

Allegories and bicategories of relations

Definition 1 ([CW87]). A (locally ordered) *cartesian bicategory* is a locally partially ordered 2-category \mathcal{C} satisfying the following:

1. \mathcal{C} is symmetric monoidal: there is a pseudofunctor $\otimes: \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ together with natural isomorphisms α, λ, ρ and σ satisfying the usual coherence conditions;
2. every object of \mathcal{C} is a commutative comonoid, that is, comes equipped with maps

$$\Delta_X: X \rightarrow X \otimes X \quad t_X: X \rightarrow I$$

whose right adjoints we write Δ_X^*, t_X^* , where I is the tensor unit, satisfying the obvious associativity, symmetry and unitality axioms, and this is the only such comonoid structure on X ;

3. every morphism $r: X \multimap Y$ is a lax comonoid morphism:

$$\Delta_Y \circ r \leq (r \otimes r) \circ \Delta_X \quad t_Y \circ r \leq t_X$$

Proposition 2 ([CW87, theorem 1.6]). *A bicategory \mathcal{C} is cartesian if and only if the following hold:*

1. $\text{Map}(\mathcal{B})$ has finite 2-products (given by \otimes and I).
2. The hom-posets of \mathcal{B} have finite products, and 1_I is the terminal object of $\mathcal{B}(I, I)$.
3. The tensor product defined as

$$r \otimes s = (\pi_1^* r \pi_1) \cap (\pi_2^* s \pi_2)$$

where the π_i are the product projections, is functorial.

Definition 3 ([CW87, def. 2.1]). An object X in a cartesian bicategory is called *Frobenius* (Carboni–Walters say *discrete*) if it satisfies

$$\Delta \circ \Delta^* = (\Delta^* \otimes 1) \circ (1 \otimes \Delta) \tag{1}$$

A *bicategory of relations* is a cartesian bicategory in which every object is Frobenius.

Proposition 4. *A bicategory of relations \mathcal{B} is compact closed, that is, there is an identity-on-objects involution $(-)^{\circ}: \mathcal{B}^{\text{op}} \rightarrow \mathcal{B}$ and a natural isomorphism*

$$\mathcal{B}(X \otimes Y, Z) \cong \mathcal{B}(X, Z \otimes Y)$$

Lemma 5 ([CW87, corollary 2.6]). *In a bicategory of relations*

1. If f is a map then $f^* = f^{\circ}$.
2. If f and g are maps and $f \leq g$ then $f = g$.

Proposition 6. *A bicategory of relations is the same thing as a unitary tabular allegory.*

Only one part of the proof is non-trivial, but we postpone the whole thing until after the necessary lemmas.

Freyd and Šcedrov give a construction [FŠ90, B.3] of the free allegory on a regular theory, which allows us to interpret any formula of regular logic in a unitary pre-tabular allegory (relative to some given interpretation of the basic sorts, terms and predicates). Suppose we have predicates $R(x, y)$ and $S(y, z)$, interpreted as $r: X \multimap Y$ and $s: Y \multimap Z$ respectively. Then their relational composite is given by the formula

$$SR(x, z) = \exists v. z^* R(x, v) \wedge x^* S(v, z)$$

This can be interpreted in two different ways: as the composite $sr: X \multimap Z$, or more ‘literally’ as

$$\begin{array}{ccccccc}
& & & X & \xrightarrow{r} & Y & \\
& & & \nearrow \pi_1 & & \searrow \pi_1^\circ & \\
X & \xrightarrow{\pi_1^\circ} & X \times Z & \xrightarrow{\pi_2} & Z & \xrightarrow{s^\circ} & Y & \xrightarrow{\pi_2^\circ} & Y \times Y & \xrightarrow{\Delta^\circ} & Y & \xrightarrow{\pi} & U & \xrightarrow{\pi^\circ} & Z \\
& & \searrow \pi_2 & & \nearrow s^\circ & & \searrow \pi_2^\circ & & \nearrow \pi_2 & & & & & &
\end{array}$$

where parallel morphisms are combined with \cap .

Proposition 7. *In a unitary pre-tabular allegory, the two interpretations above of a relational composite are equal; that is,*

$$sr = (\pi^\circ \pi \Delta^\circ (\pi_1^\circ r \pi_1 \cap \pi_2^\circ s^\circ \pi_2) \cap \pi_2) \pi_1^\circ$$

Proof. First note that on the left-hand side $sr = sr \cap \top = sr \cap \pi_2 \pi_1^\circ$, and that on the right

$$(\pi^\circ \pi \Delta^\circ (\pi_1^\circ r \pi_1 \cap \pi_2^\circ s^\circ \pi_2) \cap \pi_2) \pi_1^\circ = (\top_{YZ} (r \pi_1 \cap s^\circ \pi_2) \cap \pi_2) \pi_1^\circ$$

where $\top_{YZ} = \pi^\circ \pi: Y \multimap U \multimap Z$ is the top element. In one direction we have

$$\begin{aligned}
sr \cap \pi_2 \pi_1^\circ &= (sr \pi_1 \cap \pi_2) \pi_1^\circ && \text{modular law} \\
&= (sr \pi_1 \cap \pi_2 \cap \pi_2) \pi_1^\circ \\
&\leq (s(r \pi_1 \cap s^\circ \pi_2) \cap \pi_2) \pi_1^\circ && \text{modular law} \\
&\leq (\top_{YZ} (r \pi_1 \cap s^\circ \pi_2) \cap \pi_2) \pi_1^\circ
\end{aligned}$$

In the other,

$$\begin{aligned}
(\top_{YZ} (r \pi_1 \cap s^\circ \pi_2) \cap \pi_2) \pi_1^\circ &\leq (\top_{YZ} s^\circ (sr \pi_1 \cap \pi_2) \cap \pi_2) \pi_1^\circ && \text{modular law} \\
&\leq (\top_{YY} (sr \pi_1 \cap \pi_2) \cap \pi_2) \pi_1^\circ \\
&= (\pi_2 \pi_1^\circ (sr \pi_1 \cap \pi_2) \cap \pi_2) \pi_1^\circ && \top = \pi_2 \pi_1^\circ \\
&= (\pi_2 (\pi_1^\circ sr \pi_1 \cap \pi_1^\circ \pi_2) \cap \pi_2) \pi_1^\circ && \text{maps distribute} \\
&= \pi_2 (\pi_1^\circ sr \pi_1 \cap \pi_1^\circ \pi_2 \cap \pi_2^\circ \pi_2) \pi_1^\circ && \text{modular law} \\
&= \pi_2 (\pi_1^\circ sr \pi_1 \cap (\pi_1^\circ \cap \pi_2^\circ) \pi_2) \pi_1^\circ && \text{maps distribute} \\
&= \pi_2 (\pi_1^\circ sr \pi_1 \cap \Delta \pi_2) \pi_1^\circ && \text{see below} \\
&= \pi_2 \Delta (\Delta^\circ \pi_1^\circ sr \pi_1 \cap \pi_2) \pi_1^\circ && \text{modular law} \\
&= (sr \pi_1 \cap \pi_2) \pi_1^\circ && \Delta \pi_1 = \Delta \pi_2 = 1 \\
&= sr \cap \pi_2 \pi_1^\circ && \text{modular law}
\end{aligned}$$

In the fourth last line we used the fact that $\pi_1 \cap \pi_2 = \Delta^\circ$, which follows from lemma 8 below and the fact that $\Delta = \langle 1, 1 \rangle$. \square

Lemma 8. *Let \mathcal{A} be a unitary pre-tabular allegory. If $A \xleftarrow{f} X \xrightarrow{g} B$ in $\text{Map}(\mathcal{A})$, then $\langle f, g \rangle = \pi_1^\circ f \cap \pi_2^\circ g$ in \mathcal{A} .*

Proof. Write $r = \pi_1^\circ f \cap \pi_2^\circ g$. From the modular law and the fact that product cones tabulate top morphisms it follows that $\pi_1 r = f$ and $\pi_2 r = g$. Thus $r = \langle f, g \rangle$ if and only if r is a map.

For the counit inequality,

$$\begin{aligned} rr^\circ &= (\pi_1^\circ f \cap \pi_2^\circ g)(f^\circ \pi_1 \cap g^\circ \pi_2) \\ &\leq \pi_1^\circ f f^\circ \pi_1 \cap \pi_2^\circ g g^\circ \pi_2 && \text{distrib.} \\ &\leq \pi_1^\circ \pi_1 \cap \pi_2^\circ \pi_2 \\ &= 1 && \text{proj'ns tabulate} \end{aligned}$$

For the unit,

$$\begin{aligned} r^\circ r &= (f^\circ \pi_1 \cap g^\circ \pi_2)(\pi_1^\circ f \cap \pi_2^\circ g) \\ &= (f^\circ \pi_1 \cap g^\circ \pi_2)\pi_1^\circ f \cap (f^\circ \pi_1 \cap g^\circ \pi_2)\pi_2^\circ g && \text{distrib.} \\ &= (f^\circ \cap g^\circ \pi_2 \pi_1^\circ) f \cap (f^\circ \pi_1 \pi_2^\circ \cap g^\circ) g && \text{modular law} \\ &= f^\circ f \cap g^\circ g && g^\circ \pi_2 \pi_1^\circ = \top, \text{ etc.} \\ &\geq 1 \cap 1 = 1 \end{aligned}$$

\square

Lemma 9. *Let $t, u: A \times B \rightleftarrows B$. Then*

$$(t \cap \pi_2)\pi_1^\circ \cap (u \cap \pi_2)\pi_1^\circ = (t \cap u \cap \pi_2)\pi_1^\circ$$

Proof. By the modular law, the left-hand side is

$$((\pi_2 \cap u)\pi_1^\circ \pi_1 \cap t \cap \pi_2)\pi_1^\circ$$

It therefore suffices to show that

$$(\pi_2 \cap u)\pi_1^\circ \pi_1 \cap \pi_2 = u \cap \pi_2$$

By lemma 10 below, we have that $\pi_1^\circ \pi_1 = \pi_{13}\pi_{12}^\circ$, so the left-hand side above is

$$\begin{aligned} (\pi_2 \cap u)\pi_{13}\pi_{12}^\circ \cap \pi_2 &= ((\pi_2 \cap u)\pi_{13} \cap \pi_2 \pi_{12})\pi_{12}^\circ && \text{modular law} \\ &= (\pi_2 \pi_{13} \cap u \pi_{13} \cap \pi_2 \pi_{12})\pi_{12}^\circ && \text{maps distribute} \\ &= (u \pi_{13} \cap \pi_2(A \times \Delta_B)^\circ)\pi_{12}^\circ && \text{lemma 11} \\ &= (u \pi_{13}(A \times \Delta_B) \cap \pi_2)(A \times \Delta_B)^\circ \pi_{12}^\circ && \text{modular law} \\ &= u \cap \pi_2 \end{aligned}$$

\square

Lemma 10. *Let $\pi_1: A \times B \dashrightarrow B$ and take*

$$\pi_{12} = \langle \pi_1, \pi_2 \rangle, \pi_{13} = \langle \pi_1, \pi_3 \rangle: A \times B \times B \dashrightarrow A \times B$$

Then $\pi_1^\circ \pi_1 = \pi_{13}\pi_{12}^\circ$.

Proof. By prop. 7, the left-hand side is

$$\begin{array}{ccccccc}
 & & & & A \times B & & \\
 & & & \nearrow^{\pi_{12}} & & \searrow & \\
 A \times B & \xrightarrow{\pi_{12}^\circ} & A \times B \times A \times B & \xrightarrow{\pi_{34}} & A \times B & \longrightarrow & A \longrightarrow A \times B \\
 & & & \searrow_{\pi_{34}} & & \nearrow_{\pi_{34}} & \\
 & & & & & &
 \end{array}$$

Factoring the double (co)projections through the canonical isomorphism $A \times B \times A \times B \cong A \times A \times B \times B$, we get

$$\begin{aligned}
 (\pi_1^\circ(\pi_1 \cap \pi_2) \cap \pi_{24})\pi_{13}^\circ &= (\pi_1^\circ\pi_1(\Delta \times B \times B)^\circ \cap \pi_{24})\pi_{13}^\circ && \text{lemma 11} \\
 &= (\pi_1^\circ\pi_1 \cap \pi_{13})\pi_{12}^\circ && \text{mod., } \pi\Delta = 1 \\
 &= (\pi_1^\circ\pi_1 \cap \pi_2^\circ\pi_3)\pi_{12}^\circ && \text{lemma 8} \\
 &= \pi_{13}\pi_{12}^\circ && \text{idem}
 \end{aligned}$$

□

Lemma 11. Let $\pi_2, \pi_3: A \times B \times B \rightarrow B$, and write $\pi_{12} = \langle \pi_1, \pi_2 \rangle$, $\pi_{13} = \langle \pi_1, \pi_2 \rangle$ as before. Then

$$\pi_2 \cap \pi_3 = \pi_2(\pi_{12} \cap \pi_{13}) = \pi_2(1_A \times \Delta_B)^\circ$$

where in the middle and on the right $\pi_2: A \times B \rightarrow B$.

Proof. For the left-hand equality we have

$$\begin{aligned}
 \pi_2(\pi_{12} \cap \pi_{13}) &= \pi_2(\pi_1^\circ\pi_1 \cap \pi_2^\circ\pi_2 \cap \pi_3^\circ\pi_3) && \text{lemma 8} \\
 &= \pi_2(\pi_1^\circ\pi_1 \cap \pi_2^\circ\pi_2) \cap \pi_3 && \text{modular law} \\
 &= \pi_2\pi_1^\circ\pi_1 \cap \pi_2 \cap \pi_3 && \text{modular law} \\
 &= \top \cap \pi_2 \cap \pi_3 \\
 &= \pi_2 \cap \pi_3
 \end{aligned}$$

For the right, $1_A \times \Delta_B = \langle \pi_1, \langle 1, 1 \rangle \pi_2 \rangle$, which is

$$\pi_1^\circ\pi_1 \cap \pi_2^\circ\pi_2 \cap \pi_3^\circ\pi_2$$

by lemma 8, the opposite of which composed with π_2 is the first line above. □

Now we may proceed with our postponed proof.

Proof of prop. 6. Both allegories and bicategories of relations are locally partially ordered 2-categories equipped with an identity-on-objects involution.

Suppose \mathcal{B} is a bicategory of relations. The Frobenius law implies the modular law [CW87, remark 2.9(ii)]. The tensor unit I , the terminal object of $\text{Map}(\mathcal{B})$, is a unit: there is a unique map $X \rightarrow I$ for any X , and 1_I is the top element of $\mathcal{B}(I, I)$ by prop. 2. The product projections tabulate the top elements, so \mathcal{B} is pre-tabular.

Conversely, let \mathcal{A} be a unitary pre-tabular allegory, and refer to prop. 2. $\text{Map}(\mathcal{A})$ has finite products, and local finite products are given by the definition of an allegory and the presence of the unit; the identity on the unit is by definition the top element of the relevant hom set.

The functoriality of the tensor product

$$r \otimes s = \pi_1^* r \pi_1 \cap \pi_2^* s \pi_2$$

is the only difficult part. Firstly, on identities we have

$$1 \otimes 1 = \pi_1^* \pi_1 \cap \pi_2^* \pi_2 = 1$$

because the projections are a tabulation. Now consider

$$sr \otimes s'r' = \pi_1^* sr \pi_1 \cap \pi_2^* s'r' \pi_2$$

as in

$$\begin{array}{ccccc} & & X & \xrightarrow{r} & Y & \xrightarrow{s} & Z & & \\ & \nearrow & & & & & & \searrow & \\ X \times X' & & & & & \cap & & & Z \times Z' \\ & \searrow & & & & & & \nearrow & \\ & & X' & \xrightarrow{r'} & Y' & \xrightarrow{s'} & Z' & & \end{array}$$

By prop. 7, this is equal to

$$\begin{array}{ccccccc} & & & & X & \xrightarrow{r} & Y & \longrightarrow & U & & \\ & & & & \nearrow & & \searrow & & & & \\ & & X \times X' \times Z \times Z' & \longrightarrow & Z & \xrightarrow{s^\circ} & Y & \longrightarrow & U & & \\ X \times X' & \nearrow & & & & & & \searrow & & & Z \times Z' \\ & \searrow & & & & & & \nearrow & & & \\ & & X \times X' \times Z \times Z' & \longrightarrow & X' & \xrightarrow{r'} & Y' & \longrightarrow & U & & \\ & & \searrow & & \nearrow & & & & & & \\ & & & & Z' & \xrightarrow{s'^\circ} & Y' & & & & \end{array}$$

where unlabelled arrows are the obvious projections, and parallel arrows are to be combined with \cap as before. By lemma 9 this is the result of precomposing with the coprojection $X \times X' \twoheadrightarrow X \times X' \times Z \times Z'$ the meet of the following morphism with π_2 :

$$\begin{array}{ccccccc} & & & & X & \xrightarrow{r} & Y & & \\ & & & & \nearrow & & \searrow & & \\ X \times X' \times Z \times Z' & \nearrow & & & Z & \xrightarrow{s^\circ} & Y & & \\ & \searrow & & & & & & \searrow & \\ & & & & X' & \xrightarrow{r'} & Y' & & \\ & & & & \nearrow & & \searrow & & \\ & & & & Z' & \xrightarrow{s'^\circ} & Y' & \longrightarrow & U & \longrightarrow & Z \times Z' \end{array}$$

We may use the modular law at U , then the equality

$$Y' \twoheadrightarrow U \twoheadrightarrow Y = Y' \twoheadrightarrow Y \times Y' \twoheadrightarrow Y$$

and then the modular law again to turn the above into

$$\begin{array}{ccccc}
 & & X & \xrightarrow{r} & Y \\
 & \nearrow & & & \searrow \\
 & & Z & \xrightarrow{s^\circ} & Y \\
 X \times X' \times Z \times Z' & & & & \\
 & \searrow & & & \nearrow \\
 & & X' & \xrightarrow{r'} & Y' \\
 & & Z' & \xrightarrow{s'^\circ} & Y'
 \end{array}
 \longrightarrow Y \times Y' \longrightarrow U \longrightarrow Z \times Z'$$

But now we may use the symmetry of \cap to swap the morphism containing s° with that containing r' : the resulting morphism (after \cap ing with π_2 and composing with the coprojection out of $X \times X'$) is exactly the interpretation after prop. 7 of

$$(\pi_1^\circ s \pi_1 \cap \pi_2^\circ s' \pi_1) \circ (\pi_1^\circ r \pi_1 \cap \pi_2^\circ r' \pi_2)$$

Thus \otimes is functorial. □

References

- [CW87] Aurelio Carboni and R. F. C. Walters. Cartesian bicategories I. *Journal of Pure and Applied Algebra*, 49:11–32, 1987.
- [FŠ90] Peter Freyd and André Šcedrov. *Categories, Allegories*. North-Holland, 1990.