

Towards Models for 2-Hilb and 3-Hilb as targets for functorial field theories

Joint with André Henriques and Dave Penneys, based on arXiv:2411.01678 and Work in Progress

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W^* -Categories

Algebraic	Unitary Topological
Vector space	Hilbert space
Algebra	von Neumann algebra
Linear Category	W^* -category
Tensor Category	Bicommutant Category

W^* - categories: “categorified” Hilbert Spaces.

Bicommutant categories: “categorified” of von Neumann algebras.

Definition (Dagger Categories)

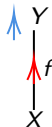
A **dagger category** is a category \mathcal{T} equipped with function $(-)^{\dagger} : \text{Hom}_{\mathcal{T}}(X, Y) \rightarrow \text{Hom}_{\mathcal{T}}(Y, X)$ for all pairs of objects X, Y of \mathcal{T} such that:

- for any object X , $\text{id}_X^{\dagger} = \text{id}_X$.
- for any morphism $f : X \rightarrow Y$, $(f^{\dagger})^{\dagger} = f$
- for any two morphisms f and g such we have $(g \circ f)^{\dagger} = f^{\dagger} \circ g^{\dagger}$.

The most important dagger category which will show up throughout this talk is the category of Hilbert spaces Hilb where the \dagger -structure is given by taking the adjoint of linear maps.

String Calculus for Dagger Categories

We represent $f: X \rightarrow Y$ as



If the direction of f matches the *up* direction (blue arrow) then we read it as f , if it is the opposite we read it as f^\dagger .

We represent composition by concatenation and it is read down to up. For example, the diagram



is read as $X \xrightarrow{f} Y \xrightarrow{g^\dagger} Z$.

W^* -categories \sim “Von-Neumann Algebra-oids”

A $*$ -category is a \mathbb{C} -linear \dagger -category such that the \dagger is anti-linear on hom-spaces.

Let \mathcal{T} be a $*$ -category, we write \mathcal{T}^\oplus for the category whose objects are formal finite sums of objects of \mathcal{T} , and whose morphisms are given by $\text{Hom}_{\mathcal{T}}(\oplus X_i, \oplus X_j) := \oplus_{i,j} \text{Hom}_{\mathcal{T}}(X_i, X_j)$.

Definition (W^* -categories)

A W^* -**category** is a $*$ -category \mathcal{T} such that $\text{End}_{\mathcal{T}}(X)$ is a von Neumann algebra for every $X \in \mathcal{T}^\oplus$.

$\overline{\mathcal{T}}$ is the category with same objects as \mathcal{T} , and complex conjugate hom-spaces. We note this as the first involution (\dagger_0).

Definition (Idempotent completion)

A W^* -category is called **idempotent complete** if whenever a morphism $p : X \rightarrow X$ satisfies $p^2 = p^* = p$, there exists an object Y and a morphism $\iota : Y \rightarrow X$ such that $\iota^* = p$ and $\iota^* \iota = \text{id}_Y$.

Definition (Generating set)

A W^* -category \mathcal{T} is said to *admit a set of generators* if there exists a set of objects such that every non-zero object admits a non-zero map from at least one of the generators. It is said to **admit a generator** if the above set may be taken to consist of a single object.

Definition (Orthogonal direct sums)

Given a collection of objects X_i in a W^* -category indexed by some set I , their **orthogonal direct sum** is an object X equipped with morphisms $\iota_i : X_i \rightarrow X$ satisfying

$$\iota_i^* \iota_j = \delta_{ij} \text{id}_{X_i} \qquad \sum \iota_i^* \iota_i = \text{id}_X.$$

The orthogonal direct sums, if it exists, is denoted $\bigoplus_{i \in I} X_i$.

Here, the infinite sum $\sum_{i \in I} \iota_i^* \iota_i$ is defined as the sup over all finite subsets $I_0 \subset I$ of the finite sums $\sum_{i \in I_0} \iota_i^* \iota_i$.

Definition (Cauchy completion)

We call a W^* -category **(Cauchy) complete** if it admits a generator, is idempotent complete, and has all direct sums.

We write $\mathcal{C}^{\hat{\oplus}}$ for the direct sum completion of the idempotent completion of a W^* -category \mathcal{C} , and call it the *Cauchy completion* of \mathcal{C} .

Theorem

Every complete W^* -category \mathcal{T} is equivalent to $R\text{-Mod}$ for some von Neumann algebra R .

Let X be a generator of \mathcal{T} and $R := \text{End}_{\mathcal{T}}(X)^{\text{op}}$.

$$BR^{\text{op}} \rightarrow \mathcal{T} \text{ and } BR^{\text{op}} \rightarrow R\text{-Mod}$$

$$\star_{R^{\text{op}}} \mapsto X \text{ and } \star_{R^{\text{op}}} \mapsto {}_R L^2 R$$

On Cauchy completion, this extends to equivalences,

$$(BR^{\text{op}})^{\hat{\oplus}} \xrightarrow{\cong} \mathcal{T} \text{ and } (BR^{\text{op}})^{\hat{\oplus}} \xrightarrow{\cong} R\text{-Mod}.$$

Therefore, $\mathcal{T} \cong R\text{-Mod}$.

Let \mathcal{T}_1 and \mathcal{T}_2 be W^* -categories.

A functor between W^* -categories is a $*$ -functor $F : \mathcal{T}_1 \rightarrow \mathcal{T}_2$ that induces normal homomorphisms $\text{End}_{\mathcal{T}_1}(T) \rightarrow \text{End}_{\mathcal{T}_2}(F(T))$ for all $T \in \mathcal{T}_1$.

Theorem

Given von Neumann algebras R_1 and R_2 , the functor

$$\begin{aligned} \text{Bim}(R_2, R_1) &\rightarrow \text{Func}(R_1\text{-Mod}, R_2\text{-Mod}) \\ {}_{R_2}X_{R_1} &\mapsto {}_{R_2}X \boxtimes_{R_1} - \end{aligned}$$

is an equivalence of categories.

The inverse of sends a functor $F : R_1\text{-Mod} \rightarrow R_2\text{-Mod}$ to the bimodule ${}_{R_2}(F({}_{R_1}L^2R_1))_{R_1}$.

Intertwiners correspond to natural transformations.

Monoidal Structure

Given Cauchy complete W^* -categories, their *completed tensor product* is given by:

$$\mathcal{T}_1 \hat{\otimes} \mathcal{T}_2 := (\mathcal{T}_1 \otimes \mathcal{T}_2)^{\hat{\otimes}}.$$

where \otimes is the tensor product of \mathbb{C} -linear categories.

$\text{Hilb} = (\mathbf{BC})^{\hat{\otimes}}$ is the unit of the above operation.

Theorem

Given von Neumann algebras R_1 and R_2 , the functor

$$\begin{aligned} (R_1\text{-Mod}) \hat{\otimes} (R_2\text{-Mod}) &\rightarrow (R_1 \bar{\otimes} R_2)\text{-Mod} \\ ({}_{R_1}H) \otimes ({}_{R_2}K) &\mapsto {}_{R_1 \bar{\otimes} R_2}(H \otimes K) \end{aligned}$$

is an equivalence of categories.

Hilb-valued inner product

Definition (Inner Product)

Every W^* -category \mathcal{T} admits a canonical **Hilb-valued inner product** $\langle -, - \rangle_{\text{Hilb}} : \overline{\mathcal{T}} \times \mathcal{T} \rightarrow \text{Hilb}$ given by:

$$\langle X, Y \rangle_{\text{Hilb}} := p_Y L^2(\text{End}(X \oplus Y)) p_X,$$

where $p_X, p_Y \in \text{End}(X \oplus Y)$ are the two projections.

Lemma

Let \mathcal{T} be a Cauchy complete W^* -category. Then

$$\begin{aligned} \mathcal{T} &\rightarrow \text{Func}(\overline{\mathcal{T}}, \text{Hilb}) \\ X &\mapsto \langle -, X \rangle \end{aligned} \tag{1}$$

is an equivalence of categories. This corresponds to a statement corresponding to the Riesz representation theorem.

Definition (Adjoint of a W^* -functor)

Given a functor $F: \mathcal{T}_1 \rightarrow \mathcal{T}_2$ between Cauchy complete W^* -categories, its **adjoint** F^\dagger is defined as the composite

$$F^\dagger: \mathcal{T}_2 \xrightarrow{\cong} \text{Hom}(\overline{\mathcal{T}_2}, \text{Hilb}) \xrightarrow{-\circ \overline{F}} \text{Hom}(\overline{\mathcal{T}_1}, \text{Hilb}) \xrightarrow{\cong} \mathcal{T}_1.$$

Equivalently, the functor $F^\dagger: D \rightarrow C$ is specified by the requirement that

$$\langle X, F^\dagger(Y) \rangle_{\text{Hilb}} \cong \langle F(X), Y \rangle_{\text{Hilb}}, \quad (2)$$

naturally in X and Y .

Definition (For a natural transformation)

Given a natural transformation $\alpha : F \Rightarrow G$ between functors $\mathcal{T}_1 \rightarrow \mathcal{T}_2$, the adjoint natural transformation $\alpha^\dagger : F^\dagger \Rightarrow G^\dagger$ is specified,

$$\begin{array}{ccc} \langle X, F^\dagger(Y) \rangle_{\text{Hilb}} & \xrightarrow{\simeq} & \langle F(X), Y \rangle_{\text{Hilb}} \\ \text{id}_X, (\alpha^\dagger)_Y \downarrow & & \downarrow \langle \alpha_X, \text{id}_Y \rangle \\ \langle X, G^\dagger(Y) \rangle_{\text{Hilb}} & \xrightarrow{\simeq} & \langle G(X), Y \rangle_{\text{Hilb}} \end{array}$$

commutes.

Lemma

The operations $F \mapsto F^\dagger$ and $\alpha \mapsto \alpha^\dagger$ assemble to an antilinear equivalence

$$\dagger : \text{Func}(\mathcal{T}_1, \mathcal{T}_2) \rightarrow \text{Func}(\mathcal{T}_2, \mathcal{T}_1), \quad (3)$$

and there are natural unitary isos $\varphi_F : F \rightarrow F^{\dagger\dagger}$ and $\nu_{F,G} : F^\dagger \circ G^\dagger \rightarrow (G \circ F)^\dagger$.

We also have $\alpha^* : G \Rightarrow F$ defined pointwise as $(\alpha^*)_X := (\alpha_X)^*$, which gives the antilinear involution $* : \text{Func}(\mathcal{T}_1, \mathcal{T}_2) \rightarrow \text{Func}(\mathcal{T}_1, \mathcal{T}_2)$. This is $(\dagger)_2$.

Theorem

The assignment $R \mapsto R\text{-Mod}$ induces an equivalence of bicategories $\text{Mor}(\text{vN}) \xrightarrow{\cong} \text{W}^*\text{Cat}$

Remark: It is in fact a equivalence of *tri-involutive monoidal* (fully-dagger monoidal) bi-categories.

	vN2	W* Cat
\dagger_0	complex conjugation or op of algebra	complex conjugation or op of category
\dagger_1	complex conjugation of underlying hilbert space	adjoint functor
\dagger_2	adjoint of linear intertwiner	pointwise * (adjoint of linear maps)
	spacial tensor product $\bar{\otimes}$	Complete tensor product of categories $\bar{\otimes}$

2-Hilb-analogies with Hilb

A Hilbert space H	A complete W^* -category \mathcal{T}
The one dimensional Hilbert space \mathbb{C}	The W^* -category Hilb
Scalar multiplication $\cdot : \mathbb{C} \times H \rightarrow H$	Canonical tensor $\text{Hilb} \times \mathcal{T} \rightarrow \mathcal{T}$
Inner product $\langle \cdot, \cdot \rangle : H \times H \rightarrow \mathbb{C}$	Hilb-valued inner product $\overline{\mathcal{T}} \times \mathcal{T} \rightarrow \text{Hilb}$

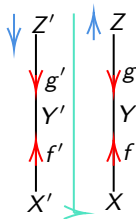
This is just a glimpse, lots more listed in the beginning of [arXiv:2411.01678](https://arxiv.org/abs/2411.01678)

Bi-involutive tensor W^* -categories

String calculus for W^* -categories

We stack the string diagrams for $\overline{\mathcal{T}}$ and \mathcal{T} , separated by a dividing line. This line comes with a co-orientation, which remembers how the inner product diagram is read. The conjugate category $\overline{\mathcal{T}}$ has the opposite local-up direction.

$$\langle X', X \rangle \xrightarrow{\langle f'^{\dagger}, f \rangle} \langle Y', Y \rangle \xrightarrow{\langle g', g^{\dagger} \rangle} \langle Z', Z \rangle$$



There is a unitary isomorphism, $J_{X,Y}: \langle X, Y \rangle \xrightarrow{\sim} \overline{\langle Y, X \rangle}$ natural in X, Y . Diagrammatically, $\overline{\langle \ , \ \rangle}$ swaps the two strings which makes the isomorphism J manifest.

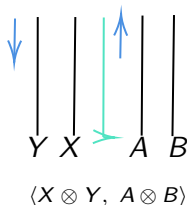
W^* -tensor categories

A * -tensor category $(\mathcal{T}, \otimes, 1, \alpha, l, r)$ is a * -category with a monoidal structure which is compatible in the sense that

$$\otimes : \mathcal{T} \times \mathcal{T} \rightarrow \mathcal{T}$$

is a bilinear functor of * -categories, and the associator α and left and right unitors l, r are unitary.

The string calculus naturally extends to monoidal W^* -categories where the tensor product of objects is represented by placing corresponding strings in parallel left to right in \mathcal{T} and right to left in $\overline{\mathcal{T}}$. The tensor product in \mathcal{T} and $\overline{\mathcal{T}}$ are read outwards starting from the dividing line.



A **bi-involutive W^* -tensor category** is a W^* -tensor category \mathcal{T} equipped with a covariant anti-linear, anti-tensor functor

$$\bar{} : \mathcal{T} \rightarrow \mathcal{T}$$

called the conjugate. The structure data of this anti-tensor functor are denoted

$$\nu_{A,B} : \bar{A} \otimes \bar{B} \xrightarrow{\simeq} \overline{B \otimes A} \quad \text{and} \quad j : \mathbf{1} \rightarrow \bar{\mathbf{1}}$$

and which satisfy some diagrams.

The functor $\bar{}$ is involutive, meaning that for every $A \in \mathcal{T}$, we are given unitary natural unitary isomorphisms

$$\varphi_A : A \rightarrow \bar{\bar{A}}$$

satisfying $\varphi_{\bar{A}} = \overline{\varphi_A}$.

$$\begin{array}{ccc}
\bar{A} \otimes (\bar{B} \otimes \bar{C}) & \xrightarrow{\text{id}_{\bar{A}} \otimes \nu_{B,C}} & \bar{A} \otimes \overline{(C \otimes B)} \\
\alpha_{\bar{A}, \bar{B}, \bar{C}} \uparrow & & \downarrow \nu_{A, C \otimes B} \\
(\bar{A} \otimes \bar{B}) \otimes \bar{C} & & \overline{(C \otimes B) \otimes A} \\
\nu_{A, B} \otimes \text{id}_{\bar{C}} \downarrow & & \downarrow \overline{\alpha_{C, B, C}} \\
\overline{(B \otimes A)} \otimes \bar{C} & \xrightarrow{\nu_{B \otimes A, C}} & \overline{C \otimes (B \otimes A)}
\end{array}$$

$$\begin{array}{ccc}
1 \otimes \bar{A} & \xrightarrow{j \otimes \text{id}_{\bar{A}}} & \bar{1} \otimes \bar{A} & \quad & \bar{A} \otimes 1 & \xrightarrow{\text{id}_{\bar{A}} \otimes j} & \bar{A} \otimes \bar{1} \\
l_{\bar{A}} \downarrow & & \downarrow \nu_{1, A} & & r_{\bar{A}} \downarrow & & \downarrow \nu_{A, 1} \\
\bar{A} & \xleftarrow{r_{\bar{A}}} & \overline{A \otimes 1} & & \bar{A} & \xleftarrow{l_{\bar{A}}} & \overline{1 \otimes A}
\end{array}$$

Finally, we require the compatibility conditions $\varphi_1 = \bar{j} \circ j$ and $\varphi_{A \otimes B} = \overline{\nu_{B, A}} \circ \nu_{\bar{A}, \bar{B}} \circ (\varphi_A \otimes \varphi_B)$.

Bi-involutive tensor functors

A tensor functor $F: \mathcal{T}_1 \rightarrow \mathcal{T}_2$ between bi-involutive $*$ -tensor categories is called a *bi-involutive tensor functor* if it comes equipped with a unitary natural transformation

$$\gamma_X : F(\overline{X}) \rightarrow \overline{F(X)}$$

satisfying the following coherences:

$$\begin{array}{ccccc}
 F(\overline{X}) \otimes F(\overline{Y}) & \xrightarrow{\mu_{\overline{X}, \overline{Y}}} & F(\overline{X} \otimes \overline{Y}) & \xrightarrow{F(\nu_{x,y}^{\mathcal{T}_1})} & F(\overline{Y \otimes X}) \\
 \gamma_X \otimes \gamma_Y \downarrow & & & & \downarrow \gamma_{Y \otimes X} \\
 \overline{F(X)} \otimes \overline{F(Y)} & \xrightarrow{\nu_{F(X) \otimes F(Y)}^{\mathcal{T}_2}} & \overline{F(Y) \otimes F(X)} & \xrightarrow{\mu_{Y \otimes X}} & \overline{F(Y \otimes X)}
 \end{array}$$

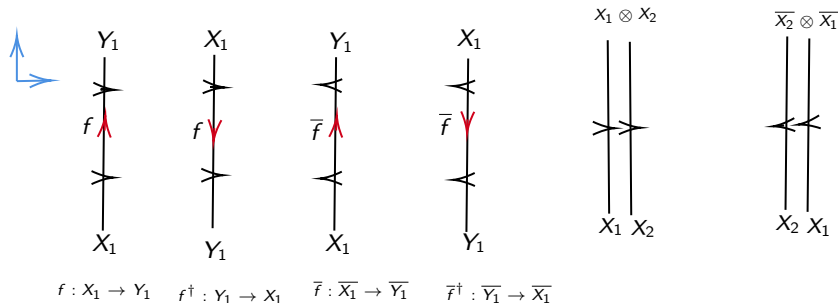
$$\begin{array}{ccc}
 F(x) & \xrightarrow{F(\varphi_X^{\mathcal{T}_1})} & F(\overline{x}) \\
 \varphi_{F(x)}^{\mathcal{T}_2} \downarrow & & \downarrow \gamma_{\overline{x}} \\
 \overline{\overline{F(x)}} & \xleftarrow{\gamma_{\overline{x}}} & \overline{F(x)}
 \end{array}$$

String calculus for bi-involutive categories

Similar to the chosen up-direction, there is a chosen right direction, together these can be thought as the chosen coordinate axes for a string diagram of a bi-involutive tensor category. We equip our objects with a normal vector or a co-orientation.

We then represent \overline{X} by reversing the co-orientation or reflecting along the up-direction.

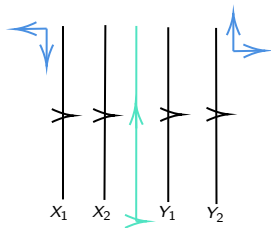
This makes ν and φ with the coherences automatically manifest. The unit object is transparent, so j is also in-built.



For bi-involutive W^* -tensor category

For a bi-involutive tensor category $\overline{\mathcal{T}}$, the local up and right directions are both reversed.

For example, the inner-product $\langle \overline{X_1} \otimes \overline{X_2}, Y_1 \otimes Y_2 \rangle \simeq \langle \overline{X_2} \otimes \overline{X_1}, Y_1 \otimes Y_2 \rangle$ is represented as:



We also have the following isomorphism defined by $\overline{\quad}$ being an anti-linear, anti-tensor functor,

$$\begin{aligned} \langle \overline{X}, \overline{Y} \rangle &= p_{\overline{Y}} L^2(\text{End}(\overline{X} \oplus \overline{Y})) p_{\overline{X}} = p_{\overline{Y}} L^2(\overline{\text{End}(X \oplus Y)}) p_{\overline{X}} \\ &= p_{\overline{Y}} \overline{L^2(\text{End}(X \oplus Y))} p_{\overline{X}} = \overline{p_Y L^2(\text{End}(X \oplus Y))} p_X \\ &= \overline{\langle X, Y \rangle}. \end{aligned}$$

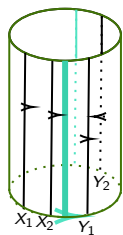
We define an isomorphism $c_{X,Y}: \langle X, Y \rangle \rightarrow \langle \overline{Y}, \overline{X} \rangle$ as the following composition,

$$\langle X, Y \rangle \xrightarrow{J} \overline{\langle Y, X \rangle} \rightarrow \langle \overline{Y}, \overline{X} \rangle$$

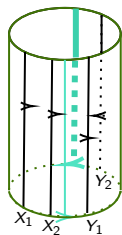
and $\tilde{c}_{X,Y}: \langle \overline{X}, Y \rangle \rightarrow \langle \overline{Y}, X \rangle$ as the following composition,

$$\langle \overline{X}, Y \rangle \rightarrow \langle \overline{Y}, \overline{\overline{X}} \rangle \rightarrow \langle \overline{Y}, X \rangle$$

using φ_X in the last arrow.



$$\begin{array}{c}
 \begin{array}{cccc}
 \text{---} & \text{---} & \text{---} & \text{---} \\
 \nearrow & \nearrow & \nearrow & \nearrow \\
 X_1 & X_2 & Y_1 & Y_2 \\
 \hline
 \langle X_1 \otimes X_2, Y_1 \otimes Y_2 \rangle
 \end{array}
 & \xrightarrow{\tilde{c}_{X_1 \otimes X_2, Y_1 \otimes Y_2}} &
 \begin{array}{cccc}
 \text{---} & \text{---} & \text{---} & \text{---} \\
 \nearrow & \nearrow & \nearrow & \nearrow \\
 Y_1 & Y_2 & X_1 & X_2 \\
 \hline
 \langle Y_1 \otimes Y_2, X_1 \otimes X_2 \rangle
 \end{array}
 \end{array}$$



Commutant of a tensor category

Relative Centre

Let $\iota : \mathcal{T} \rightarrow \mathcal{E}$ be a tensor functor between W^* -tensor categories. The *unitary commutant* $\mathcal{Z}(\iota : \mathcal{T} \rightarrow \mathcal{E})$ of \mathcal{T} inside \mathcal{E} (denoted, $\mathcal{Z}(\iota)$) is the category whose objects are pairs $(X, \{e_X\})$, where X is an object of \mathcal{E} and

$$\{e_X\} = (e_{X,Y} : X \otimes \iota Y \rightarrow \iota Y \otimes X)_{Y \in \mathcal{T}}$$

is a collection of unitary isomorphisms, called a half-braiding. The half-braiding is required to be natural in Y , and to satisfy the following in \mathcal{E} for every $Y, Z \in \mathcal{T}$:

$$\begin{array}{ccc}
 & \iota Y \otimes X \otimes \iota Z & \\
 e_{X,Y} \otimes \text{id}_{\iota Z} \nearrow & & \searrow \text{id}_{\iota Y} \otimes e_{X,Z} \\
 X \otimes \iota Y \otimes \iota Z & & \iota Y \otimes \iota Z \otimes X \\
 \downarrow \sim & & \downarrow \sim \\
 X \otimes \iota(Y \otimes Z) & \xrightarrow{e_{X,Y \otimes Z}} & \iota(Y \otimes Z) \otimes X
 \end{array}$$

Absorbing objects

An object $\Omega \in \mathcal{T}$ is *left absorbing* (*right absorbing*) if it is a non-left(right)-zero-divisor and for every non-left(right)-zero-divisor $X \in \mathcal{T}$ we have $\Omega \otimes X \cong \Omega$ ($X \otimes \Omega \cong \Omega$).

The object is *absorbing* if it is both left absorbing and right absorbing.

The absorbing subcategory \mathcal{T}^{abs} is the completion of the full subcategory on absorbing objects.

Lemma

Given a tensor functor $\iota: \mathcal{T} \rightarrow \mathcal{E}$, for a complete W^* -tensor category which admits weakly absorbing objects, the map $\mathcal{Z}(\iota) \rightarrow \mathcal{Z}(\iota|_{\mathcal{T}^{\text{abs}}})$ is fully-faithful.

Bicommutant Categories

Definition of Bicommutant Category

Let \mathcal{T} be a Cauchy complete bi-involutive W^* -tensor category that admits absorbing objects. This gives us two dagger-monoidal functors,

$$\begin{aligned} L: \mathcal{T} &\rightarrow \text{End}(\mathcal{T}^{\text{abs}}) \text{ given by } X \mapsto X \otimes - \\ R: \mathcal{T}^{\text{mop}} &\rightarrow \text{End}(\mathcal{T}^{\text{abs}}) \text{ given by } X \mapsto - \otimes X \end{aligned}$$

A *bicommutant category* \mathcal{T} is such a category equipped with a unitary natural isomorphism $\gamma^L: L(-) \Rightarrow L(-)^\dagger$ which make L a bi-involutive tensor W^* -functor. Using the associators of \mathcal{T} , L and R induce maps

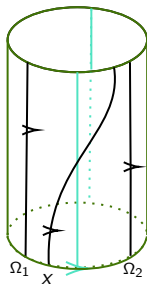
$$\begin{aligned} \mathcal{T} &\rightarrow \mathcal{Z}(R: \mathcal{T}^{\text{mop}} \rightarrow \text{End}(\mathcal{T}^{\text{abs}})) \\ \mathcal{T}^{\text{mop}} &\rightarrow \mathcal{Z}(L: \mathcal{T} \rightarrow \text{End}(\mathcal{T}^{\text{abs}})) \end{aligned}$$

We require these to be equivalences. We require a lot of diagram to commute, which we now list.

For all $X \in \mathcal{T}$, $\Omega_1, \Omega_2 \in \mathcal{T}^{\text{abs}}$, γ^L induces a map

$$\gamma_X^L : \langle \overline{\Omega_1 \otimes X}, \Omega_2 \rangle_{\text{Hilb}} \rightarrow \langle \overline{\Omega_1}, X \otimes \Omega_2 \rangle_{\text{Hilb}}$$

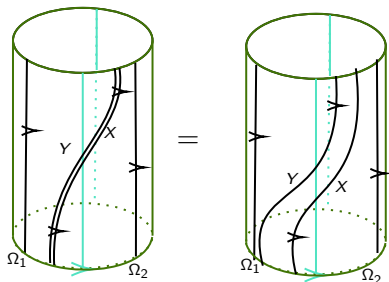
We represent this map diagrammatically by:



The coherences for L being bi-involutive tensor functors are the following:

$$\begin{array}{ccc}
 \langle \overline{X} \otimes \overline{Y} \otimes \overline{\Omega}_1, \Omega_2 \rangle & \xrightarrow{\gamma_X^L} & \langle \overline{Y} \otimes \overline{\Omega}_1, X \otimes \Omega_2 \rangle \\
 \nu_{X,Y} \downarrow & & \downarrow \gamma_Y^L \\
 \langle \overline{Y} \otimes \overline{X} \otimes \overline{\Omega}_1, \Omega_2 \rangle & \xrightarrow{\gamma_{Y \otimes X}^L} & \langle \overline{\Omega}_1, Y \otimes X \otimes \Omega_2 \rangle
 \end{array}$$

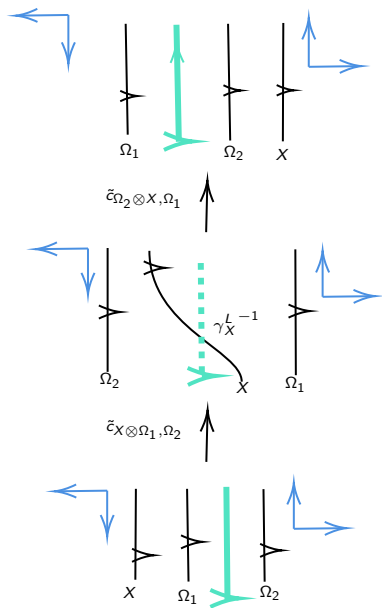
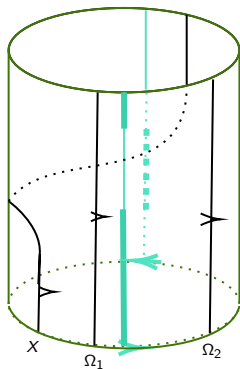
which is manifestly encoded by the diagrammatic calculus.



Using γ^L , we *define* γ^R :

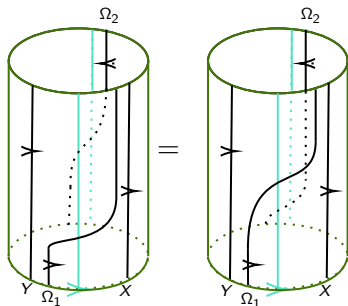
$$\begin{array}{ccc}
 \langle \overline{X \otimes \Omega_1}, \Omega_2 \rangle & \xrightarrow{\gamma_X^R} & \langle \overline{\Omega_1}, \Omega_2 \otimes X \rangle \\
 \tilde{c}_{X \otimes \Omega_1, \Omega_2} \downarrow & & \uparrow \tilde{c}_{\Omega_2 \otimes X, \Omega_1} \\
 \langle \overline{\Omega_2}, X \otimes \Omega_1 \rangle & \xrightarrow{(\gamma_X^L)^{-1}} & \langle \overline{\Omega_2 \otimes X}, \Omega_1 \rangle
 \end{array}$$

Diagrammatically, this involves changing the preferred dividing line, performing γ_X^L and changing the preferred line again.



This makes R a bi-involutive tensor functor.

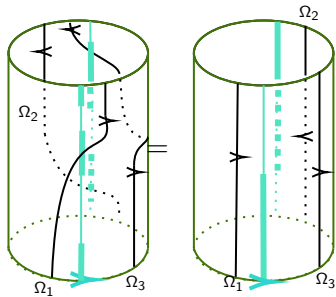
$$\begin{array}{ccc}
 \langle \overline{Y \otimes \Omega_1 \otimes X}, \Omega_2 \rangle & \xrightarrow{\gamma_X^L} & \langle \overline{Y \otimes \Omega_1}, X \otimes \Omega_2 \rangle \\
 \downarrow \gamma_Y^R & & \downarrow \gamma_Y^R \\
 \langle \overline{\Omega_1 \otimes X}, \Omega_2 \otimes Y \rangle & \xrightarrow{\gamma_X^L} & \langle \overline{\Omega_1}, X \otimes \Omega_2 \otimes Y \rangle
 \end{array}$$



Finally, when all three of $\Omega_1, \Omega_2, \Omega_3$ are absorbing, we require the following diagram to commute.

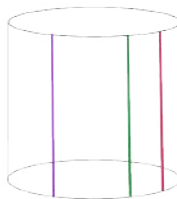
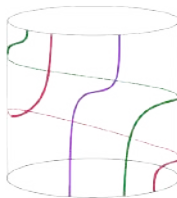
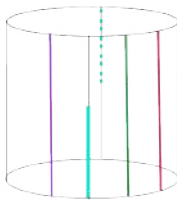
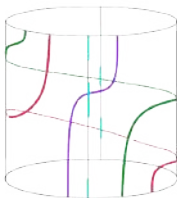
$$\begin{array}{ccc}
 \langle \overline{\Omega_1}, \Omega_3 \otimes \Omega_2 \rangle & \xrightarrow{\gamma_{\Omega_2}^{R^{-1}}} & \langle \overline{\Omega_2 \otimes \Omega_1}, \Omega_3 \rangle \\
 \tilde{c}_{\Omega_1, \Omega_2 \otimes \Omega_3} \downarrow & & \downarrow \gamma_{\Omega_1}^L \\
 \langle \overline{\Omega_3 \otimes \Omega_2}, \Omega_1 \rangle & \xleftarrow{\gamma_{\Omega_3}^{R^{-1}}} & \langle \overline{\Omega_2}, \Omega_1 \otimes \Omega_3 \rangle
 \end{array}$$

We represent this by two morphisms, the LHS of the diagram below indicates the morphism $\gamma_{\Omega_3}^{R^{-1}} \circ \gamma_{\Omega_1}^L \circ \gamma_{\Omega_2}^{R^{-1}}$, and the RHS is the morphism $\tilde{c}_{\Omega_1, \Omega_2 \otimes \Omega_3}$.



Strings on Cylinder

We can put the dividing line anywhere and evaluate, as long as there is at least one string on each side of it.



Concluding Remarks

Examples

$\mathcal{T} = (\mathcal{T}, \otimes, \mathbb{I}, \alpha, l, r, -, \nu, j, \varphi, \gamma^L)$ be a bicommutant category. Some examples include:

(Hilb, \otimes)

$(\text{Bim}(R), \boxtimes_R)$ for some von Neumann algebra R .

$(\text{Hilb}[G])$ for a discrete group G

$(\text{Rep}_{\text{soliton}}(\mathcal{A}), \boxtimes_{\mathcal{A}})$ for a conformal net \mathcal{A}

- When \mathcal{A} is the WZW net for a compact connected group G at level k , this is $\text{Rep}_k(\Omega G)$
- When \mathcal{A} is the Virasoro net, this is $\text{Rep}_c(\text{Diff}(S^1)), \boxtimes$ at some fixed central charge $c \in \{16/m(m+1) : m \geq 2\} \cup [1, \infty)$.

Work in Progress includes *understanding modules* over Bicommutant categories and their “categorified” Connes-fusion. A candidate definition is, given M, N right and left \mathcal{T} -modules respectively, we can define their fusion as

$$M \boxtimes_{\mathcal{T}} N = p_N \text{End}_{\mathcal{T}\text{-Mod}}(\overline{M} \oplus N)^{\text{abs}} p_{\overline{M}}$$

Constructing the 0 and 1 piece of a Segal (Functorial) Chiral CFT, as a functor,

$$\text{Cob}_{0,1,2}^{\text{conf}} \rightarrow \text{Mor}(\text{BicommCat})$$

where $\text{Mor}(\text{BicommCat})$ may serve as 3-Hilb in place of $\text{Mor}(\text{TensCat})$ which is usually taken as 3-Vect. We hope to have strictly more fully-dualisable objects and hence it can serve as targets for unitary 3d-TQFTs which were previously known to not fully-extend to a point.

Thank You

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