Algebras, Operads, Combinads

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HOGT Lille, 23 mars 2012

I. Types of algebras, types of operads

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- II. An example: planar trees and ns operads

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- IV. Another example: surjective maps and permutads

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 - a) Feynman diagrams
 - b) higher operads: opetopes
 - (relationship with quantum groups)

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(thanks to María Ronco and Bruno Vallette)
[JLL-MR] JLL, M.Ronco, Permutads, soumis à J.Algebraic
Combinatorics A.
[JLL-BV] JLL and B.Vallette, Algebraic operads,

Grundlehren Math.Wiss. 346, Springer, Heidelberg, 2012.

I. Types of algebras, types of operads

Types of algebras

Algebra usually means: unital associative algebra

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There are other types of algebras: Lie, Poisson, Jordan, dendriform, Zinbiel, Leibniz, Batalin-Vilkovisky, magmatic, A_{∞} , etc.

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 $\mathcal{P}:\mathrm{Mod}\to\mathrm{Mod}$

1.
$$\mathcal{P}(V) = \bigoplus_{n} \mathcal{P}_{n} \otimes V^{\otimes n}$$

2. $\mathcal{P}(V) = \bigoplus_{n} \mathcal{P}(n) \otimes_{\mathbb{S}_{n}} V^{\otimes n}$
3. $\Gamma \mathcal{P}(V) = \bigoplus_{n} (\mathcal{P}(n) \otimes V^{\otimes n})^{\mathbb{S}_{n}}$ (B.Fresse)

There are several types of operads:

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- ns operads,
- symmetric operads,
- divided power operads
- cyclic operads,
- shuffle operads,
- wheeled operads,
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Motivation: Koszul duality.

$$A =$$
 associative alg $\Rightarrow A^! =$ associative alg
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Given an operad of some type, which type is the Koszul dual?,

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Koszul duality for types of operads

In order to perform Koszul duality for types of operads, we need to know how to encode a type of operads, and, then, to construct a Koszul duality theory for these objects.

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Obviously the relevant object encoding (ns) operads is self-dual, since the Koszul dual of an (ns) operad is an (ns) operad.

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The strategy: look for the free object, discard the variable ex: T(V) gives As

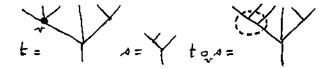
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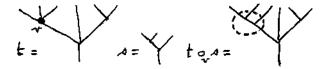
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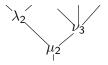
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Key point: substitution is associative.

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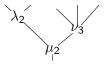
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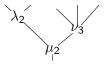


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THM The substitution process defines a monoid structure $\Gamma : \mathbb{PT} \circ \mathbb{PT} \to \mathbb{PT}$ on the endofunctor \mathbb{PT} , hence (\mathbb{PT}, Γ) is a monad on the category of N-modules.

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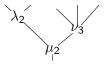


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An algebra over the monad (\mathbb{PT}, Γ) is a **ns operad** (this is the *combinatorial definition*, see for instance [JLL-BV]). **PROP** $\mathbb{PT}(M)$ is the free ns operad on M Other examples of types of operads

Construction of the free object:

- ns operads use planar rooted trees
- symmetric operads use nonplanar rooted trees
- cyclic operads use nonplanar nonrooted trees
- strictly cyclic operads use planar nonrooted trees
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Question: what is the general framework for all these examples?

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Definition of a **combinatorial pattern** X over \mathbb{N} :

X is a set, whose elements are called *trees* (abuse of terminology)

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In the planar case all the sets leaves(t), in(v) are completely determined by their cardinality $n \in \mathbb{N}$.

Definition of a combinad

A combinad is a monad on \mathbb{N} -modules (\mathcal{X}, Γ)

 $\mathcal{X}: \mathbb{N}\text{-}\textit{mod} \to \mathbb{N}\text{-}\textit{mod} \quad, \quad \Gamma: \mathcal{X} \circ \mathcal{X} \to \mathcal{X}$

where ${\mathcal X}$ is induced by a given combinatorial pattern ${\mathbb X}$

 $\mathcal{X}(M)_n =$ span of the \mathbb{X} -trees with n leaves, vertices decorated by M

and the composition Γ is induced by the substitution of the combinatorial pattern.

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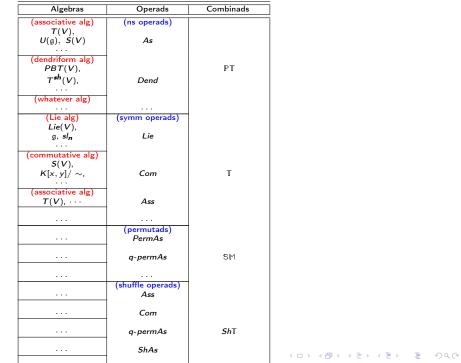
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[JLL-MR]

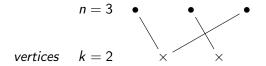
[JLL-MR]

The combinatorial pattern of surjective maps \mathbb{X} : $\mathbb{X}_n = \text{surjective maps } t : \underline{n} \to \underline{k}$ vertices of t: the elements of \underline{k} (shown as \times) inputs of a vertex v of t: the sectors around v (the number is $\#t^{-1}(v) + 1$)

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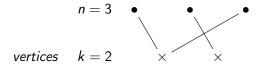


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Substitution: given by composition of surjective maps

[JLL-MR]

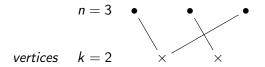
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Substitution: given by composition of surjective maps

An algebra over this combinad is called a *permutad*. This is the *combinatorial definition* of a permutad.

Permutad vs shuffle algebra

Prop A permutad is an \mathbb{N} -module \mathcal{P} equipped with linear maps

$$\bullet_{\gamma}:\mathcal{P}_{n+1}\otimes\mathcal{P}_{m+1}\to\mathcal{P}_{n+m+1}, \text{ for } \gamma\in Sh(n,m),$$

verifying:

$$x \bullet_{\gamma} (y \bullet_{\delta} z) = (x \bullet_{\sigma} y) \bullet_{\lambda} z,$$

whenever $(1_n \times \delta) \cdot \gamma = (\sigma \times 1_r) \cdot \lambda$ in Sh(n, m, r).

For some shuffles: $ullet_{\gamma} = \circ_i$ (consecutive elements in the second set)

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For some shuffles: $\bullet_{\gamma} = \circ_i$ (consecutive elements in the second set) This is the *partial definition* of a permutad. This is essentially the definition given by M. Ronco of a *shuffle algebra*, which are in fact *colored algebras*.

M. Ronco, Shuffle bialgebras, Ann.Inst.Fourier 61 (2011), 799-850.

PermAs and the permutohedron

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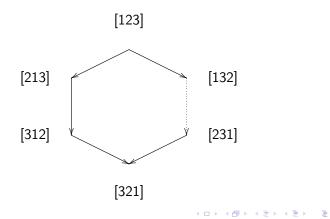
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PermAs and the permutohedron

The analog of As is the permutad generated by a binary operation and the associativity relation, denoted *PermAs*

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Key lemma: connectedness of a subgraph of the weak Bruhat order graph:



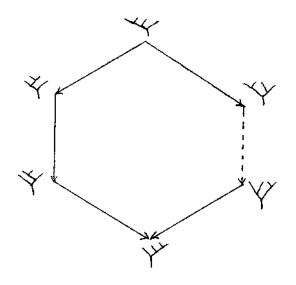
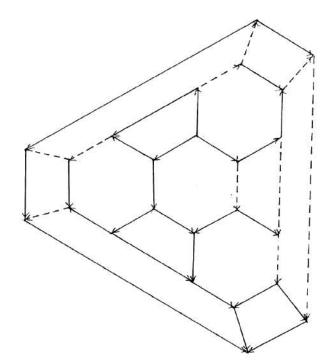


Figure: P^2 and trees



Minimal model of PermAs

 \bullet Minimal model of the ns operad As is A_∞ where

 $(A_{\infty})_n = C_*(associahedron)$

 \bullet Minimal model of the permutad PermAs is \textit{PermAs}_∞ where

 $(PermAs_{\infty})_n = C_*(\text{permutohedron})$

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PROP The permutad with one binary operation and relation

$$(xy)z = q x(yz)$$

is Koszul for any q. (In the operad case, only for $q = 0, 1, \infty$)

V. Further research (work in progress):

a) Feynman diagrams

One can construct a combinad from finite graphs with various decorations **stable by substitution**,

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for instance Feynman graphs (QED, $\varphi^4)$

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Symmetry choice in a combinatorial pattern at vertices with n + 1 flags:

- planar rooted \Rightarrow symmetry group = {1},
- ▶ planar nonrooted \Rightarrow symmetry group = C_{n+1} (cyclic group),
- nonplanar rooted \Rightarrow symmetry group $= S_n$,
- nonplanar nonrooted \Rightarrow symmetry group = S_{n+1} ,
- *p* inputs, *q* outputs \Rightarrow symmetry group $= \mathbb{S}_p \times \mathbb{S}_q$.

More general combinatorial patterns

Let \mathbb{Y} be a combinatorial pattern, for instance •, \mathbb{N} (ladders). A combinatorial pattern \mathbb{X} over \mathbb{Y} is a set X of elements such that each $t \in X$ has an underlying set $|t| \in Y$ (replaces n) Any $t \in X$ comes with its set of vertices $v \in vert(t)$ and its set of *leaves* $leav(t) \in Y$

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Any vertex v comes with its set of *inputs* in(v) \in Y

More general combinatorial patterns

Let \mathbb{Y} be a combinatorial pattern, for instance •, \mathbb{N} (ladders). A combinatorial pattern \mathbb{X} over \mathbb{Y} is a set X of elements such that each $t \in X$ has an underlying set $|t| \in Y$ (replaces n) Any $t \in X$ comes with its set of vertices $v \in vert(t)$ and its set of *leaves* leav $(t) \in Y$ Any vertex v comes with its set of *inputs* in $(v) \in Y$ For any trees t, s and $v \in vert(t)$ and an isomorphism in $(v) \cong leav(s)$ there is given a new tree denoted $t \circ_v s$ such that $|t \circ_v s| = |t|$, $vert(t \circ_v s) = (vert(t) \setminus \{v\}) \cup vert(s)$

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More general combinatorial patterns

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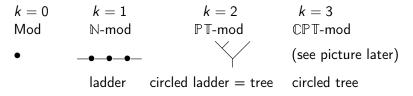
We assume associativity of this composition, that is

I. sequential axiom, for $w \in vert(s)$: $(t \circ_v s) \circ_w r = t \circ_v (s \circ_w r)$ II. parallell axiom, for $w \in vert(t)$: $(t \circ_v s) \circ_w r = (t \circ_w r) \circ_v s$

Combinatorics

ladder = Dynkin diagram A_n

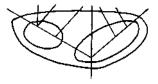
How to continue the sequence:



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Free combinad

The combinatorial pattern is made of circled trees

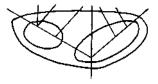


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leaves:	elements	in	Algebras
vertices:	operations	in	Operads
circles:	compositions	in	Combinads

Free combinad

The combinatorial pattern is made of circled trees



leaves:	elements	in	Algebras
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Substitution: as usual

(related to the dendroidal sets of leke Moerdijk and Ittay Weiss)

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b) higher operads: opetopes

Alg=0-op	Op=1-op	Comb=2-op	
		$\mathbb{CPT}(W)$	CPT
	$Mag(M) = \mathbb{PT}(M)$	PT	
T(V) = As(V)	$As = {As_n}_{n \ge 1}$		

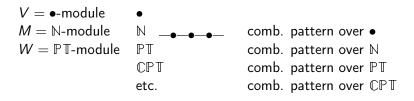
 $V = \bullet \text{-module} \qquad \bullet$ $M = \mathbb{N} \text{-module} \qquad \mathbb{N} \qquad \bullet \text{-} \bullet \text{-} \bullet$ $W = \mathbb{P}\mathbb{T} \text{-module} \qquad \mathbb{P}\mathbb{T} \qquad \mathbb{C}\mathbb{P}\mathbb{T} \qquad \mathbb{C}\mathbb{P}\mathbb{T} \qquad \mathbb{C}\mathbb{P}\mathbb{T}$ etc.

comb. pattern over \bullet comb. pattern over \mathbb{N} comb. pattern over \mathbb{PT} comb. pattern over \mathbb{CPT}

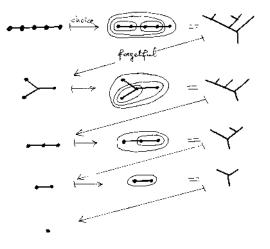
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b) higher operads: opetopes

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			•••
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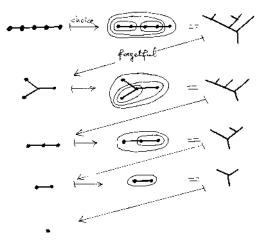
 $T(V), Mag(M), \mathbb{CPT}(W)$ are free objects, As, \mathbb{PT} are associative objects A tower of binary trees



graphs circled graphs rooted trees

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A tower of binary trees = binary opetope



graphs circled graphs rooted trees

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Opetopes

Similar objects already appeared in the literature in the work:

Baez, J.C.; Dolan, J., Higher-dimensional algebra. III. n-categories and the algebra of opetopes. Adv. Math. 135 (1998)

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under the name opetopes.

Opetopes

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under the name opetopes.

Not surprizing since the philosophy is the same, except that Baez, Dolan (and then Joyal, Kock, Batanin, etc.) work in the set environment, while we are working in the linear environment. Main differences: a set is a co-object, because of the diagonal (duplication is possible),

in the linear case there is more freedom (*Lie* is not set-theoretic).

Number of binary opetopes a(n), for A_n ,

n	1	2	3	4	5	6	7
a(n)	1	1	2	10	144		

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 $10=5\times2,\quad 144=(12\times5+2\times6)\times2$

Number of binary opetopes a(n), for A_n ,

n	1	2	3	4	5	6	7
a(n)	1	1	2	10	144		

This is, up to n = 5, the number of regions of linearity for Lusztig's piecewise-linear function in type A_{n-1} appearing in the study the canonical basis of

$$U_q^-(sl_n)$$

Canonical bases for quantized enveloping algebras were introduced by George Lusztig and Masaki Kashiwara in the nineties. Lusztig, G. Canonical bases arising from quantized enveloping algebras. J.Amer.Math.Soc.3 (1990)

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Number of binary opetopes a(n), for A_n ,

n	1	2	3	4	5	6	7
a(n)	1	1	2	10	144	6 608	

This is, up to n = 6, the number of regions of linearity for Lusztig's piecewise-linear function in type A_{n-1} appearing in the study the canonical basis of

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Number of binary opetopes a(n), for A_n ,

n	1	2	3	4	5	6	7
a(n)	1	1	2	10	144	6 608	1 044 736

This is, up to n = 6, the number of regions of linearity for Lusztig's piecewise-linear function in type A_{n-1} appearing in the study the canonical basis of

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CONJECTURE The number of regions of linearity for A_6 is

 $1\,044\,736$

Number of binary opetopes a(n), for A_n , d(n), for D_n , etc.:

n	1	2	3	4	5	6	7
a(n)	1	1	2	10	144	6 608	1 044 736
d(n)	1	1	2	12	184	8 704	1 395 456

This is, up to n = 6, the number of regions of linearity for Lusztig's piecewise-linear function in type A_{n-1} appearing in the study the canonical basis of

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CONJECTURE The number of regions of linearity for A_6 is

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More: a helpful lemma in order to compute a(n) and d(n) is the following. A *partition* of the graph G is: I and J nonempty connected subgraphs of G such that each vertex is either in I or in J. Define $\varphi(G) = \text{sum of "derived graphs"}$ (forgetful \circ choice)

Lemma For any graph G we have

$$\varphi(G) = \sum_{(I,J)} \varphi(I) \vee \varphi(J) ,$$

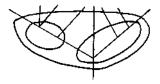
where the sum is over all partitions of G.

This decomposition is close to Carter's method to construct regions of linearity in

Carter, R.W. Canonical bases, reduced words, and Lusztig's piecewise-linear function. ...(1997).

Thanks to Richard Green and Robert Marsh for their help.

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