures”.**

(see [nLab:quantum+reader+monad](#))

\circ_B

\pitchfork

Monad(LType)

Aside: Linear indefiniteness monad recovers Coecke's “classical structures”.

(see [nLab:quantum+reader+monad](#))

\circ_B

!!

B-Reader

\mho

Monad(LType)

Aside: Linear indefiniteness monad recovers Coecke’s “classical structures”.

(see [nLab:quantum+reader+monad](#))

\circlearrowleft_B
 \Downarrow
 B -Reader
 \Downarrow
 $\mathbb{1}^B$ -Writer

$$\mathbb{1}^B\text{-Writer}(D) := \mathbb{1}^B \otimes D$$

$$\text{bind}_{\mathbb{1}^B\text{Writer}}(D_1 \xrightarrow{\text{prog}} \mathbb{1}^B \otimes D_2) := \mathbb{1}^B \otimes D_1 \xrightarrow{\mathbb{1}^B \otimes \text{prog}} \mathbb{1}^B \otimes (\mathbb{1}^B \otimes D_2) \xrightarrow{\mu \otimes \text{id}_{D_2}} \mathbb{1}^B \otimes D_2$$

$B : \text{FinType} \vdash$

$\text{Monad}(\text{LType})$

Where $\mathbb{1}^B = \bigoplus_{b:B} \mathbb{C} \cdot P_b \in \text{CMon}(\text{LType})$ is

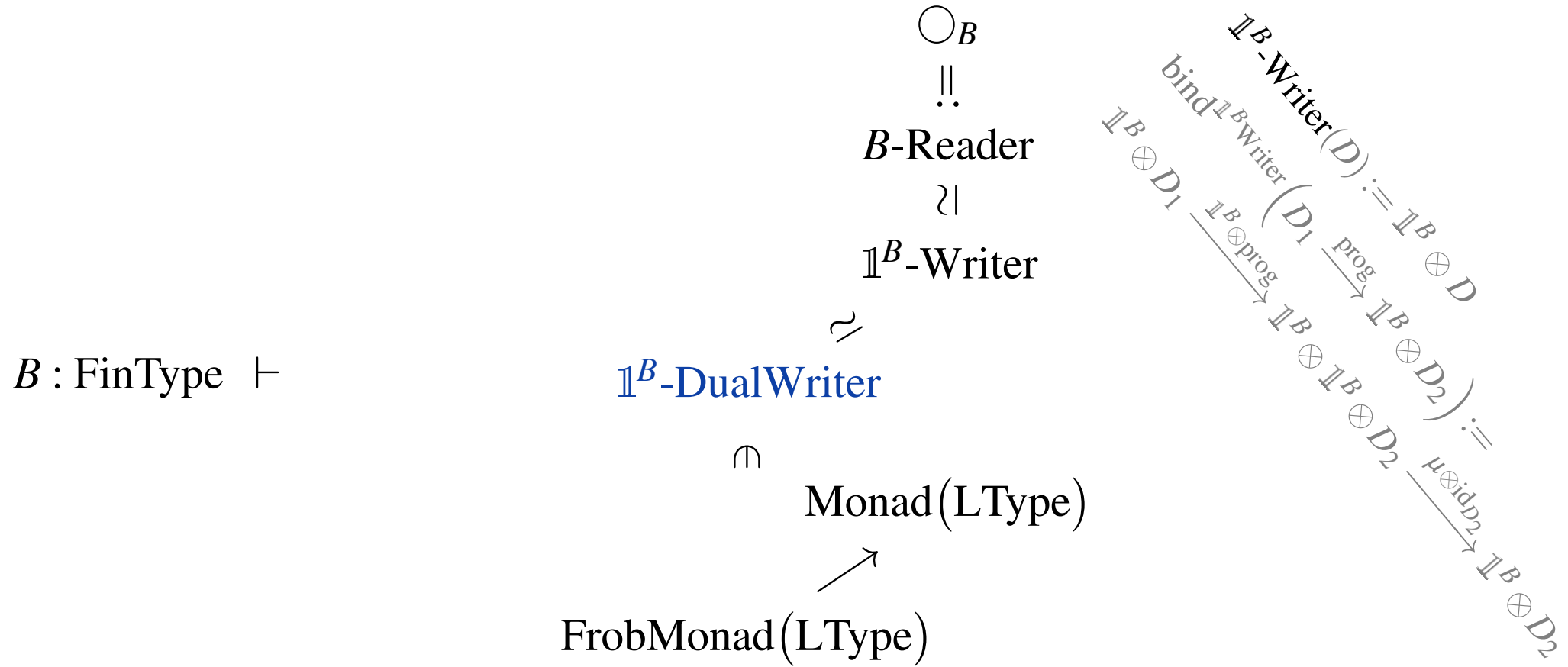
algebra of B -projection operators :

$$\mathbb{1} \xrightarrow[\substack{1 \mapsto \sum_{b:B} P_b}]{\text{unit } \eta} \mathbb{1}^B$$

$$\mathbb{1}^B \otimes \mathbb{1}^B \xrightarrow[\substack{P_b \otimes P_{b'} \mapsto \begin{cases} P_b & \text{if } b=b' \\ 0 & \text{else} \end{cases}}]{\text{product } \mu} \mathbb{1}^B$$

Aside: Linear indefiniteness monad recovers Coecke’s “classical structures”.

(see [nLab:quantum+reader+monad](#))

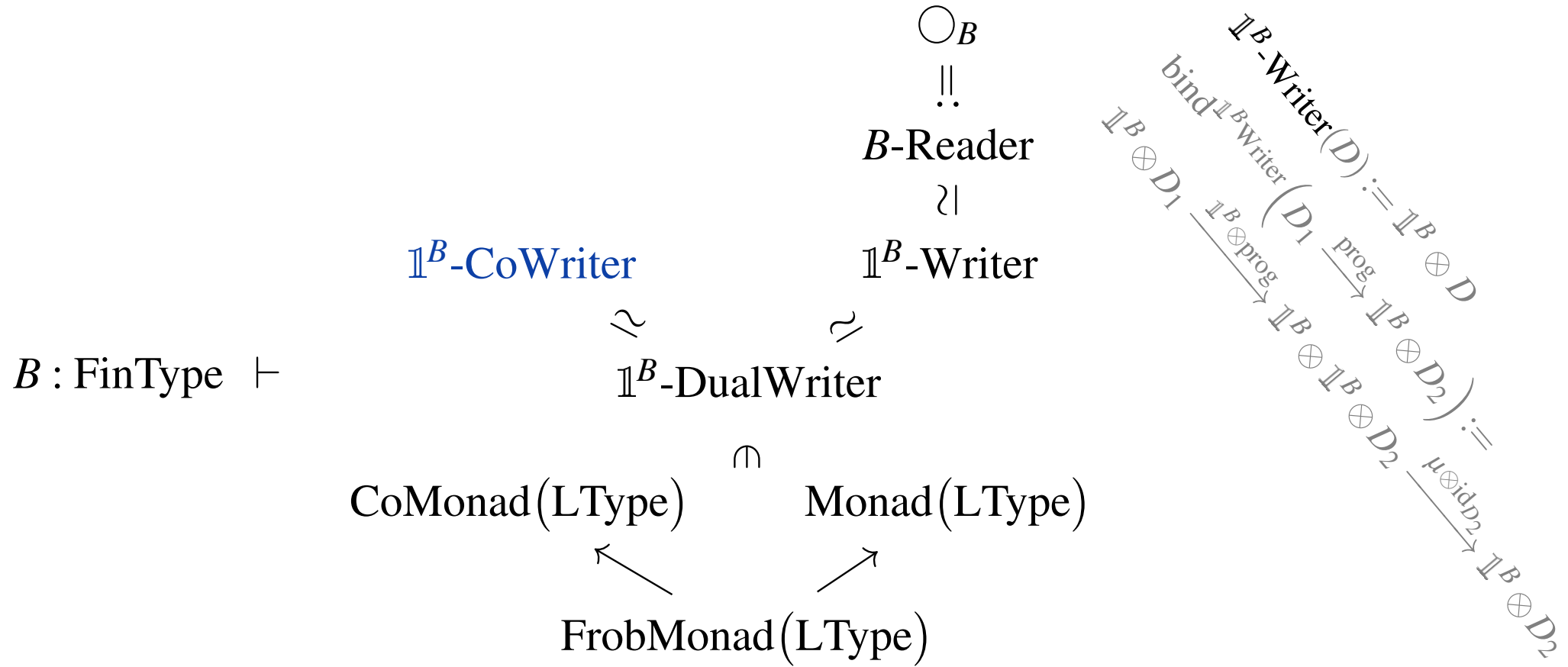


Where $\mathbb{1}^B = \bigoplus_{b:B} \mathbb{C} \cdot P_b \in \text{CMon}(\text{LType})$ is **Frobenius** algebra of B -projection operators :

$$\mathbb{1} \xrightarrow[\substack{\eta \\ 1 \mapsto \sum_{b:B} P_b}]{\text{unit}} \mathbb{1}^B \xrightarrow[\substack{\delta \\ P_b \mapsto P_b \otimes P_b}]{\text{co-product}} \mathbb{1}^B \otimes \mathbb{1}^B \xrightarrow[\substack{\mu \\ P_b \otimes P_{b'} \mapsto \begin{cases} P_b & \text{if } b=b' \\ 0 & \text{else} \end{cases}}]{\text{product}} \mathbb{1}^B \xrightarrow[\substack{\varepsilon \\ P_b \mapsto 1}]{\text{co-unit}} \mathbb{1}$$

Aside: Linear indefiniteness monad recovers Coecke’s “classical structures”.

(see [nLab:quantum+reader+monad](#))

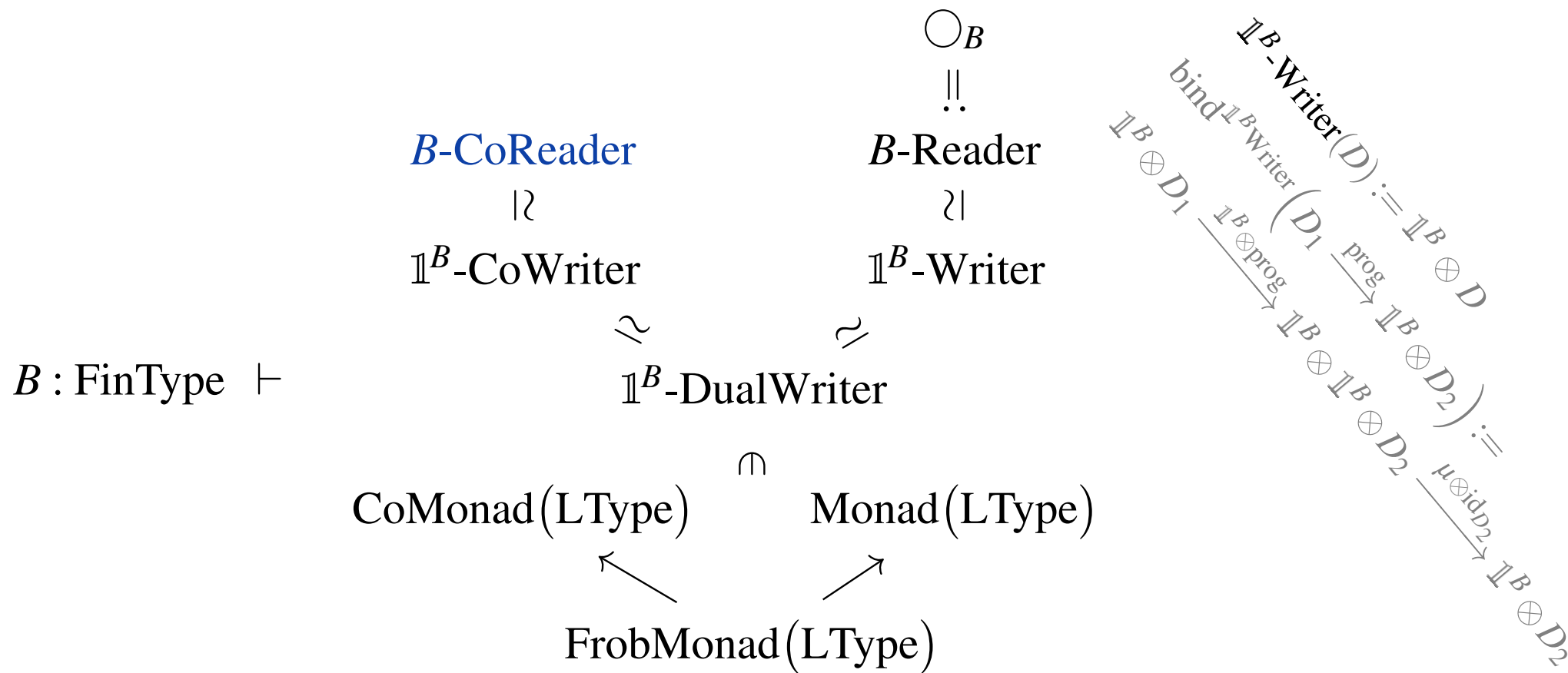


Where $\mathbb{1}^B = \bigoplus_{b:B} \mathbb{C} \cdot P_b \in \text{CMon}(\text{LType})$ is Frobenius algebra of B -projection operators :

$$\begin{array}{ccccccc}
 \mathbb{1} & \xrightarrow[\substack{1 \mapsto \sum_{b:B} P_b}]{\text{unit } \eta} & \mathbb{1}^B & \xrightarrow[\substack{P_b \mapsto P_b \otimes P_b}]{\text{co-product } \delta} & \mathbb{1}^B \otimes \mathbb{1}^B & \xrightarrow[\substack{P_b \otimes P_{b'} \mapsto \begin{cases} P_b & \text{if } b=b' \\ 0 & \text{else} \end{cases}}]{\text{product } \mu} & \mathbb{1}^B & \xrightarrow[\substack{P_b \mapsto 1}]{\text{co-unit } \varepsilon} & \mathbb{1}
 \end{array}$$

Aside: Linear indefiniteness monad recovers Coecke's “classical structures”.

(see [nLab:quantum+reader+monad](#))

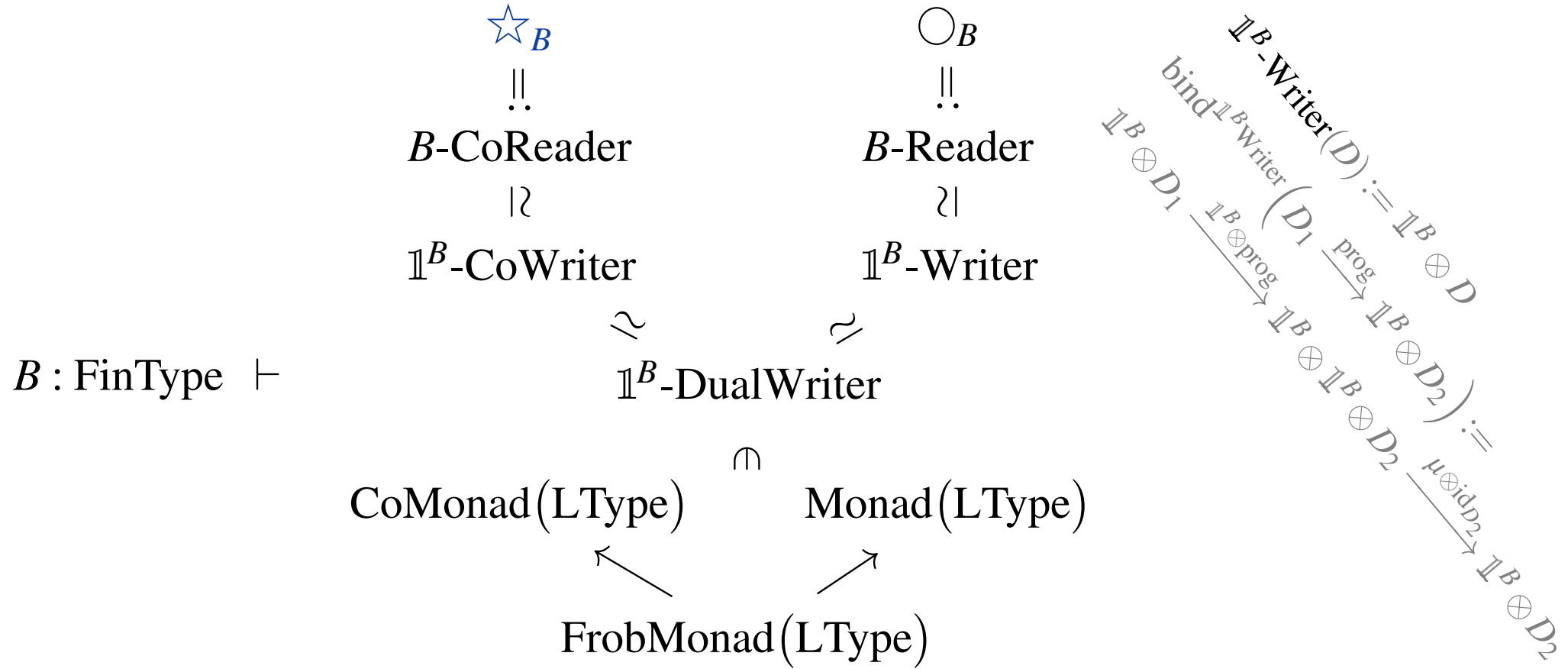


Where $\mathbb{1}^B = \bigoplus_{b:B} \mathbb{C} \cdot P_b \in \text{CMon}(\text{LType})$ is Frobenius algebra of B -projection operators :

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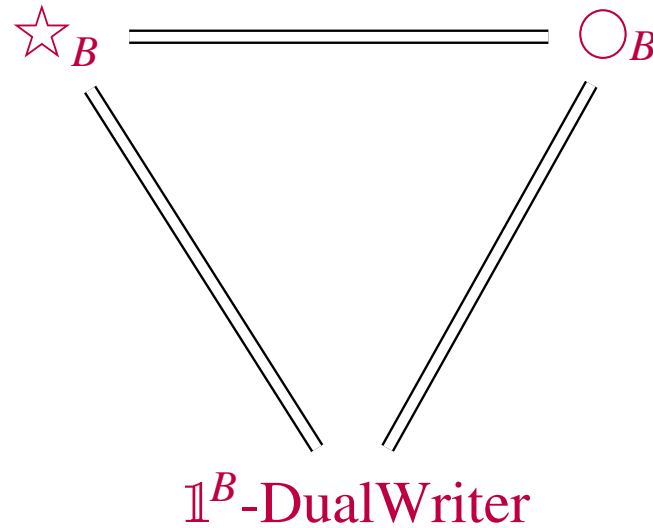


Where $\mathbb{1}^B = \bigoplus_{b:B} \mathbb{C} \cdot P_b \in \text{CMon}(\text{LType})$ is Frobenius algebra of B -projection operators :

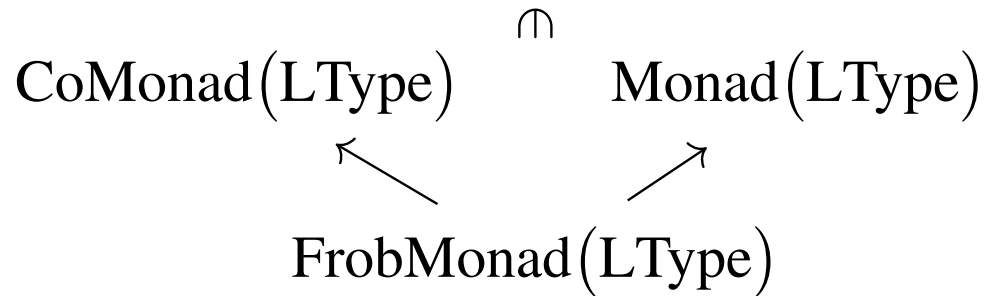
$$\mathbb{1} \xrightarrow[\substack{\eta \\ 1 \mapsto \sum_{b:B} P_b}]{\text{unit}} \mathbb{1}^B \xrightarrow[\substack{\delta \\ P_b \mapsto P_b \otimes P_b}]{\text{co-product}} \mathbb{1}^B \otimes \mathbb{1}^B \xrightarrow[\substack{\mu \\ P_b \otimes P_{b'} \mapsto \begin{cases} P_b & \text{if } b=b' \\ 0 & \text{else} \end{cases}}]{\text{product}} \mathbb{1}^B \xrightarrow[\substack{\varepsilon \\ P_b \mapsto 1}]{\text{co-unit}} \mathbb{1}$$

Aside: Linear indefiniteness monad recovers Coecke’s “classical structures”.

(see [nLab:quantum+reader+monad](#))



$B : \text{FinType} \vdash$



Where $\mathbb{1}^B = \bigoplus_{b:B} \mathbb{C} \cdot P_b \in \text{CMon}(\text{LType})$ is Frobenius algebra of B -projection operators :

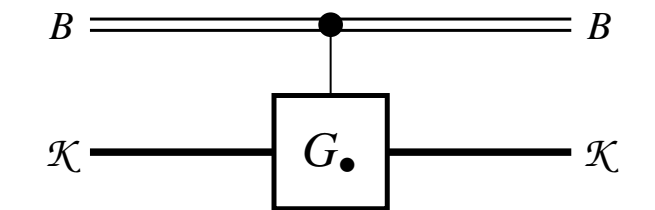
$$\begin{array}{ccccccc}
 \mathbb{1} & \xrightarrow[\substack{\eta \\ 1 \mapsto \sum_{b:B} P_b}]{\text{unit}} & \mathbb{1}^B & \xrightarrow[\substack{\delta \\ P_b \mapsto P_b \otimes P_b}]{\text{co-product}} & \mathbb{1}^B \otimes \mathbb{1}^B & \xrightarrow[\substack{\mu \\ P_b \otimes P_{b'} \mapsto \begin{cases} P_b & \text{if } b=b' \\ 0 & \text{else} \end{cases}}]{\text{product}} & \mathbb{1}^B & \xrightarrow[\substack{\varepsilon \\ P_b \mapsto 1}]{\text{co-unit}} & \mathbb{1}
 \end{array}$$

Exmp: Deferred measurement principle – Proven by monadic effect logic.



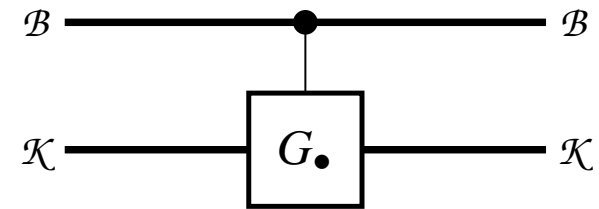
classically controlled gate

quantumly controlled gate



$$\mathcal{B} \boxtimes \mathcal{K} \xrightarrow{G} \mathcal{B} \boxtimes \mathcal{K}$$

$$b : B \vdash \mathcal{K} \xrightarrow{G_b} \mathcal{K}$$

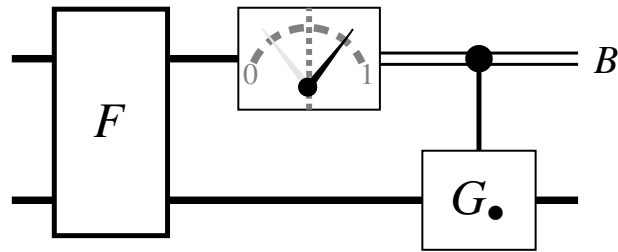


$$\square_B \mathcal{B} \boxtimes \mathcal{K} \xrightarrow{\square_B G} \square_B \mathcal{B} \boxtimes \mathcal{K}$$

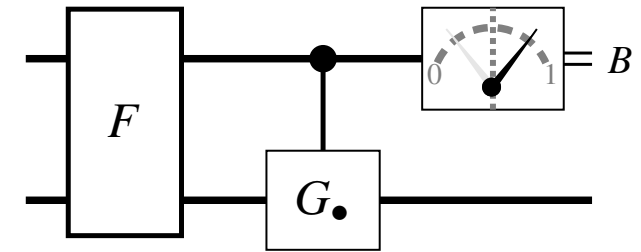
$$b : B \vdash \bigoplus_{b' : B} \mathcal{K} \xrightarrow{\bigoplus_{b' : B} G_{b'}} \bigoplus_{b' : B} \mathcal{K}$$

Exmp: Deferred measurement principle – Proven by monadic effect logic.

$$\begin{array}{ccccc}
 \square_B \mathcal{H}_\bullet \xrightarrow{F} \square_B \mathcal{H}_\bullet \xrightarrow{\varepsilon_{\mathcal{H}_\bullet}^{\square_B}} \mathcal{H}_\bullet \xrightarrow{G_\bullet} \mathcal{H}_\bullet & \mapsto & \square_B \mathcal{H}_\bullet \xrightarrow{\text{diag}_B(G_\bullet) \circ F} \square_B \mathcal{H}_\bullet & \mapsto & \square_B \mathcal{H}_\bullet \xrightarrow{\text{diag}_B(G_\bullet) \circ F} \square_B \mathcal{H}_\bullet \xrightarrow{\varepsilon_{\mathcal{H}_\bullet}^{\square_B}} \mathcal{H}_\bullet \\
 \text{measurement-controlled quantum gate} & & \text{quantum-controlled quantum gate...} & & \text{...followed by measurement}
 \end{array}$$

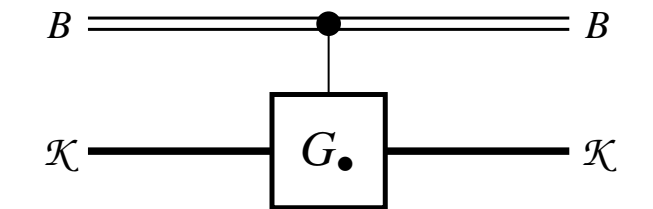


← Deferred Measurement Principle →



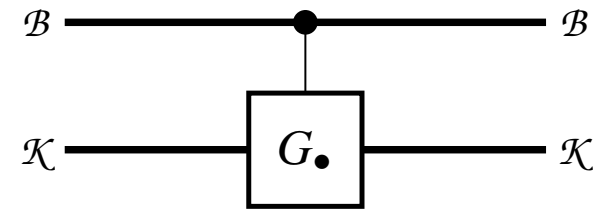
classically controlled gate

quantumly controlled gate



$$B_\bullet \boxtimes K \xrightarrow{G_\bullet} B_\bullet \boxtimes K$$

$$b : B \vdash K \xrightarrow{G_b} K$$



$$\square_B B_\bullet \boxtimes K \xrightarrow{\square_B G_\bullet} \square_B B_\bullet \boxtimes K$$

$$b : B \vdash \bigoplus_{b' : B} K \xrightarrow{\bigoplus_{b' : B} G_{b'}} \bigoplus_{b' : B} K$$

Exmp: Deferred measurement principle – Proven by monadic effect logic.

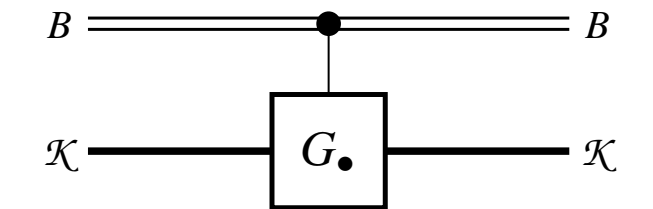
$$\begin{array}{c}
 \text{id} \\
 \downarrow \\
 \text{Kl}(\square_B) \xrightarrow[\delta^B \circ \square_B(-)]{\sim} \text{LType}_{B\square_B} \xrightarrow[\varepsilon^{\square_B \circ (-)}]{\sim} \text{Kl}(\square_B) \\
 \text{\scriptsize } \square_B\text{-Kleisli morphisms} \qquad \qquad \qquad \text{\scriptsize } \square_B\text{-coalgebra homomorphisms} \qquad \qquad \qquad \text{\scriptsize } \square_B\text{-Kleisli morphisms} \\
 \text{Kleisli equivalence}
 \end{array}$$

$$\begin{array}{c}
 \square_B\mathcal{H}_\bullet \xrightarrow{F} \square_B\mathcal{H}_\bullet \xrightarrow{\varepsilon^{\square_B}_{\mathcal{H}_\bullet}} \mathcal{H}_\bullet \xrightarrow{G_\bullet} \mathcal{H}_\bullet \quad \mapsto \quad \square_B\mathcal{H}_\bullet \xrightarrow{\text{diag}_B(G_\bullet) \circ F} \square_B\mathcal{H}_\bullet \quad \mapsto \quad \square_B\mathcal{H}_\bullet \xrightarrow{\text{diag}_B(G_\bullet) \circ F} \square_B\mathcal{H}_\bullet \xrightarrow{\varepsilon^{\square_B}_{\mathcal{H}_\bullet}} \mathcal{H}_\bullet \\
 \text{measurement-controlled quantum gate} \qquad \qquad \qquad \text{quantum-controlled quantum gate...} \qquad \qquad \qquad \text{...followed by measurement}
 \end{array}$$



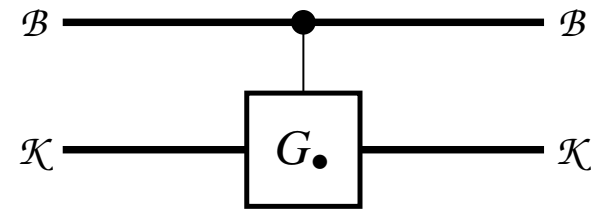
classically controlled gate

quantumly controlled gate



$$\mathcal{B}_\bullet \boxtimes \mathcal{K} \xrightarrow{G_\bullet} \mathcal{B}_\bullet \boxtimes \mathcal{K}$$

$$b : B \vdash \mathcal{K} \xrightarrow{G_b} \mathcal{K}$$

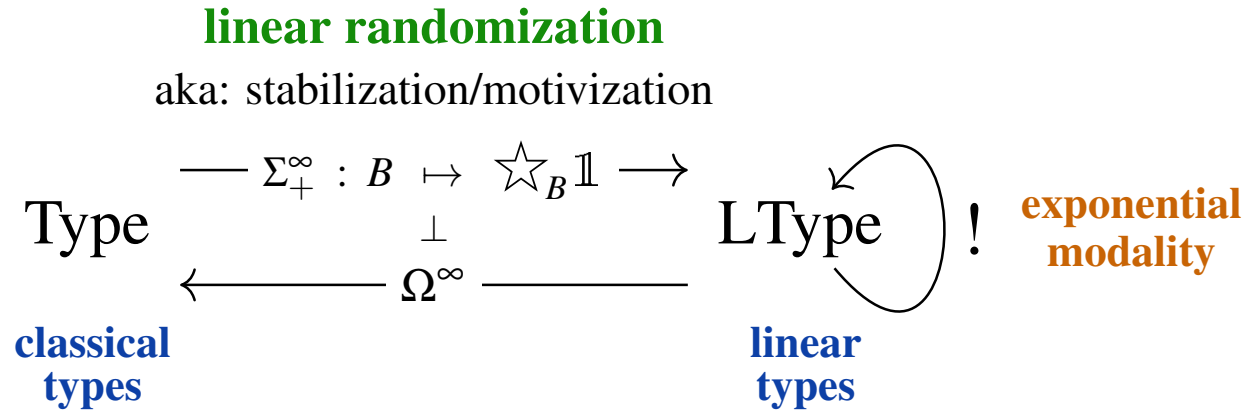


$$\square_B \mathcal{B}_\bullet \boxtimes \mathcal{K} \xrightarrow{\square_B G_\bullet} \square_B \mathcal{B}_\bullet \boxtimes \mathcal{K}$$

$$b : B \vdash \bigoplus_{b' : B} \mathcal{K} \xrightarrow{\bigoplus_{b' : B} G_{b'}} \bigoplus_{b' : B} \mathcal{K}$$

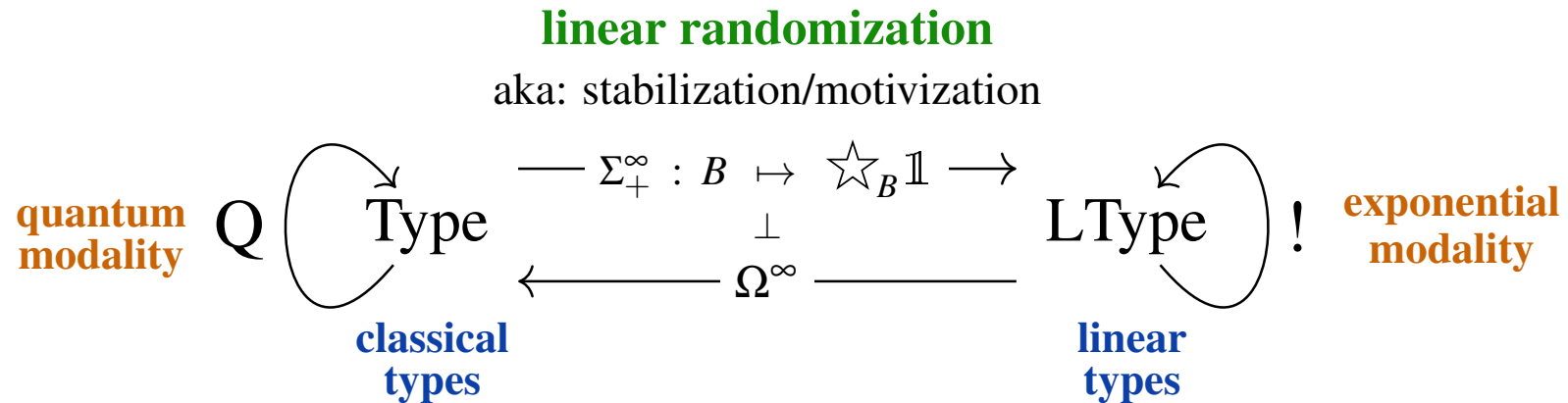
The Quantum modality.

Also the *exponential modality* traditionally postulated in linear logic is an emergent effect in dLHoTT,



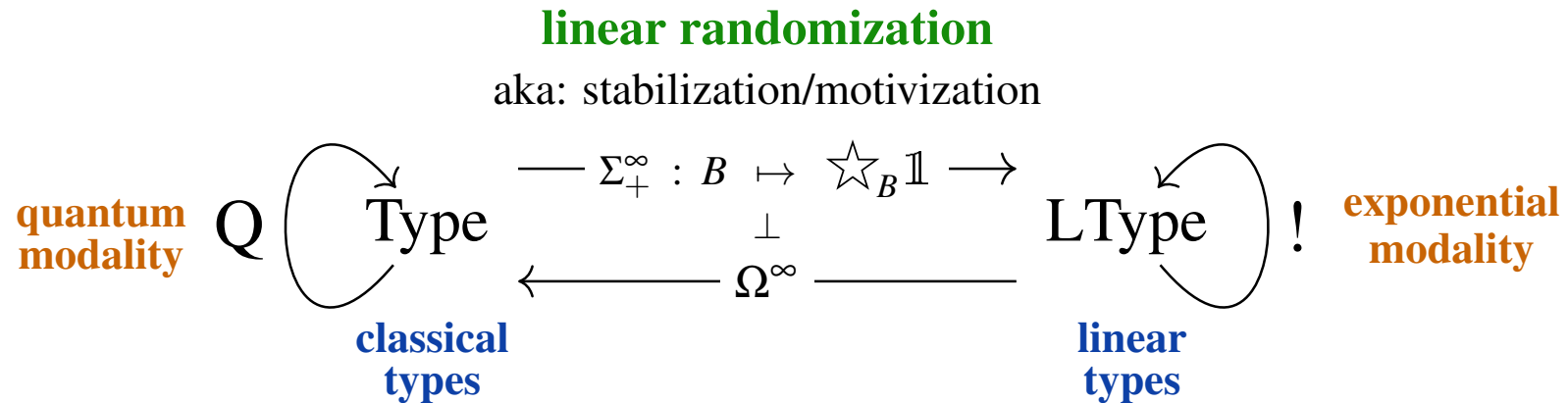
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The Quantum modality.

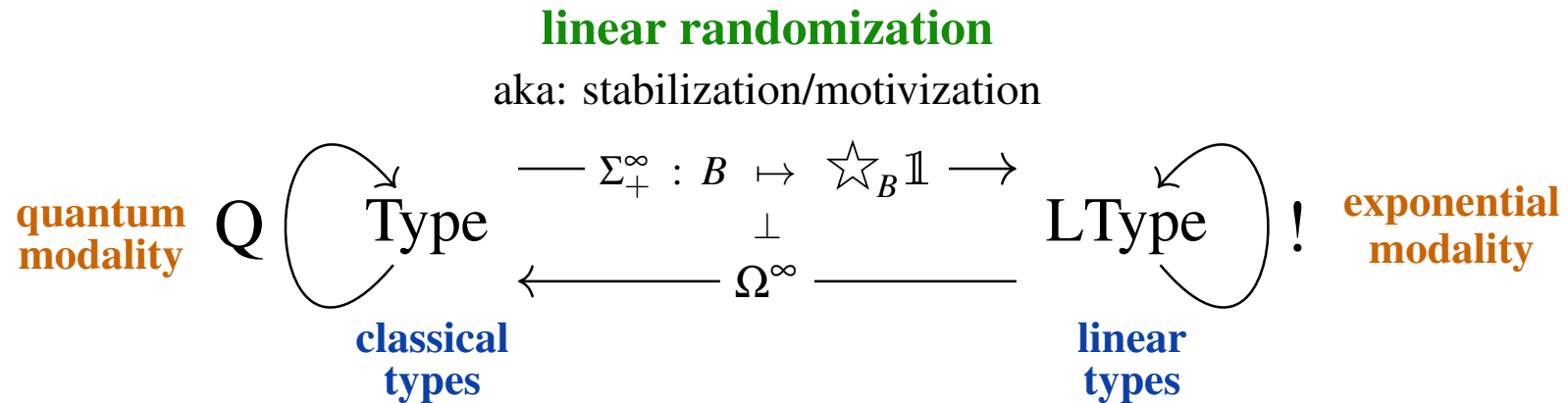
Also the *exponential modality* traditionally postulated in linear logic is an emergent effect in dLHoTT, as is the crucial *Quantum Modality*, not considered before:



The Q-monad plays a crucial role in the full formulation of the QS-language. It is the secret actor behind $\text{QBit} = Q(\text{Bit})\dots$

The Quantum modality.

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Quantum Circuits

Quantum effects are compatible with tensor product.

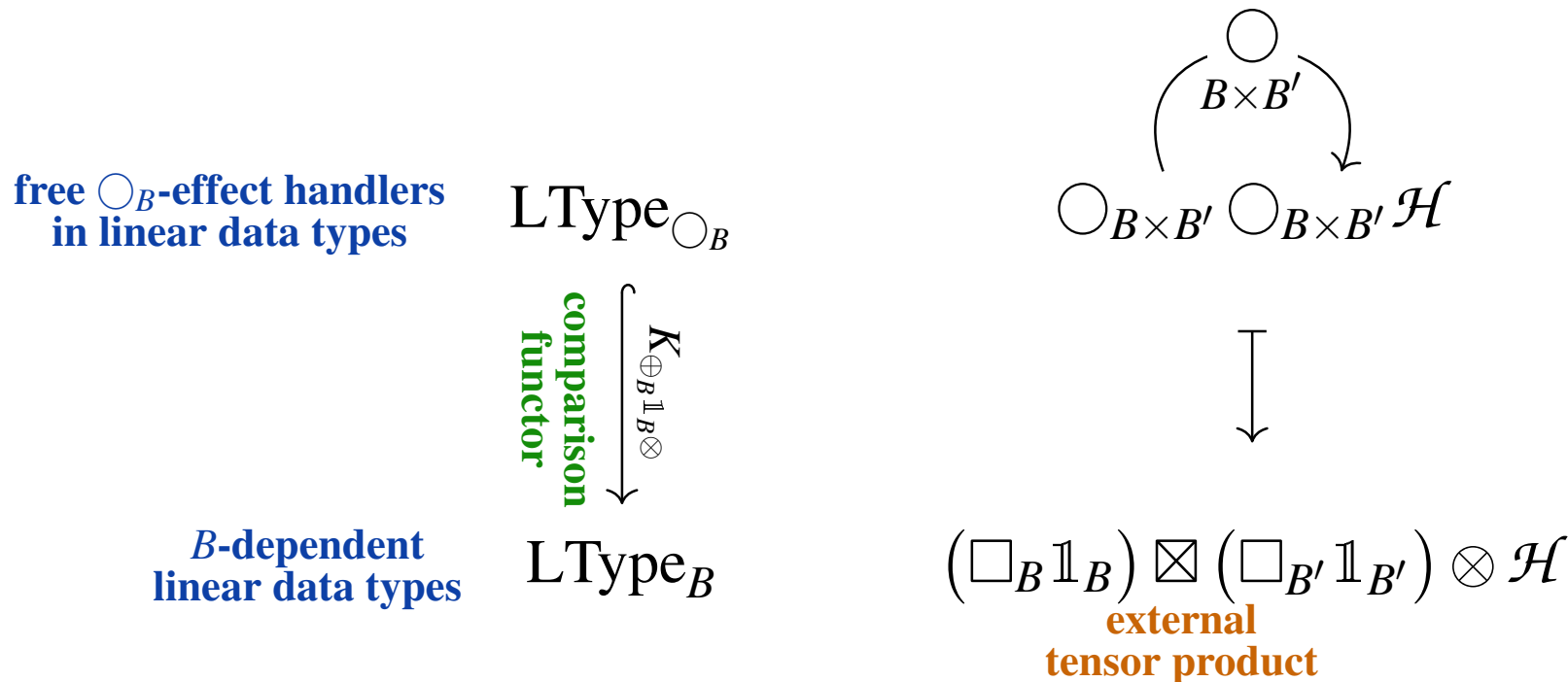
Linear Randomness and Indefiniteness are “very strong” effects, in that:

$$\circlearrowleft_B(D \otimes D') \simeq (\circlearrowleft_B D) \otimes D', \quad \star_B(D \otimes D') \simeq (\star_B D) \otimes D'$$

There is a whole system of them:

$$\circlearrowleft_B \circlearrowleft_{B'} \simeq \circlearrowleft_{B \times B'}, \quad \text{NB: } \circlearrowleft_B \circlearrowleft'_B \simeq \circlearrowleft_B \mathbb{1} \otimes \circlearrowleft'_B$$

which under dynamic lifting (monadicity comparison functor) gives the external tensor product of dependent linear types:

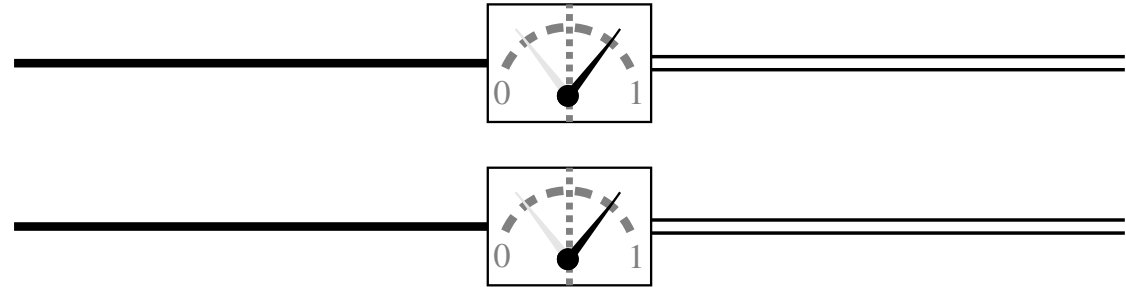


Quantum circuits with classical control & effects

are the *effectful* string diagrams in the linear type system

E.g.

The dependent linear type of a measurement on a pair of qbits:



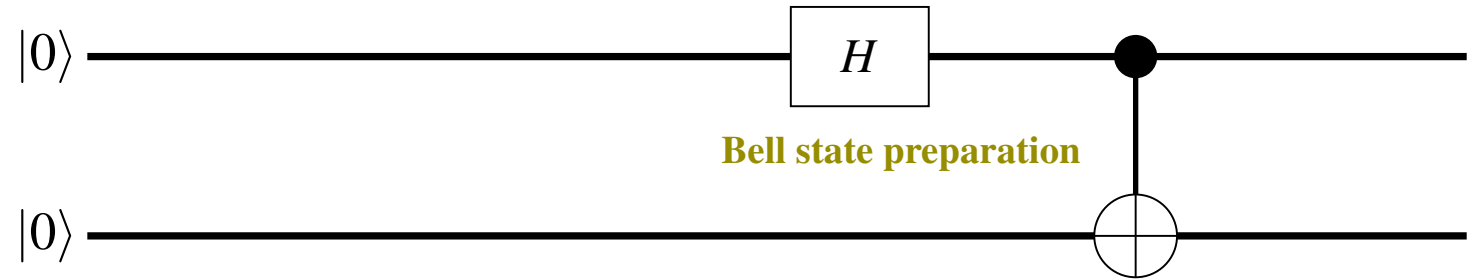
$$\begin{array}{ccc}
 \text{type of a pair of coherent qbits} & \text{pair of measurements} & \text{type of collapsed qbits dependent on measured bits } b, b' \\
 \square_{\text{Bit}^2}(\text{QBit}_\bullet \boxtimes \text{QBit}_\bullet) & \xrightarrow{\varepsilon_{\text{Bit}^2}(\text{QBit}_\bullet \boxtimes \text{QBit}_\bullet)} & \text{QBit}_\bullet \boxtimes \text{QBit}_\bullet
 \end{array}$$

measured bits

$$(b, b') : \text{Bit}^2 \vdash \square_{\text{Bit}^2}(\text{QBit}_\bullet \boxtimes \text{QBit}_\bullet)_{(b, b')} \simeq \mathbb{C}^2 \otimes \mathbb{C}^2 \xrightarrow{\sum_{d, d'} q_{dd'} |d\rangle \otimes |d'\rangle \mapsto q_{bb'} |b\rangle \otimes |b'\rangle} \mathbb{C}.$$

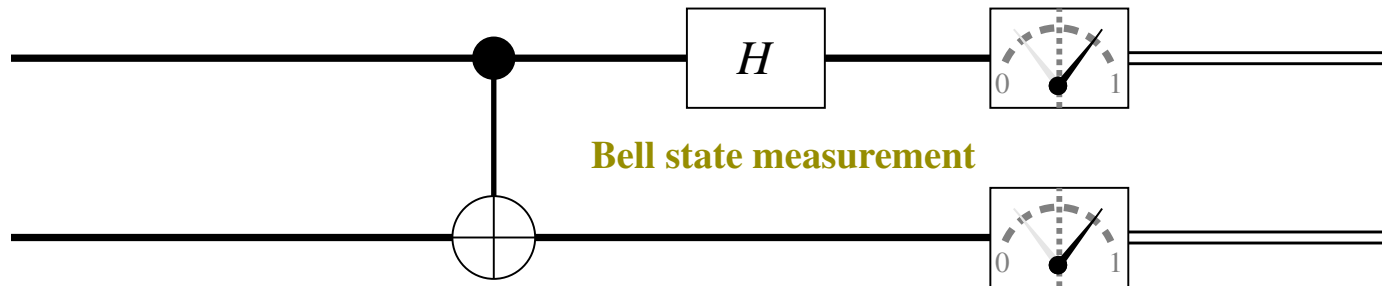
collapse of the quantum state

Example: Bell states of q-bits are typed as follows (regarded in $LType_{\text{Bit} \times \text{Bit}}$):



$$\text{QBit}_{\bullet} \boxtimes \text{QBit}_{\bullet} \rightarrow (\diamond_{\text{Bit}} \text{QBit}_{\bullet}) \boxtimes (\diamond_{\text{Bit}} \text{QBit}_{\bullet}) \simeq \square_{\text{Bit}^2} (\text{QBit}_{\bullet} \boxtimes \text{QBit}_{\bullet}) \rightarrow \square_{\text{Bit}^2} (\text{QBit}_{\bullet} \boxtimes \text{QBit}_{\bullet})$$

$$b, b' : \text{Bit} \vdash \mathbb{C} \xrightarrow{1 \mapsto |0\rangle \otimes |0\rangle} \xrightarrow{\frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \otimes |0\rangle} \xrightarrow{\frac{1}{\sqrt{2}} (|0\rangle \otimes |0\rangle + |1\rangle \otimes |1\rangle)} \mathbb{C}^2 \otimes \mathbb{C}^2$$



$$\square_{\text{Bit}^2} (\text{QBit}_{\bullet} \boxtimes \text{QBit}_{\bullet}) \longrightarrow \text{QBit}_{\bullet} \boxtimes \text{QBit}_{\bullet}$$

$$b_1, b_2 : \text{Bit} \vdash \mathbb{C}^2 \otimes \mathbb{C}^2 \xrightarrow{\sum_{b'_1, b'_2} q_{b'_1, b'_2} \cdot |b'_1\rangle \otimes |b'_2\rangle} \xrightarrow{(q_{0, b_2} + (-1)^{b_1} \cdot q_{1, (1-b_2)}) \cdot |b_1\rangle \otimes |b_2\rangle} \mathbb{C}$$

QS – Quantum Systems language @ CQTS

↪ full-blown Quantum Systems language emerges embedded in dLHoTT

Dependent Linear Homotopy Type Theory (dLHoTT)
for universal algorithmic quantum computation

Homotopy Type Theory (HoTT)
for topological logic gates

Quantum Systems Language (QS)
for quantum logic circuits

Topological Quantum Gate Circuits
for realistic quantum computation

*discussed
elsewhere*

*discussed in
this talk*



Center for
Quantum &
Topological
Systems

Quantum Data Types via Linear HoTT

presentation at:

Workshop on Quantum Software @ QTML 2022

Urs Schreiber (NYU Abu Dhabi)

on joint work at CQTS with

D. J. Myers, M. Riley,

and Hisham Sati

slides and further pointers at: ncatlab.org/schreiber/show/QDataInLHoTT#QTML2022