RECONSTRUCTION THEORY

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ABSTRACT

Suppose that G is a compact group. Denote by <u>Rep</u> G the category whose objects are the continuous finite dimensional unitary representations of G and whose morphisms are the intertwining operators -- then <u>Rep</u> G is a monoidal *-category with certain properties P_1, P_2, \ldots . Conversely, if <u>C</u> is a monoidal *-category possessing properties P_1, P_2, \ldots , can one find a compact group G, unique up to isomorphism, such that <u>Rep</u> G "is" <u>C</u>? The central conclusion of reconstruction theory is that the answer is affirmative.

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§1. MONOIDAL CATEGORIES

Given categories C,D, their product is the category C \times D defined by

 $Ob(\underline{C} \times \underline{D}) = Ob \underline{C} \times Ob \underline{D}$ $Mor((X,Y), (X',Y')) = Mor(X,X') \times Mor(Y,Y')$ $id_{X} \times Y = id_{X} \times id_{Y'}$

with composition

$$(f',g') \circ (f,g) = (f' \circ f,g' \circ g).$$

Now take $\underline{C} = \underline{D}$ -- then a <u>monoidal category</u> is a category \underline{C} equipped with a functor $\underline{\Omega}:\underline{C} \times \underline{C} \rightarrow \underline{C}$ (the <u>multiplication</u>) and an object $e \in Ob \underline{C}$ (the <u>unit</u>), together with natural isomorphisms R, L, and A, where

and

$$A_{X,Y,Z}: X \otimes (Y \otimes Z) \rightarrow (X \otimes Y) \otimes Z,$$

subject to the following assumptions.

(MC₁) The diagram

commutes.

(MC₂) The diagram

commutes.

[Note: The "coherency" principle then asserts that "all" diagrams built up from instances of R, L, A (or their inverses), and id by repeated application of @ necessarily commute. In particular, the diagrams

commute and $L_e = R_e : e \otimes e \rightarrow e.$]

N.B. Technically, the categories

$$\begin{bmatrix} - & \underline{C} \times (\underline{C} \times \underline{C}) \\ - & (\underline{C} \times \underline{C}) \times \underline{C} \end{bmatrix}$$

are not the same so it doesn't quite make sense to say that the functors

$$\underline{C} \times (\underline{C} \times \underline{C}) \rightarrow \underline{C} \begin{bmatrix} - & (X, (Y, Z)) \rightarrow X & (Y & Z) \\ & & (f, (g, h)) \rightarrow f & (g & h) \end{bmatrix}$$

$$(\underline{C} \times \underline{C}) \times \underline{C} \rightarrow \underline{C}$$

$$\begin{bmatrix} - & ((X,Y),Z) \rightarrow (X \otimes Y) \otimes Z \\ & \\ & \\ & \\ & \\ & \\ & \\ & ((f,g),h) \rightarrow (f \otimes g) \otimes h \end{bmatrix}$$

are naturally isomorphic. However, there is an obvious isomorphism

$$\underline{\mathbf{C}} \times (\underline{\mathbf{C}} \times \underline{\mathbf{C}}) \xrightarrow{\iota} (\underline{\mathbf{C}} \times \underline{\mathbf{C}}) \times \underline{\mathbf{C}}$$

and the assumption is that A:F \rightarrow G \circ ι is a natural isomorphism, where

$$\begin{array}{c} \underline{C} \times (\underline{C} \times \underline{C}) & \stackrel{F}{\rightarrow} \underline{C} \\ & \iota \downarrow \\ (\underline{C} \times \underline{C}) \times \underline{C} & \stackrel{+}{\rightarrow} \underline{C} \\ & & & & \\ \end{array}$$

Accordingly,

 $\forall (X, (Y, Z)) \in Ob \subseteq \times (\underline{C} \times \underline{C})$

and

$$\forall (f,(g,h)) \in Mor \ \underline{C} \times (\underline{C} \times \underline{C}),$$

the square

$$\begin{array}{c|c} X \otimes (Y \otimes Z) & \xrightarrow{A_{X,Y,Z}} \\ f \otimes (g \otimes h) & \downarrow & \downarrow & (f \otimes g) \otimes h \\ & X' \otimes (Y' \otimes Z') & \xrightarrow{A_{X',Y',Z'}} & (X' \otimes Y') \otimes Z' \end{array}$$

commutes.

Interchange Principle If

$$f \in Mor(X, X^{t})$$
$$g \in Mor(Y, Y^{t}),$$

then

$$(f \otimes id_{Y'}) \circ (id_{X} \otimes g) = f \otimes g = (id_{X'} \otimes g) \circ (f \otimes id_{Y'})$$

[Note: Since $\mathfrak{Q}: \underline{C} \times \underline{C} \rightarrow \underline{C}$ is a functor, in general

$$(\mathbf{f} \circ \mathbf{f}') \otimes (\mathbf{g} \circ \mathbf{g}') = (\mathbf{f} \otimes \mathbf{g}) \circ (\mathbf{f}' \otimes \mathbf{g}'),$$

1.1 <u>EXAMPLE</u> Given a field <u>k</u>, let <u>VEC</u> be the category whose objects are the vector spaces over <u>k</u> and whose morphisms are the linear transformations -- then $\underline{\text{VEC}}_{\underline{k}}$ is monoidal: Take X **Q** Y to be the algebraic tensor product and let e be <u>k</u>.

[Note: If

then

$$\otimes$$
 (f,q) = f \otimes q:X \otimes Y \rightarrow X' \otimes Y'

sends $x \otimes y$ to $f(x) \otimes g(y)$.]

Let H and K be complex Hilbert spaces -- then their algebraic tensor product $H \otimes K$ can be equipped with an inner product given on elementary tensors by

$$\langle \mathbf{x}_1 \otimes \mathbf{y}_1, \mathbf{x}_2 \otimes \mathbf{y}_2 \rangle = \langle \mathbf{x}_1, \mathbf{x}_2 \rangle \langle \mathbf{y}_1, \mathbf{y}_2 \rangle$$

and its completion $H \ \underline{\otimes} \ K$ is a complex Hilbert space.

N.B. If

$$\begin{bmatrix} \mathbf{A} \in \mathcal{B}(\mathcal{H}_1, \mathcal{H}_2) \\ \\ \mathbf{B} \in \mathcal{B}(\mathcal{K}_1, \mathcal{K}_2), \end{bmatrix}$$

then

$$A \otimes B: H_1 \otimes K_1 \rightarrow H_2 \otimes K_2.$$

extends by continuity to a bounded linear operator

$$A \stackrel{\alpha}{=} B: H_1 \stackrel{\alpha}{=} K_1 \rightarrow H_2 \stackrel{\alpha}{=} K_2.$$

Denote by <u>HILB</u> the category whose objects are the complex Hilbert spaces and whose morphisms are the bounded linear operators.

1.2 <u>EXAMPLE HILB</u> is a monoidal category. PROOF Define a functor

Q:HILB × HILB → HILB

by

 $\mathfrak{Q}(H,K) = H \mathfrak{Q} K$

amī

$$\mathbf{\underline{a}}(H_1 \xrightarrow{A} H_2, K_1 \xrightarrow{B} K_2) = A \underline{\underline{a}} B$$

and let e be C.

1.3 <u>REMARK</u> Both $\underline{\text{Vec}}_{\underline{k}}$ and $\underline{\text{HILB}}$ admit a second monoidal structure: Take for the multiplication the direct sum Θ and take for the unit the zero object {0}.

Put

$$M(C) = Mor(e, e)$$
.

Then M(C) is a monoid with categorical composition as monoid multiplication.

1.4 LEMMA The monoid M(C) is commutative.

PROOF Take s,t \in M(C) and consider the commutative diagram



Then

$$R_e^{-1} \circ (s \circ t) \circ R_e = R_e^{-1} \circ (t \circ s) \circ R_e$$

=>
 $s \circ t = t \circ s.$

Given $f\in Mor\left(X,Y\right)$ and $s\in \underline{M}(\underline{C})\,,$ define s f to be the composition

$$\begin{array}{ccc} L^{-1} & \text{sof} & L \\ X & \longrightarrow & e \text{ os } X & \longrightarrow & e \text{ os } Y & \longrightarrow & Y. \end{array}$$

1.5 LEMMA We have

$$id_e \cdot f = f$$

$$s \cdot (t \cdot f) = (s \circ t) \cdot f$$

$$(t \cdot g) \circ (s \cdot f) = (t \circ s) \cdot (g \circ f)$$

$$(s \cdot f) \otimes (t \cdot g) = (s \circ t) \cdot (f \otimes g).$$

A monoidal category <u>C</u> is said to be <u>strict</u> if R, L, and A are identities. So, if <u>C</u> is strict, then

 and

 $\begin{bmatrix} X & e & e & x \\ e & x & x & e & x \end{bmatrix}$

[Note: While monoidal, neither $\underline{VEC}_{\underline{k}}$ nor <u>HILB</u> is strict monoidal.] <u>N.B.</u> Take <u>C</u> strict and consider $\underline{M}(\underline{C})$ -- then \forall f,g $\in \underline{M}(\underline{C})$,

$$\mathbf{f} \otimes \mathbf{g} = \mathbf{f} \circ \mathbf{g} = \mathbf{g} \circ \mathbf{f} = \mathbf{g} \otimes \mathbf{f}.$$

1.6 <u>EXAMPLE</u> Let § be the category whose objects are the nonnegative integers and whose morphisms are specified by the rule

composition in Mor(n, n) being group multiplication in S_n . Define

on objects by

Q(n,m) = n + m

and on morphisms by

$$\underline{\boldsymbol{\sigma}} (\mathbf{n} \rightarrow \mathbf{n}, \ \mathbf{m} \rightarrow \mathbf{m}) = \rho_{\mathbf{n},\mathbf{m}} (\boldsymbol{\sigma}, \tau) ,$$

where

$$\rho_{n,m}:\mathfrak{F}_n \times \mathfrak{F}_m \to \mathfrak{F}_{n+m}$$

is the canonical map, i.e.,

 $\rho_{n,m}(\sigma,\tau) = \begin{bmatrix} -1 & 2 & \dots & n & n+1 & n+2 & \dots & n+m \\ & & & & & \\ & & & & \\ & & & & &$

and let e = 0 -- then with these choices, 3 is a strict monoidal category.

[Note: *\$* is equivalent to the category whose objects are the finite sets and whose morphisms are the bijective maps.]

1.7 EXAMPLE Let $\underline{MAT}_{\underline{k}}$ be the category whose objects are the positive integers and whose morphisms are specified by the rule

$$Mor(n,m) = M_{n,m}(\underline{k}),$$

the n-by-m matrices with coefficients in <u>k</u>. Here $id_n: n \rightarrow n$ is the unit diagonal n-by-n matrix and composition

$$\circ:Mor(n,m) \times Mor(m,p) \rightarrow Mor(n,p)$$

is

$$B \circ A = AB,$$

the product on the right being ordinary multiplication of matrices. Define

$$\mathbf{Q}: \underline{\mathsf{MAT}}_{\underline{k}} \times \underline{\mathsf{MAT}}_{\underline{k}} \to \underline{\mathsf{MAT}}_{\underline{k}}$$

on objects by

$$Q(n,m) = mm$$

and on morphisms by

$$= \begin{vmatrix} a_{11}^{B} & \cdots & a_{1m}^{B} \\ \vdots & \vdots \\ a_{n1}^{B} & \cdots & a_{m}^{B} \end{vmatrix} \in Mor(np,mq)$$

and let e = 1 -- then with these choices, $\underline{MAT}_{\underline{k}}$ is a strict monoidal category. [Note: Write $\underline{FDVEC}_{\underline{k}}$ for the full subcategory of $\underline{VEC}_{\underline{k}}$ whose objects are finite dimensional -- then there is an equivalence $\underline{MAT}_{\underline{k}} \neq \underline{FDVEC}_{\underline{k}}$. Thus assign to each object n the vector space \underline{k}^n and to each morphism $n \neq m$ the linear map from \underline{k}^n to \underline{k}^m that sends $(x_1, \dots, x_n) \in \underline{k}^n$ to $(y_1, \dots, y_m) \in \underline{k}^m$, where y_i is the i^{th} entry of the 1-by-m matrix $[x_1, \dots, x_n]A$.]

1.8 <u>EXAMPLE</u> Given a C*-algebra A, let End A be the category whose objects are the unital *-homomorphisms $\Phi: A \rightarrow A$ and whose arrows $\Phi \rightarrow \Psi$ are the intertwiners, i.e.,

$$Mor(\Phi, \Psi) = \{ \mathbf{T} \in A : \mathbf{T}\Phi(\mathbf{A}) = \Psi(\mathbf{A})\mathbf{T} \forall \mathbf{A} \in \mathbf{A} \}.$$

Here, the composition of arrows, when defined, is given by the product in A and $1_A \in Mor(\Phi, \Phi)$ is 1_{Φ} . Define

Q:End A
$$\times$$
 End A \rightarrow End A

on objects by

$$\mathbf{Q}(\Phi, \Phi^{\mathsf{I}}) = \Phi \circ \Phi^{\mathsf{I}}$$

and on morphisms by

$$\begin{aligned} & \mathbf{T} & \mathbf{T'} \\ & \mathbf{\mathfrak{G}} & \stackrel{\mathbf{T}}{\rightarrow} \Psi, \ \Phi' & \stackrel{\mathbf{T}}{\longrightarrow} \Psi') \\ & = \mathbf{T} \Phi(\mathbf{T'}) (= \Psi(\mathbf{T'}) \mathbf{T}) \in \operatorname{Mor} (\Phi \circ \Phi', \ \Psi \circ \Psi') \end{aligned}$$

and let $e = id_A$ -- then with these choices, End A is a strict monoidal category. [Note: $\forall A \in A$, we have

 $T\Phi(T') (\Phi \circ \Phi') (A)$ $= T\Phi(T')\Phi(\Phi'(A))$ $= T\Phi(T'\Phi'(A))$ $= \Psi(T'\Phi'(A))T$ $= \Psi(\Psi'(A)T')T$ $= \Psi(\Psi'(A))\Psi(T')T$ $= (\Psi \circ \Psi')(A)T\Phi(T').]$

1.9 <u>EXAMPLE</u> Given a category <u>C</u>, let [<u>C</u>,<u>C</u>] be the metacategory whose objects are the functors $F:\underline{C} \rightarrow \underline{C}$ and whose morphisms are the natural transformations Ξ from F to G. Define

$$\&: [\underline{C}, \underline{C}] \rightarrow [\underline{C}, \underline{C}]$$

on objects by

 $\otimes(\mathbf{F},\mathbf{F}') = \mathbf{F} \circ \mathbf{F}'$

and on morphisms by

$$\Xi = \Xi'$$

$$Q(F \rightarrow G, F' \rightarrow G')$$

$$= \Xi Q \Xi',$$

where

$$(E \otimes E')_{X}$$

= E • FE' (= GE' • E),
G'X • FX

and let $e = id_{\underline{C}}$ (the identity functor) -- then with these choices, [C,C] is a strict monoidal category.

[Note: If

$$\begin{bmatrix} \exists \in Nat(F,G) \\ \\ \exists ' \in Nat(F',G'), \end{bmatrix}$$

then

.___. .

$$\begin{bmatrix} \forall X, Y \in Ob C \\ and \\ \forall X', Y' \in Ob C \end{bmatrix} \begin{bmatrix} \forall f \in Mor(X, Y) \\ \forall f' \in Mor(X', Y'), \end{bmatrix}$$

there are commutative diagrams

$$\begin{array}{c} FX \xrightarrow{\Xi_X} GX \\ Ff \downarrow & \qquad \downarrow Gf \\ FY \xrightarrow{\Xi_Y} GY \end{array}$$



In particular: The diagram



commutes. This said, the claim is that

 $\Xi \otimes \Xi' \in Nat(F \circ F', G \circ G'),$

i.e., that the diagram



commutes. In fact,

$$GG'f \circ (\Xi \boxtimes \Xi')_X$$
$$= GG'f \circ \Xi \circ F\Xi'_X$$
$$= GG'f \circ G\Xi'_X \circ \Xi$$
$$F'X$$

$$= G(G'f \circ E_X') \circ E_F'X$$

$$= G(E_Y' \circ F'f) \circ E_F'X$$

$$= GE_Y' \circ GF'f \circ E_F'X$$

$$= GE_Y' \circ E_F'f \circ FF'f$$

$$= E_GY' \circ FE_Y' \circ FF'f$$

$$= (E \otimes E')_Y \circ FF'f.]$$

1.10 LEMMA Suppose that C is monoidal and let e,e' be units -- then e and e' are isomorphic.

[There is an isomorphism $\phi: e \rightarrow e'$ for which the diagrams

commute, viz.

$$\phi = L \circ (R'_e)^{-1} \quad (e \to e \ a \ e' \to e').]$$

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§2. MONOIDAL FUNCTORS

Let C, C' be monoidal categories -- then a monoidal functor is a triple (F,ξ,Ξ) , where $F:\underline{C} \rightarrow \underline{C}'$ is a functor, $\xi:\underline{e}' \rightarrow Fe$ is an isomorphism, and the

$$\Xi_{X,Y}: FX \otimes' FY \rightarrow F(X \otimes Y)$$

are isomorphisms, natural in X,Y, subject to the following assumptions.

(MF1) The diagram

commutes.

commute.

<u>N.B.</u> A monoidal functor is said to be strict if ξ and Ξ are identities.

2.1 <u>EXAMPLE</u> Write <u>FDHILB</u> for the full subcategory of <u>HILB</u> whose objects are finite dimensional -- then the forgetful functor

is strict monoidal.

[Take for

$$E_{X,Y}: \mathbf{U}X \otimes \mathbf{U}Y \to \mathbf{U}(X \otimes Y)$$

the identity $id_{X \otimes Y}$ and let $\xi = id_{C^*}$]

[Note: A forgetful functor need not be monoidal, let alone strict monoidal. E.g.: Give <u>AB</u> its monoidal structure per the tensor product, give <u>SET</u> its monoidal structure per the cartesian product, and consider $U:\underline{AB} \rightarrow \underline{SET}$ — then the canonical maps

are not isomorphisms.]

Let

be monoidal functors -- then a monoidal natural transformation

$$(\mathbf{F}, \boldsymbol{\xi}, \boldsymbol{\Xi}) \rightarrow (\mathbf{G}, \boldsymbol{\theta}, \boldsymbol{\Theta})$$

is a natural transformation $\alpha: F \rightarrow G$ such that the diagrams

commute.

Write $[\underline{C},\underline{C}']^{\bigotimes}$ for the metacategory whose objects are the monoidal functors $\underline{C} \neq \underline{C}'$ and whose morphisms are the monoidal natural transformations.

N.B. A monoidal natural transformation is a monoidal natural isomorphism if α is a natural isomorphism.

2.2 <u>REMARK</u> Some authorities assume outright that Fe = e', the rationale being that this can always be achieved by replacing $F \in Ob [\underline{C},\underline{C'}]^{\underline{Q}}$ by an isomorphic $F' \in Ob [\underline{C},\underline{C'}]^{\underline{Q}}$ such that F'e = e' (on objects $X \neq e$, F'X = FX).

2.3 LEMMA Let

$$(\mathbf{F}, \Xi, \xi) \quad (\mathbf{F}; \underline{C} \to \underline{C}')$$
$$(\mathbf{F}', \Xi', \xi') \quad (\mathbf{F}': \underline{C}' \to \underline{C}'')$$

be monoidal functors -- then their composition F^{\ast} $\,\circ\,\,F$ is a monoidal functor.

[Consider the arrows e'' \longrightarrow F'e' \longrightarrow F'Fe and

$$\mathbf{F'FX} \otimes \mathbf{''} \quad \mathbf{F'FY} \xrightarrow{\Xi'} FX, FY \rightarrow F'(FX \otimes \mathbf{'FY}) \xrightarrow{F'\Xi} X, Y \rightarrow F'F(X \otimes Y).]$$

Write <u>MONCAT</u> for the metacategory whose objects are the monoidal categories and whose morphisms are the monoidal functors.

2.4 <u>RAPPEL</u> Let C, D be categories -- then a functor $F: C \rightarrow D$ is said to be an <u>equivalence</u> if there exists a functor $G: D \rightarrow C$ such that $G \circ F \approx id_{\underline{C}}$ and $F \circ G \approx id_{\underline{D}}$, the symbol \approx standing for natural isomorphism. 2.5 <u>LEMMA</u> A functor $F: \underline{C} \rightarrow \underline{D}$ is an equivalence iff it is full, faithful, and has a representative image (i.e., for any $Y \in Ob \underline{D}$, there exists an $X \in Ob \underline{C}$ such that FX is isomorphic to Y).

<u>N.B.</u> Categories <u>C</u>, <u>D</u> are said to be <u>equivalent</u> provided there is an equivalence $F:\underline{C} \rightarrow \underline{D}$. The object isomorphism types of equivalent categories are in a one-to-one correspondence.

2.6 RAPPEL Given categories
$$\begin{bmatrix} C \\ D \\ D \end{bmatrix}$$
, functors $\begin{bmatrix} F:C \rightarrow D \\ G:D \rightarrow C \end{bmatrix}$ are said to be an $G:D \rightarrow C$
adjoint pair if the functors $\begin{bmatrix} Mor \circ (F^{OP} \times id_D) \\ & & & \\ & & & \\ Mor \circ (id_{C^{OP}} \times G) \end{bmatrix}$ from $C^{OP} \times D$ to SET are naturally

isomorphic, i.e., if it is possible to assign to each ordered pair $\begin{bmatrix} x \in Ob \ C \\ y \in Ob \ D \end{bmatrix}$

a bijective map $E_{X,Y}$:Mor(FX,Y) \rightarrow Mor(X,GY) which is functorial in X and Y. When this is so, F is a <u>left adjoint</u> for G and G is a <u>right adjoint</u> for F. Any two left (right) adjoints for G (F) are naturally isomorphic. In order that (F,G) be an adjoint pair, it is necessary and sufficient that there exist natural trans-

The data (F,G,μ,ν) is referred to as an adjoint situation, the natural trans-

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formations

$$\mu: id_{\underline{C}} \rightarrow G \circ F$$

being the arrows of adjunction. An adjoint equiv-
 $v: F \circ G \rightarrow id_{\underline{D}}$

<u>alence</u> of categories is an adjoint situation (F,G,μ,ν) in which both μ and ν are natural isomorphisms.

2.7 <u>LETMA</u> A functor $F:\underline{C} \rightarrow \underline{D}$ is an equivalence iff F is part of an adjoint equivalence.

Let C, C' be monoidal categories — then C, C' are monoidally equivalent if there are monoidal functors

and monoidal natural isomorphisms

$$\begin{array}{c} F' \circ F \approx \operatorname{id}_{\underline{C}} \\ F \circ F' \approx \operatorname{id} \\ C' \end{array}$$

2.8 LEMMA Suppose that $F:\underline{C} \rightarrow \underline{C}'$ is a monoidal functor. Assume: F is an equivalence -- then F is a monoidal equivalence.

2.9 REMARK Embed F in an adjoint situation (F,F', μ , μ '), where

$$\begin{bmatrix} \mu: id_{\underline{C}} \to F' \circ F \\ \underline{C} & \mu': F \circ F' \to id \\ \underline{C'} & \underline{C'} & \underline{C'} \end{bmatrix}$$

are the arrows of adjunction (cf. 2.7) -- then one can equip F' with the structure
of a monoidal functor in such a way that the natural isomorphisms
$$\mu$$
, μ' are
monoidal natural isomorphisms. Thus first specify $\xi':e \rightarrow F'e'$ by taking it to

 $\begin{array}{l} \Xi^{\prime} & : \mathbf{F}^{\prime}\mathbf{X}^{\prime} \otimes \mathbf{F}^{\prime}\mathbf{Y}^{\prime} \rightarrow \mathbf{F}^{\prime}\left(\mathbf{X}^{\prime} \otimes^{\prime} \mathbf{Y}^{\prime}\right), \\ & \mathbf{X}^{\prime}, \mathbf{Y}^{\prime} \end{array}$

be the composition e $\xrightarrow{\mu_e}$ F'Fe $\xrightarrow{F'\xi^{-1}}$ F'e'. As for

1. $F'X' \otimes F'Y' \rightarrow F'F(F'X' \otimes F'Y');$

3. $FF^{1}X^{1} \xrightarrow{\mu} X^{1}$

2. F'F(F'X' @ F'Y') ------> F'(FF'X' @' FF'Y');

 $\begin{array}{cccc} \mu^{1} & \& & \mu^{1} & :FF^{1}X^{1} & \&^{1} & FF^{1}Y^{1} \rightarrow X^{1} & \&^{1} & Y^{1} \\ & X^{1} & & Y^{1} \end{array}$

 $\begin{array}{cccc} \mathbf{F}^{*}\left(\mu^{*} & \underline{\otimes} & \mu^{*}\right): \mathbf{F}^{*}\left(\mathbf{F}\mathbf{F}^{*}\mathbf{X}^{*} & \underline{\otimes}^{*} & \mathbf{F}\mathbf{F}^{*}\mathbf{Y}^{*}\right) & \rightarrow & \mathbf{F}^{*}\left(\mathbf{X}^{*} & \underline{\otimes}^{*} & \mathbf{Y}^{*}\right), \\ \mathbf{X}^{*} & \mathbf{Y}^{*} \end{array}$

If C is monoidal, then C^{OP} is monoidal when equipped with the same Ω and e,

 $\begin{bmatrix} - & R^{OP} = & R^{-1} \\ & L^{OP} = & L^{-1} \\ & & A^{OP} = & A^{-1}. \end{bmatrix}$

build it in three stages:

=>

=>

taking

6.

§3. STRICTIFICATION

A <u>strictification</u> of a monoidal category \underline{C} is a strict monoidal category which is monoidally equivalent to \underline{C} .

3.1 EXAMPLE $\underline{MAT}_{\underline{k}}$ is a strictification of $\underline{FDVEC}_{\underline{k}}$.

[The equivalence $\underline{MAT}_{\underline{k}} \rightarrow \underline{FDVEC}_{\underline{k}}$ constructed in 1.7 is a monoidal functor, hence is a monoidal equivalence (cf. 2.8).]

3.2 <u>THEOREM</u> Every monoidal category <u>C</u> is monoidally equivalent to a strict monoidal category \underline{C}_{str} .

The proof is constructive and best broken up into steps.

<u>Step 1</u>: Let S be the class of all finite sequences $S = (X_1, \dots, X_n)$ of objects of C, including the empty sequence Ø. Given nonempty

$$T = (X_{1}, \dots, X_{n})$$
$$T = (Y_{1}, \dots, Y_{m}),$$

let

$$S * T = (X_1, \dots, X_n, Y_1, \dots, Y_m)$$

and write

$$S \star \emptyset = S = \emptyset \star S.$$

Step 2: The claim is that \underline{S} is the object class of a strict monoidal

category $\underline{c}_{\underline{str}}$, i.e., $\underline{s} = 0b \underline{c}_{\underline{str}}$. In any event, the multiplication

 $*:S \times S \rightarrow S$

is associative, so we can take A to be the identity. Also, \emptyset serves as the unit and

$$= R_{S}: S \star \emptyset \to S$$
$$= L_{S}: \emptyset \star S \to S$$

are the identities.

<u>Step 3</u>: Given S, T, we need to specify Mor(S,T). For this purpose, define a map $\Gamma: \underline{S} \neq Ob \subseteq Dy \Gamma \emptyset = e$, $\Gamma((\underline{X})) = \underline{X}$, and $\Gamma(\underline{S} * (\underline{X})) = \Gamma \underline{S} \otimes \underline{X}$, thus

$$\Gamma(X_{1}, ..., X_{n}) = (... (X_{1} \otimes X_{2}) \otimes ...) \otimes X_{n},$$

where all opening parentheses are to the left of X_1 . Definition:

$$Mor(S,T) = Mor(\Gamma S, \Gamma T)$$
.

This prescription then gives rise to a category $C_{\underline{str}}$ with $Ob C_{\underline{str}} = S$.

<u>Step 4</u>: We shall now define a functor $*:\underline{C}_{\underline{str}} \times \underline{C}_{\underline{str}} \to \underline{C}_{\underline{str}}$ that serves to render $\underline{C}_{\underline{str}}$ strict monoidal, the issue being the meaning of

 $u \quad u^{*}$ $u \star u^{*} = \star (S \rightarrow T, S^{*} \rightarrow T^{*})$ $\in Mor(S \star S^{*}, T \star T^{*})$ $= Mor(\Gamma(S \star S^{*}), \Gamma(T \star T^{*})).$

Bearing in mind that

$$Mor(S,T) = Mor(TS,TT)$$

$$u \iff f$$

$$Mor(S',T') = Mor(TS',TT')$$

$$u' \iff f',$$

let u * u' be the composite

$$\begin{array}{c} \mathbf{f} \underline{\mathbf{0}} \mathbf{f}^{*} \\ \Gamma(\mathbf{S} \star \mathbf{S}^{*}) \to \Gamma \mathbf{S} \ \underline{\mathbf{0}} \ \Gamma \mathbf{S}^{*} & \longrightarrow \\ \Gamma \mathbf{T} \ \underline{\mathbf{0}} \ \Gamma \mathbf{T}^{*} \to \Gamma \langle \mathbf{T} \star \mathbf{T}^{*} \rangle , \end{array}$$

where the outer arrows are the obvious canonical morphisms in <u>C</u>. Accordingly, with this agreement, C_{str} is strict monoidal.

<u>Step 5:</u> It is clear from its very construction that $\Gamma:\underline{C}_{\underline{str}} \rightarrow \underline{C}$ is a functor which, moreover, is full, faithful, and is isomorphism dense. But $\Gamma \emptyset = e$ and there are isomorphisms

$$\mathbb{E}_{\mathbf{S},\mathbf{T}}:\Gamma\mathbf{S} \ \boldsymbol{\otimes} \ \Gamma\mathbf{T} \rightarrow \Gamma(\mathbf{S} \ast \mathbf{T}),$$

natural in S, T and satisfying $M\Gamma_1$, $M\Gamma_2$ of §2. Therefore Γ is monoidal. To finish, it remains only to quote 2.8.

[Note: It is not necessary to quote 2.8: Simply observe that there is an inclusion functor $\gamma: \underline{C} \rightarrow \underline{C}_{str}$ and

4.

Detail: From

$$Mor(\gamma \Gamma S, S) = Mor(\Gamma S, \Gamma S)$$

let

$$\alpha_{\rm S} \iff {\rm id}_{\rm \Gamma S}$$
,

thus $\alpha_{S}:\gamma\Gamma S \rightarrow S$ and $\alpha:\gamma \circ \Gamma \rightarrow id_{\underline{C}}$ is a monoidal natural isomorphism.]

3.3 <u>REMARK</u> Let C, C' be monoidal categories -- then each monoidal functor $F:C \rightarrow C'$ induces a strict monoidal functor $F_{\underline{str}}:C_{\underline{str}} \rightarrow C'_{\underline{str}}$ and there is a commutative diagram



Here, on an object S,

$$\mathbf{F}_{\underline{str}} \mathbf{S} = (\mathbf{FX}_1, \dots, \mathbf{FX}_n),$$

while on a morphism u:S \rightarrow T,

$$(FX_1, \dots, FX_n) \xrightarrow{F \underline{str}^u} (FY_1, \dots, FY_m)$$

is that element of Mor(FFS,FFT) defined by requiring commutativity of the square



where $f \in Mor(\Gamma S, \Gamma T)$ corresponds to u.

[Note: Composition of monoidal functors is preserved by this construction.]

There are five ingredients figuring in the definition of a monoidal category: Ω , e, R, L, A. Keeping track of R, L, A in calculations can be annoying and one way out is to pass from <u>C</u> to <u>C</u><u>str</u>. But this too has its downside since <u>C</u><u>str</u> is a more complicated entity than <u>C</u>. So, in what follows, we shall stick with <u>C</u> and determine to what extent R, L, A can be eliminated from consideration (i.e., are identities).

Suppose that

are monoidal structures on \underline{C} -- then these structures are deemed <u>isomorphic</u> if \exists a monoidal equivalence of the form $(id_{\underline{C}}, \xi, \Xi)$ between them.

N.B. Therefore $\xi:e' \rightarrow e$ is an isomorphism and the

$$\Xi_{X,Y}: X \otimes Y \to X \otimes Y$$

are isomorphisms, subject to the coherence conditions of §2.

3.4 <u>REMARK</u> The philosophy is that replacing a given monoidal structure on \underline{C} by another isomorphic to it is of no consequence for the underlying mathematics.

3.5 LEMMA Let (Q, e, R, L, A) be a monoidal structure on C. Suppose given a map Q':Ob C × Ob C → Ob C, an object e' \in Ob C, an isomorphism $\xi':e \rightarrow e'$, and isomorphisms

$$E_{X,Y}^{1}:X \otimes Y \to X \otimes^{!} Y.$$

Then there is a unique monoidal structure (Q', e', R', L', A') on C such that

$$(\operatorname{id}_{\underline{C}}, \xi', \Xi'): (\underline{C}, \mathbf{Q}', e', R', L', A') \rightarrow (\underline{C}, \mathbf{Q}, e, R, L, A)$$

is an isomorphism.

<u>PROOF</u> Extend Q' to a functor $Q': C \times C \rightarrow C$ by the prescription

so $\Omega \approx \Omega'$ (via $E' \in Nat(\Omega, \Omega')$). This done, define R', L', A' by the diagrams



3.6 <u>THEOREM</u> Let (\mathfrak{G} , e, R, L, A) be a monoidal structure on <u>C</u>. Suppose that e' is an object isomorphic to e, say $\xi:e' \rightarrow e$ — then there is an isomorphic monoidal structure (\mathfrak{G}' , e', R', L', A') on <u>C</u> in which R', L' are identities.

PROOF Bearing in mind 3.5, put

$$X \otimes Y = X \otimes Y \text{ if } X \neq e^1 \neq Y$$

and

$$X Q' Y = \begin{bmatrix} - & Y & \text{if } X = e' \\ & \\ & X & \text{if } Y = e' \end{bmatrix}$$

Define

 $\Xi'_{X,Y}$: X Q Y \rightarrow X Q' Y

by stipulating that $E'_{X,Y}$ is to be the identity if $X \neq e^* \neq Y$, otherwise let

$$\begin{bmatrix} \exists ' &= R_X \circ (id_X \otimes \xi) \\ X, e' & \end{bmatrix}$$
$$\begin{bmatrix} \exists ' &= L_Y \circ (\xi \otimes id_Y) \\ e', Y & \end{bmatrix}$$

To establish consistency, i.e., that

 $\begin{array}{cccc} R & \circ (id & \& \xi) = L & \circ (\xi \& id), \\ e' & e' & e' & e' \end{array}$

set $\xi' = \xi^{-1}$ -- then

is an isomorphism and due to the naturality of R, L, the diagrams

A,1



commute. Therefore

$$R \circ (id \otimes \xi) \circ (\xi' \otimes \xi')$$

$$= R \circ (\xi' \otimes id_{e})$$

$$= \xi' \circ R_{e} = \xi' \circ L_{e} (R_{e} = L_{e})$$

$$= L \circ (id_{e} \otimes \xi')$$

$$= L \circ (\xi \otimes id_{e}) \circ (\xi' \otimes \xi)$$

$$e' e'$$

from which the contention. Finally, by construction (cf. 3.5), R', L' are identities. E.g.:

$$\begin{array}{ccc} R'_X \circ \Xi' & \circ \text{ id}_X \otimes \xi' = R_X \\ & X, e' \end{array}$$

or still,

$$R_X' \circ R_X \circ (id_X \otimes \xi) \circ id_X \otimes \xi' = R_X$$

or still,

$$R'_X \circ R_X = R_X = R'_X = id_X.$$

[Note: If A is the identity and e' is not in the image of Q, then A' is

also the identity. Proof:

$$\begin{bmatrix} e' \in \{X,Y,Z\} \Rightarrow A_{XYZ}^{t} = id \\ e' \notin \{X,Y,Z\} \& e' \notin Im @ \Rightarrow A_{XYZ}^{t} = A_{XYZ}.\end{bmatrix}$$

3.7 <u>REMARK</u> Take e' = e — then the preceding result implies that by passing to an isomorphic monoidal structure, it is always possible to arrange that $\forall X \in Ob C$,

The situation for the associativity constraint is more complicated and it will be necessary to impose some conditions on C.

Definition: A construct is a pair (C,U), where

$$U:C \rightarrow SET$$

is a faithful functor.

3.8 EXAMPLE Define a functor $Q:\underline{SET}^{OP} \rightarrow \underline{SET}$ as follows: On objects, $QX = 2^X$ and on morphisms, $Q(A \rightarrow B):QA \rightarrow QB$ sends $X \subset A$ to the inverse image $f^{-1}(X) \subset B$. In this connection, recall that

$$\begin{array}{c} \mathbf{f} \\ \mathbf{A} \neq \mathbf{B} \in \mathsf{Mor} \ \underline{\mathbf{SET}}^{\mathsf{OP}} \end{array}$$

means that

$$\begin{array}{c} f \\ B \rightarrow A \in Mor \ \underline{SET}. \end{array}$$

Therefore $(\underline{SET}^{OP}, Q)$ is a construct.

Let (\underline{C}, U) be a construct -- then (\underline{C}, U) is <u>ammestic</u> if a <u>C</u>-isomorphism f is a <u>C</u>-identity whenever Uf is a <u>SET</u>-identity, i.e., if $X, Y \in Ob \underline{C}$, if $f:X \rightarrow Y$ is an isomorphism, if Uf = id, then X = Y and f = id.

Let (C,U) be a construct -- then (C,U) is <u>transportable</u> if \forall C-object X and every bijection UX \Rightarrow S, \exists a C-object Y with UY = S and an isomorphism $\Phi: X \Rightarrow Y$ such that $U\Phi = \phi$.

3.9 LEMMA If (C,U) is amnestic and transportable, then the pair (Y, ϕ) is unique.

PROOF Say we have

$$\mathbf{Y}_1 \xrightarrow{\Phi_1^{-1}} \mathbf{X} \xrightarrow{\Phi_2} \mathbf{Y}_2.$$

Then $\Phi_2 \circ \Phi_1^{-1}$ is an isomorphism and

 $U(\Phi_2 \circ \Phi_1^{-1}) = U\Phi_2 \circ U\Phi_1^{-1} = \phi \circ \phi^{-1} = id.$

Therefore by amnesticity, $Y_1 = Y_2$ and $\Phi_2 \circ \Phi_1^{-1} = id => \Phi_2 = \Phi_1$.

3.10 <u>EXAMPLE</u> The construct <u>FDVEC</u> is annestic and transportable but the full subcategory of <u>FDVEC</u> whose objects are the \underline{k}^n , while annestic, is not transportable.

3.11 <u>LEMMA</u> If $\zeta:\underline{SET} \rightarrow \underline{SET}$ is an isomorphism and if (\underline{C}, U) is amnestic and transportable, then $(\underline{C}, \zeta \circ U)$ is amnestic and transportable.

3.12 <u>THEOREM</u> Suppose that (C,U) is amnestic and transportable. Let (Ω, e, R, L, A) be a monoidal structure on <u>C</u> -- then there is an isomorphic strict monoidal structure (Ω', e, R', L', A') on <u>C</u>.

The proof is lengthy, the point of departure being 3.2:

$$\Gamma:\underline{C}_{\underline{str}} \to \underline{C}$$
$$\gamma:\underline{C} \to \underline{C}_{\underline{str}},$$

where

$$\begin{bmatrix} \Gamma \circ \gamma = id_{\underline{C}} \\ \gamma \circ \Gamma \approx id_{\underline{C}} \\ \underline{\zeta} \\$$

<u>Step 1</u>: Given $s \in Ob C_{\underline{str}'}$ consider

$$\{S\} \times UTS \in Ob SET.$$

Then the projection

$$\{s\} \times UTS \xrightarrow{\pi} S UTS$$

is bijective, so there exists a unique $[S] \in Ob \subseteq$ with $U[S] = \{S\} \times U\Gamma S$ and a unique isomorphism $\Pi_S: [S] \rightarrow \Gamma S$ such that $U\Pi_S = \Pi_S$.

<u>Step 2</u>: There is a functor $\overline{\Gamma:C}_{\underline{str}} \rightarrow \underline{C}$ which on objects is the prescription

$$\overline{\Gamma}S = [S]$$

and on morphisms is dictated by requiring that $\Pi \in \operatorname{Nat}(\overline{\Gamma}, \Gamma)$:



<u>Step 3</u>: $\overline{\Gamma}:\underline{C}_{\underline{str}} \to \underline{C}$ is an equivalence of categories $(\Pi:\overline{\Gamma} \to \Gamma$ being a natural isomorphism). In addition, $\overline{\Gamma}$ is injective on objects.

<u>Step 4</u>: Define a functor $\overline{\gamma}: \underline{C} \to \underline{C}_{\underline{str}}$ on objects by taking $\overline{\gamma}X = \gamma X$ if X is not in the image of $\overline{\Gamma}$ and letting $\overline{\gamma}[S] = S$ otherwise. Next, define

$$v_{\mathbf{X}} : \overline{\Gamma} \overline{\gamma} \mathbf{X} \to \mathbf{X}$$

by

if X is not in the image of $\overline{\Gamma}$ and let $v_X = id_X$ if X = [S] for some S. Since $\overline{\Gamma}$ is fully faithful, we can then define $\overline{\gamma}$ on morphisms by requiring that $v:\overline{\Gamma} \circ \overline{\gamma} \to id_{\underline{C}}$ be a natural isomorphism.

Step 5: The arrow

$$\mu = \text{id:id}_{\underline{C}} \rightarrow \overline{\gamma} \circ \overline{\Gamma}$$

is a natural isomorphism.

<u>Step 6</u>: The data $(\overline{\Gamma}, \overline{\gamma}, \mu, \nu)$ is an adjoint situation:

$$\begin{vmatrix} \overline{\Gamma} & (\nu \overline{\Gamma}) & \circ (\overline{\Gamma} \mu) = \text{id} \\ & \overline{\Gamma} & \\ (\overline{\gamma} \nu) & \circ (\mu \overline{\gamma}) = \text{id} \\ & \overline{\gamma} & \end{vmatrix}$$

Explicated:

$$\begin{bmatrix} v & \circ \bar{\Gamma}\mu_{\mathbf{S}} = \mathbf{id} \\ \bar{\Gamma}\mathbf{S} & \bar{\Gamma}\mathbf{S} \end{bmatrix}$$
$$\begin{bmatrix} \bar{\gamma}\nu_{\mathbf{X}} & \circ \mu_{\mathbf{Y}} = \mathbf{id} \\ \bar{\gamma}X & \sigma_{\mathbf{Y}}X & \sigma_{\mathbf{Y}}X \end{bmatrix}$$

Claim:

$$\begin{bmatrix} v &= id \\ \overline{\Gamma}S & \overline{\Gamma}S \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & &$$

But

$$\overline{\Gamma}S = [S] \Rightarrow v = id (\Xi \overline{\Gamma}\mu_S).$$

As for the relation

$$\overline{\gamma}v_{\mathbf{X}} = \operatorname{id}_{\overline{\gamma}\mathbf{X}} (\equiv \mu_{\overline{\gamma}\mathbf{X}}),$$

since $\overline{\Gamma}$ is faithful, it suffices to show that

$$\overline{\Gamma}_{Y}v_{X} = \operatorname{id}_{\overline{\Gamma}_{Y}X}$$

for all $X \in Ob \ \underline{C}$. But from the definitions, $\forall \ f \in Mor(\overline{\Gamma\gamma}X,X)$, there is a commutative diagram

$$\begin{array}{c} \overline{\Gamma}\overline{\gamma}\overline{\Gamma}\overline{\gamma}x \xrightarrow{\overline{\Gamma}\overline{\gamma}f} & \overline{\Gamma}\overline{\gamma}x \\ \nu & & \downarrow & \downarrow & \nu_{X} \\ \overline{\Gamma}\overline{\gamma}x & \downarrow & & \downarrow & \nu_{X} \\ & & \overline{\Gamma}\overline{\gamma}x \xrightarrow{f} & x \end{array}$$

Now take $f = v_X$ to get

$$v_{\mathbf{X}} \circ \overline{\Gamma} \overline{\gamma} v_{\mathbf{X}} = v_{\mathbf{X}} \circ v_{\overline{\Gamma} \overline{\gamma} \mathbf{X}}$$

or still,

$$\overline{\Gamma}\overline{\gamma}\nu_{X} = \nu_{\overline{\Gamma}\overline{\gamma}X}$$

or still,

$$\overline{\Gamma}\overline{\gamma}v_{X} = \operatorname{id}_{\overline{\Gamma}\overline{\gamma}X},$$

as desired.

<u>Step 7</u>: The adjoint situation $(\overline{\Gamma}, \overline{\gamma}, \mu, \nu)$ is an adjoint equivalence of categories (μ and ν are natural isomorphisms).

Step 8: Put

$$\mathbf{X} \ \mathbf{\otimes}' \ \mathbf{Y} = \overline{\Gamma} (\overline{\gamma} \mathbf{X} \ \star \ \overline{\gamma} \mathbf{Y})$$

and let $e' = \overline{\Gamma} \emptyset$ -- then

$$\overline{\gamma} (\mathbf{X} \otimes' \mathbf{Y}) = \overline{\gamma} \overline{\Gamma} (\overline{\gamma} \mathbf{X} \star \overline{\gamma} \mathbf{Y})$$
$$= \overline{\gamma} \mathbf{X} \star \overline{\gamma} \mathbf{Y}$$

anđ

.....

$$\overline{\gamma}e' = \overline{\gamma}\overline{\Gamma}\emptyset = \emptyset.$$

Step 9: We have

$$\begin{array}{rcl} X \ {} {\boldsymbol{ \Theta}}^{*} & (Y \ {\boldsymbol{ \Theta}}^{*} \ Z) \end{array} &= \ \overline{\Gamma} \left(\overline{\gamma} X \ \star \ \overline{\gamma} \left(Y \ {\boldsymbol{ \Theta}}^{*} \ Z \right) \right) \\ \\ &= \ \overline{\Gamma} \left(\overline{\gamma} X \ \star \ \overline{\gamma} Y \ \star \ \overline{\gamma} Z \right) \\ \\ &= \ \overline{\Gamma} \left(\overline{\gamma} \left(X \ {\boldsymbol{ \Theta}}^{*} \ Y \right) \ \star \ \overline{\gamma} Z \right) \\ \\ &= \ (X \ {\boldsymbol{ \Theta}}^{*} \ Y) \ {\boldsymbol{ \Theta}}^{*} \ Z, \end{array}$$
so A' = id will work.

Step 10: Let

$$\mathbf{R}_{\mathbf{X}}^{\bullet} = \mathbf{v}_{\mathbf{X}}$$

: $\overline{\Gamma}\overline{\gamma}\mathbf{X} \neq \mathbf{X}$
$$\mathbf{L}_{\mathbf{X}}^{\bullet} = \mathbf{v}_{\mathbf{X}}$$

Then this makes sense:

$$\begin{array}{c} \mathbf{X} \ \mathbf{\hat{\omega}}' \ \mathbf{e}' \ = \ \overline{\Gamma} \left(\overline{\gamma} \mathbf{X} \ \star \ \overline{\gamma} \overline{\Gamma} \mathbf{\hat{\beta}} \right) \ = \ \overline{\Gamma} \left(\overline{\gamma} \mathbf{X} \ \star \ \mathbf{\hat{\beta}} \right) \ = \ \overline{\Gamma} \overline{\gamma} \mathbf{X} \\ \\ \mathbf{e}' \ \mathbf{\hat{\omega}}' \ \mathbf{X}' \ = \ \overline{\Gamma} \left(\overline{\gamma} \overline{\Gamma} \mathbf{\hat{\beta}} \ \star \ \overline{\gamma} \mathbf{X} \right) \ = \ \overline{\Gamma} \left(\mathbf{\hat{\beta}} \ \star \ \overline{\gamma} \mathbf{X} \right) \ = \ \overline{\Gamma} \overline{\gamma} \mathbf{X}. \end{array}$$

Furthermore, the diagram

$$\begin{array}{c|c} X & @' & e' & @' & Y \\ \hline X & @' & e' & @' & Y \\ \hline & & & & & & & \\ id & @' & L' & & & & & & \\ X & @' & Y & & & & & \\ \hline & & & & & & & \\ X & @' & Y & & & & & \\ \hline & & & & & & & \\ X & @' & Y & & & & \\ \hline \end{array}$$

commutes. To see this, note first that

$$X \ \mathfrak{B}^{\mathbf{i}} \ \mathbf{e}^{\mathbf{i}} \ \mathfrak{B}^{\mathbf{i}} \ \mathbf{Y} = \overline{\Gamma} (\overline{\gamma} X \ \star \ \overline{\gamma} \mathbf{e}^{\mathbf{i}} \ \star \ \overline{\gamma} \mathbf{Y})$$
$$= \overline{\Gamma} (\overline{\gamma} X \ \star \ \mathfrak{I} \ \star \ \overline{\gamma} \mathbf{Y})$$
$$= \overline{\Gamma} (\overline{\gamma} X \ \star \ \overline{\gamma} \mathbf{Y})$$
$$= \overline{\Gamma} (\overline{\gamma} X \ \star \ \overline{\gamma} \mathbf{Y})$$
$$= X \ \mathfrak{B}^{\mathbf{i}} \ \mathbf{Y}.$$

And the arrows

$$= \begin{array}{c} R_1' \otimes ' \operatorname{id}_Y : X \otimes ' e' \otimes ' Y \to X \otimes ' Y \\ \operatorname{id}_X \otimes ' L_Y' : X \otimes ' e' \otimes ' Y \to X \otimes ' Y \\ - \end{array}$$

are identities. E.g.:

$$R_{X}' \ \mathfrak{Q}' \ id_{Y} = \overline{\Gamma}(\overline{\gamma}v_{X} * \overline{\gamma} \ id_{Y})$$

$$= \overline{\Gamma}(id * id_{1})$$

$$= \overline{\Gamma}(id * id_{1})$$

$$= \overline{\Gamma}(id)$$

$$\overline{\gamma}X * \overline{\gamma}Y$$

$$= id$$

$$\overline{\Gamma}(\overline{\gamma}X * \overline{\gamma}Y)$$

$$= id$$

$$X \ \mathfrak{Q}' Y$$

Step 11: It is clear that

$$\overline{\gamma}: (\underline{C}, \underline{\omega}^{\dagger}, e^{\dagger}, R^{\dagger}, L^{\dagger}, A^{\dagger}) \rightarrow (\underline{C}_{\underline{str}}, \star, \emptyset, R, L, A)$$

is a monoidal equivalence (cf. 2.8), thus the same is true of

$$\mathbb{F}\widetilde{\gamma}: (\underline{C},\underline{\omega}',\underline{e}',\underline{R}',\underline{L}',\underline{A}') \rightarrow (\underline{C},\underline{\omega},\underline{e},\underline{R},\underline{L},\underline{A}) \quad (\texttt{cf. 2.3}).$$

But there is a monoidal natural isomorphism $\Gamma \overline{\gamma} \approx id_{C}: \forall X \in Ob \underline{C}$,

$$\Gamma \overline{\gamma} x \xrightarrow{\overline{\gamma} x} \overline{\overline{\gamma} x} \rightarrow \overline{\Gamma} \overline{\gamma} x \xrightarrow{\nu} x \xrightarrow{\nu} x.$$

Therefore the monoidal structure (@',e',R',L',A') is isomorphic to (@,e,R,L,A).

<u>Step 12</u>: To complete the proof, it is necessary to fine tune (@',e',R',L',A') by an application of 3.6:

$$(\Omega', e', R', L', A') \rightarrow (\Omega'', e'', R'', L'', A''),$$

choosing e'' = e (cf. 1.10). So, R'', L'' are identities. However, by construction, A' is the identity, thus if e is not in the image of Ω ', then A'' is also the identity. To ensure that e is not in the image of Ω ', it is enough that e is not in the image of $\overline{\Gamma}$. Suppose it were -- then

$$Ue = \{S\} \times UIS \quad (\exists S \in Ob \subseteq_{str}).$$

Now use 3.11 and replace U by ζU , where ζ has the property that ζUe is not a cartesian product of two sets.

3.13 EXAMPLE Consider the construct $\underline{FDVEC}_{\underline{k}}$ — then the failure of the tensor product to be associative "on the nose" is an artifact of its definition by a universal property which determines it only up to isomorphism. While the usual procedures do not lead to an associative tensor product, the lesson to be drawn from 3.12 is that it is possible to find a tensor product on $\underline{FDVEC}_{\underline{k}}$ such that

$$\begin{array}{c} x & \underline{\mathbf{a}}_{\underline{k}} & \underline{\mathbf{k}} = x \\ \underline{\mathbf{k}} & \underline{\mathbf{a}}_{\underline{k}} & x = x \end{array}$$

and

$$(X \otimes_{\underline{k}} Y) \otimes_{\underline{k}} Z = X \otimes_{\underline{k}} (Y \otimes_{\underline{k}} Z) = X \otimes_{\underline{k}} Y \otimes_{\underline{k}} Z.$$

§4. SYMMETRY

A symmetry for a monoidal category C is a natural isomorphism T, where

$$T_{X,Y}: X \otimes Y \rightarrow Y \otimes X,$$

such that

$$T_{Y,X} \circ T_{X,Y} : X \otimes Y \to X \otimes Y$$

is the identity, $R_{X} = L_{X} \circ T_{X,e}$, and the diagram

commutes. A <u>symmetric monoidal category</u> is a monoidal category <u>C</u> endowed with a symmetry τ . A monoidal category can have more than one symmetry (or none at all).

[Note: The "coherency" principle then asserts that "all" diagrams built up from instances of R, L, A, τ (or their inverses), and id by repeated application of Ω necessarily commute.]

N.B. Let

$$\mathbf{f}: \underline{\mathbf{C}} \times \underline{\mathbf{C}} \to \underline{\mathbf{C}} \times \underline{\mathbf{C}}$$

be the interchange -- then f is an isomorphism and $\tau: \Omega \to \Omega \circ f$ is a natural isomorphism.

E.g.: $\underline{\text{VEC}}_{k}$ and $\underline{\text{HILB}}$ are symmetric monoidal.

4.1 EXAMPLE Let C*ALG be the category whose objects are the C*-algebras

and whose morphisms are the *-homomorphisms -- then under the minimal tensor product or the maximal tensor product, <u>C*ALG</u> is a symmetric monoidal category.

4.2 EXAMPLE Let CHX be the category of chain complexes of abelian groups and chain maps — then CHX is monoidal: Take X \otimes Y to be the tensor product and let $e = \{e_n\}$ be the chain complex defined by $e_0 = \underline{Z}$ and $e_n = 0$ (n $\neq 0$). Further-

more, if $\begin{vmatrix} & X = \{X_p\} \\ & \text{ and if } \\ & Y = \{Y_q\} \end{vmatrix} \begin{vmatrix} & X \in X_p \\ & \text{ , then the assignment } \\ & y \in Y_q \end{vmatrix}$ $\begin{vmatrix} & X \otimes Y \neq Y \otimes X \\ & \text{ is a symmetry for } \underline{CHX}. \\ & X \otimes Y \neq (-1)^{pq} (Y \otimes X) \end{vmatrix}$

4.3 <u>REMARK</u> In the strict situation, matters reduce to the relations $T_{e,X} = T_{X,e} = id_X$ and

$$\mathsf{T}_{X \otimes Y, Z} = (\mathsf{T}_{X, Z} \otimes \mathrm{id}_{Y}) \circ (\mathrm{id}_{X} \otimes \mathsf{T}_{Y, Z}).$$

[Note: Therefore

$$^{\mathsf{T}}_{\mathsf{X} \otimes \mathsf{Y}, \mathsf{Z}} \circ ^{\mathsf{T}}_{\mathsf{Y} \otimes \mathsf{Z}, \mathsf{X}} \circ ^{\mathsf{T}}_{\mathsf{Z} \otimes \mathsf{X}, \mathsf{Y}} = \mathrm{id.}]$$

4.4 <u>EXAMPLE</u> Let \$ be the permutation category introduced in 1.6 -- then \$ is symmetric monoidal. To establish this, one must exhibit isomorphisms

$$\tau_{n,m} \in Mor(n \otimes m, m \otimes n)$$

= \$_{n+m}

fulfilling the various conditions. Definition:

$$\tau_{n,m} = \begin{vmatrix} -1 & 2 & \dots & n & n+1 & n+2 & \dots & n+m \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & &$$

with the understanding that $\tau_{n,0} = id_n = \tau_{0,n'}$ thus

$$T_{m,n} \circ T_{n,m} = id_n \otimes m^*$$

As for the remaining details, it is simplest to work with permutation matrices, so take n > 0, m > 0, and note that

$$\tau_{n,m} = \begin{vmatrix} 0 & I_m \\ & m \\ & & \\ I_n & 0 \end{vmatrix}.$$

Then

....

$$(\mathsf{T}_{n,p} \ {\mbox{ad}} \ {\mbox{id}}_{\mathfrak{m}}) \circ (\mathrm{id}_{n} \ {\mbox{ad}} \ {\mbox{T}}_{\mathfrak{m},p})$$

$$= \begin{vmatrix} 0 & \mathbf{I}_{p} & 0 \\ \mathbf{I}_{n} & 0 & 0 \\ 0 & 0 & \mathbf{I}_{m} \end{vmatrix} \begin{vmatrix} \mathbf{I}_{n} & 0 & 0 \\ 0 & 0 & \mathbf{I}_{p} \\ 0 & \mathbf{I}_{m} & 0 \end{vmatrix}$$
$$= \begin{vmatrix} 0 & 0 & \mathbf{I}_{p} \\ \mathbf{I}_{n} & 0 & 0 \\ 0 & \mathbf{I}_{m} & 0 \end{vmatrix} = \mathbf{T}_{n \otimes m, p},$$

[Note:

$$\begin{bmatrix} 0 & \mathbf{I}_{\mathbf{m}} \\ \mathbf{I}_{\mathbf{n}} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \sigma & 0 \\ 0 & \tau \end{bmatrix} = \begin{bmatrix} 0 & \tau \\ \sigma & 0 \end{bmatrix}$$
$$= \begin{bmatrix} \tau & 0 \\ 0 & \sigma \end{bmatrix} \begin{bmatrix} 0 & \mathbf{I}_{\mathbf{m}} \\ \mathbf{I}_{\mathbf{n}} & \mathbf{0} \end{bmatrix}.$$

Therefore naturality is manifest, i.e.,

$$\tau_{n,m} \circ (\sigma \otimes \tau) = (\tau \otimes \sigma) \circ \tau_{n,m}$$
.]

Let C, C' be symmetric monoidal categories -- then a symmetric monoidal functor is a monoidal functor (F, ξ ,E) such that the diagram



commutes.

<u>N.B.</u> The monoidal natural transformations between symmetric monoidal functors are, by definition, "symmetric monoidal" (i.e., no further conditions are imposed

that reflect the presence of a symmetry).

[Note: Therefore the subcategory $[\underline{C},\underline{C}']^{\underline{\otimes},\top}$ of $[\underline{C},\underline{C}']^{\underline{\otimes}}$ whose objects are the symmetric monoidal natural transformations is, by definition, a full subcategory.]

4.5 EXAMPLE Recall that $\$_n$ has the following presentation: It is generated by $\sigma_1, \dots, \sigma_{n-1}$ subject to the relations

$$\sigma_{\mathbf{i}}^2 = \mathbf{1}, \ \sigma_{\mathbf{i}}\sigma_{\mathbf{i}+1}\sigma_{\mathbf{i}} = \sigma_{\mathbf{i}+1}\sigma_{\mathbf{i}}\sigma_{\mathbf{i}+1}, \ \sigma_{\mathbf{i}}\sigma_{\mathbf{j}} = \sigma_{\mathbf{j}}\sigma_{\mathbf{i}} \ (|\mathbf{i}-\mathbf{j}| > 1).$$

Suppose now that <u>C</u> is symmetric strict monoidal and fix $X \in Ob$ <u>C</u>. Define automorphisms $\mathbb{I}^1, \ldots, \mathbb{I}^{n-1}$ of $x^{\otimes n}$ by

$$\Pi^{\perp} = \operatorname{id}_{X^{\mathfrak{Q}}(i-1)} \overset{\mathfrak{Q}}{\xrightarrow{}} T_{X,X} \overset{\mathfrak{Q}}{\xrightarrow{}} \operatorname{id}_{X^{\mathfrak{Q}}(n-i-1)}$$

Then there exists a unique homomorphism

$$\Pi_n^X:\mathfrak{F}_n \to \operatorname{Aut} X^{\otimes n}$$

of groups such that

$$\Pi_n^X(\sigma_i) = \Pi^i$$
 (i = 1,...,n-1).

Combining the Π_n^X then leads to a symmetric monoidal functor $F: \mathfrak{F} \to \underline{C}$ such that $Fn = x^{\otimes n}$.

and recall that F has a representative image (cf. 2.5).]

.

4.7 <u>EXAMPLE</u> If <u>C</u> is symmetric monoidal, then <u>C</u><u>str</u> is symmetric monoidal and $\gamma:\underline{C} \rightarrow \underline{C}_{\underline{str}}$ is a symmetric monoidal equivalence.

4.8 <u>LEMMA</u> Let C, C' be symmetric monoidal and let (F,F',μ,μ') be an adjoint equivalence. Assume: F is symmetric monoidal — then F' is symmetric monoidal (cf. 2.9).

§5. DUALITY

Let \underline{C} be a monoidal category — then each $X \in Ob \ \underline{C}$ defines functors

$$- \underline{\otimes} X: \underline{C} \to \underline{C}$$
$$\underline{X } \underline{\otimes} - : \underline{C} \to \underline{C}.$$

Definition: \underline{C} is

if $\forall X \in Ob \underline{C}$,

$$\otimes$$
 X admits a right adjoint, denoted lhom(X,--)
 $-$ X \otimes -- admits a right adjoint, denoted rhom (X,--).

[Note: <u>C</u> is <u>closed</u> if it is both left closed and right closed.] So:

$$\underline{C} \text{ left closed} \Rightarrow \text{Mor}(Y \otimes X, Z) \approx \text{Mor}(Y, \text{lhom}(X, Z))$$
$$\underline{C} \text{ right closed} \Rightarrow \text{Mor}(X \otimes Y, Z) \approx \text{Mor}(Y, \text{rhom}(X, Z))$$

for all $Y,Z \in Ob \underline{C}$.

 $\underline{N.B.}$ The functor

is called the
$$\begin{bmatrix} - & \text{lhom}(X, --) \\ & \text{rhom}(X, --) \end{bmatrix}$$

5.1 <u>REMARK</u> If <u>C</u> is symmetric monoidal, then left and right internal homs are naturally isomorphic and if <u>C</u> is left or right closed, then <u>C</u> is closed.

5.2 <u>EXAMPLE</u> Given a commutative ring <u>k</u>, let $\underline{MOD}_{\underline{k}}$ be the category whose objects are the left <u>k</u>-modules and whose morphisms are the <u>k</u>-linear maps — then $\underline{MOD}_{\underline{k}}$ is symmetric monoidal. Moreover, $\underline{MOD}_{\underline{k}}$ is closed and

Thom $(X,Z) \approx \operatorname{Hom}_{\underline{k}} (X,Z)$ rhom $(X,Z) \approx \operatorname{Hom}_{\underline{k}} (X,Z)$.

5.3 <u>LEMMA</u> Suppose that <u>C</u> is left closed — then $\forall X \in Ob \underline{C}$, the functor — $\underline{O} X$ preserves colimits (being a left adjoint) and the functor lhom(X, -)preserves limits (being a right adjoint).

5.4 <u>LEMMA</u> Suppose that <u>C</u> is left closed -- then $\forall Z \in Ob \underline{C}$, the cofunctor lhom(-,Z) converts colimits to limits.

<u>PROOF</u> Let <u>I</u> be a small category, $\Delta: \underline{I} \rightarrow \underline{C}$ a diagram for which colim<u>I</u> $\Delta_{\underline{I}}$ exists -- then $\forall Y \in Ob \underline{C}$,

Mor (Y, lhom (colim_I Δ_i , Z)) $\approx Mor(Y \otimes colim_I \Delta_i, Z)$ $\approx Mor(colim_I (Y \otimes \Delta_i), Z)$

$$\approx \lim_{\underline{I}} \operatorname{Mor}(\underline{Y} \otimes \Delta_{\underline{i}}, \underline{Z})$$
$$\approx \lim_{\underline{I}} \operatorname{Mor}(\underline{Y}, \operatorname{lhom}(\Delta_{\underline{i}}, \underline{Z}))$$
$$\approx \operatorname{Mor}(\underline{Y}, \operatorname{lim}_{\underline{I}} \operatorname{lhom}(\Delta_{\underline{i}}, \underline{Z}))$$

$$1hom(colim_{\underline{I}} \Delta_{\underline{i}}, Z) \approx \lim_{\underline{I}} 1hom(\Delta_{\underline{i}}, Z).$$

Let <u>C</u> be a monoidal category. Given $X \in Ob \underline{C}$, an object ${}^{\vee}X \in Ob \underline{C}$ is said to be a <u>left dual</u> of X if \exists morphisms

$$\begin{bmatrix} & & \\ & & \\ & & \\ & & \\ & & \\ & &$$

and commutative diagrams

=>

N.B. When C is strict, these diagrams reduce to the relations

$$\begin{array}{c} - (\operatorname{id}_X \otimes \varepsilon_X) \circ (\operatorname{n}_X \otimes \operatorname{id}_X) = \operatorname{id}_X \\ (\varepsilon_X \otimes \operatorname{id}_V) \circ (\operatorname{id}_V \otimes \operatorname{n}_X) = \operatorname{id}_V \\ X & X & X \end{array}$$

5.5 <u>LEMMA</u> Suppose that ^VX is a left dual of X — then the functor — \otimes ^VX is a right adjoint for the functor — \otimes X and the functor ^VX \otimes — is a left adjoint for the functor X \otimes —.

In brief: $\forall Y, Z \in Ob C$,

$$\begin{array}{|c|c|c|c|c|} \hline & \operatorname{Mor}(Y \ \mathfrak{B} \ X, Z) & \approx \operatorname{Mor}(Y, Z \ \mathfrak{B} \ \ ^{\vee} X) \\ & \operatorname{Mor}(^{\vee} X \ \mathfrak{B} \ Y, Z) & \approx \operatorname{Mor}(Y, X \ \mathfrak{B} \ Z) \, . \end{array}$$

<u>PROOF</u> It will be enough to show that $- \mathbf{Q}^{\vee} X$ is a right adjoint for $- \mathbf{Q} X$, the proof that $^{\vee} X \mathbf{Q}$ - is a left adjoint for $X \mathbf{Q}$ - being similar. So let

$$\begin{bmatrix} F = - & X \\ & (cf. 2.6) \\ G = - & X \end{bmatrix}$$

and to simplify the writing, take C strict. Define

$$\downarrow \in \operatorname{Nat}(\operatorname{id}_{\underline{C}}, \mathbf{G} \circ \mathbf{F})$$
$$\lor \in \operatorname{Nat}(\mathbf{F} \circ \mathbf{G}, \operatorname{id}_{\underline{C}})$$

by

$$\begin{array}{|c|c|c|c|} & \mu_{W} \in \operatorname{Mor}(W, W \otimes X \otimes VX) \\ & \mu_{W} = \operatorname{id}_{W} \otimes \eta_{X} \\ & &$$

$$v_{W} \in Mor(W \otimes X \otimes X, V)$$

 $v_{W} = id_{W} \otimes e_{X}.$

Consider

Thus

$$((\nabla F) \circ (F\mu))_{W} = (\nabla F)_{W} \circ (F\mu)_{W}$$

And

$$(\mu) \begin{vmatrix} F \mu \in Nat(F,FGF) \\ (F\mu)_{W}:FW \rightarrow FGFW \end{vmatrix}$$

or still,

(v)
$$\bigvee_{VF \in Nat(FGF,F)}^{VF \in Nat(FGF,F)}$$

(v) $\bigvee_{VF}^{VF}_{W}^{VFGFW \rightarrow FW}$

or still,

$$(vF)_{W} \overset{\mathrm{id}_{X} & \mathrm{id}_{X} & \varepsilon_{X} \\ (vF)_{W} \overset{\mathrm{v}_{X} & \mathrm{v}_{X} & \mathrm{v}_{X} & \mathrm{w} & \mathrm{w}_{X} \\ \end{array}$$

Therefore

$$(\forall \mathbf{F})_{W} \circ (\mathbf{F}\mu)_{W} \in Mor(W \otimes \mathbf{X}, W \otimes \mathbf{X})$$

is the composition

$$(id_{W} \otimes id_{X} \otimes \varepsilon_{X}) \circ (id_{W} \otimes n_{X} \otimes id_{X})$$

$$= (id_{W} \circ id_{W}) \otimes ((id_{X} \otimes \varepsilon_{X}) \circ (n_{X} \otimes id_{X}))$$

$$= id_{W} \otimes id_{X}$$

$$= id_{W} \otimes x$$

$$= id_{FW}$$

$$= (id_{F})_{W}$$

I.e.:

 $(\nabla F) \circ (F\mu) = id_{\overline{F}}$.

The verification that

 $(Gv) \circ (\mu G) = id_{G}$

is analogous.

5.6 <u>LEMMA</u> A left dual of X, if it exists, is unique up to isomorphism. <u>PROOF</u> Suppose that $\begin{bmatrix} & & \\ & & \\ & & \\ & & \\ & & & & \\ & & & \\ & & & & \\ & & & \\ & &$

are two left duals of X -- then the functors

are naturally isomorphic (both being right adjoints for — @ X), so $\forall \ W \in Ob \ \underline{C},$

$$w \otimes x_1 \approx w \otimes x_2$$
.

Now specialize and take W = e to get

=>

$$x_1 \approx x_2$$

[Note: Explicated,



$$\varepsilon_{\mathbf{X}}^{\prime} = \varepsilon_{\mathbf{X}}^{\prime} \circ (\phi^{-1} \ \boldsymbol{\Omega} \ \mathrm{id}_{\mathbf{X}})$$
$$\eta_{\mathbf{X}}^{\prime} = (\mathrm{id}_{\mathbf{X}} \ \boldsymbol{\Omega} \ \phi) \circ \eta_{\mathbf{X}}^{\prime}.$$

Then the triple $({}^{\vee}X^{*}, \epsilon_{X}^{*}, \eta_{X}^{*})$ is a left dual of X.

[Consider first the case when \underline{C} is strict, thus, e.g.,

$$(\operatorname{id}_X \mathfrak{Q} \mathfrak{e}_X') \circ (\mathfrak{n}_X' \mathfrak{Q} \operatorname{id}_X)$$

$$= \operatorname{id}_{X} \mathfrak{Q} (\varepsilon_{X} \circ (\phi^{-1} \mathfrak{Q} \operatorname{id}_{X})) \circ ((\operatorname{id}_{X} \mathfrak{Q} \phi) \circ \eta_{X}) \mathfrak{Q} \operatorname{id}_{X}$$
$$= \operatorname{id}_{X} \mathfrak{Q} \varepsilon_{X} \circ \operatorname{id}_{X} \mathfrak{Q} (\phi^{-1} \mathfrak{Q} \operatorname{id}_{X}) \circ (\operatorname{id}_{X} \mathfrak{Q} \phi) \mathfrak{Q} \operatorname{id}_{X} \circ \eta_{X} \mathfrak{Q} \operatorname{id}_{X}.$$

But

_....

$$(\mathrm{id}_X \otimes \phi) \otimes \mathrm{id}_X = \mathrm{id}_X \otimes (\phi \otimes \mathrm{id}_X)$$

=>

$$id_{X} \ \emptyset \ (\phi^{-1} \ \emptyset \ id_{X}) \ \circ \ (id_{X} \ \emptyset \ \phi) \ \emptyset \ id_{X}$$

$$= id_{X} \ \emptyset \ (\phi^{-1} \ \emptyset \ id_{X}) \ \circ \ id_{X} \ \emptyset \ (\phi \ \emptyset \ id_{X})$$

$$= id_{X} \ \emptyset \ (\phi^{-1} \ \emptyset \ id_{X}) \ \circ \ (\phi \ \emptyset \ id_{X})$$

$$= id_{X} \ \emptyset \ id_{V_{X} \ \emptyset \ X}$$

$$= id_{X} \ \emptyset \ id_{V_{X} \ \emptyset \ X}$$

=>

$$\operatorname{id}_X \otimes \varepsilon'_X \circ n'_X \otimes \operatorname{id}_X$$

$$= (\mathrm{id}_X \, \otimes \, \varepsilon_X) \circ (n_X \, \otimes \, \mathrm{id}_X) = \mathrm{id}_X.$$

In general, the claim is that $\operatorname{id}_{\boldsymbol{X}}$ equals

$$R \circ (id_X \otimes \epsilon_X') \circ A^{-1} \circ (\eta_X' \otimes id_X) \circ L^{-1}$$

or still,

$$\mathbf{R} \circ \operatorname{id}_{\mathbf{X}} \boldsymbol{\boldsymbol{\omega}} \ (\mathbf{e}_{\mathbf{X}} \circ (\phi^{-1} \ \boldsymbol{\boldsymbol{\omega}} \ \operatorname{id}_{\mathbf{X}})) \circ \mathbf{A}^{-1} \circ ((\operatorname{id}_{\mathbf{X}} \ \boldsymbol{\boldsymbol{\omega}} \ \phi) \circ \mathbf{\eta}_{\mathbf{X}}) \ \boldsymbol{\boldsymbol{\omega}} \ \operatorname{id}_{\mathbf{X}} \circ \mathbf{L}^{-1}$$

or still,

$$\mathbf{R} \circ \operatorname{id}_{\mathbf{X}} \otimes \varepsilon_{\mathbf{X}} \circ \operatorname{id}_{\mathbf{X}} \otimes (\phi^{-1} \otimes \operatorname{id}_{\mathbf{X}}) \circ \mathbf{A}^{-1} \circ (\operatorname{id}_{\mathbf{X}} \otimes \phi) \otimes \operatorname{id}_{\mathbf{X}} \circ \eta_{\mathbf{X}} \otimes \operatorname{id}_{\mathbf{X}} \circ \mathbf{L}^{-1}.$$

Here

$$A^{-1}: (X \otimes ^{\vee}X') \otimes X \to X \otimes (^{\vee}X' \otimes X).$$

So, to complete the verification, one has only to show that the composition

$$(X \otimes ^{\vee}X) \otimes X \xrightarrow{(\text{id } \otimes \phi) \otimes \text{id}} (X \otimes ^{\vee}X') \otimes X$$

$$\xrightarrow{A^{-1}} X \otimes (^{\vee}X' \otimes X)$$

$$\xrightarrow{\text{id } \otimes (\phi^{-1} \otimes \text{id})} X \otimes (^{\vee}X \otimes X)$$

$$A^{-1}$$

is

$$(X \otimes ^{\vee} X) \otimes X \xrightarrow{A} X \otimes (^{\vee} X \otimes X).$$

However, due to the naturality of the associativity constraint, there is a

commutative diagram

And

$$(id \otimes (\phi \otimes id))^{-1} = id \otimes (\phi^{-1} \otimes id).$$

A monoidal category <u>C</u> is said to be <u>left autonomous</u> if each object in <u>C</u> admits a left dual.

N.B. Suppose that C is left autonomous. Given $f\in Mor(X,Y)\,,$ define ${}^{\vee}f\in Mor({}^{\vee}Y,{}^{\vee}X)$ by

$$\begin{array}{c} {}^{v}\mathbf{Y} \xrightarrow{\mathbf{R}^{-1}} {}^{v}\mathbf{Y} \ \mathbf{\Theta} \ \mathbf{e} \\ & \begin{array}{c} \mathrm{id} \ \mathbf{\Theta} \ \eta_{\mathbf{X}} \\ & \begin{array}{c} - & \end{array} \\ & \begin{array}{c} \mathrm{id} \ \mathbf{\Theta} \ \eta_{\mathbf{X}} \\ & \begin{array}{c} - & \end{array} \\ & \begin{array}{c} \mathrm{id} \ \mathbf{\Theta} \ \eta_{\mathbf{X}} \\ \end{array} \end{array} \right) {}^{v}\mathbf{Y} \ \mathbf{\Theta} \ (\mathbf{X} \ \mathbf{\Theta} \ \mathbf{V} \mathbf{X}) \\ & \begin{array}{c} \mathbf{A} \\ & \begin{array}{c} - & \end{array} \\ & \begin{array}{c} \mathrm{id} \ \mathbf{\Theta} \ \mathbf{f} \right) \ \mathbf{\Theta} \ \mathrm{id} \\ \end{array} \right) {}^{v}\mathbf{Y} \ \mathbf{\Theta} \ \mathbf{X} \right) \ \mathbf{\Theta} \ \mathbf{V} \mathbf{X} \\ & \begin{array}{c} \mathrm{(id} \ \mathbf{\Theta} \ \mathbf{f} \right) \ \mathbf{\Theta} \ \mathrm{id} \\ \end{array} \\ & \begin{array}{c} \mathrm{e} \ \mathbf{\Theta} \ \mathbf{V} \\ \mathbf{\Phi} \\ \end{array} \right) {}^{v}\mathbf{Y} \ \mathbf{\Theta} \ \mathbf{Y} \right) \ \mathbf{\Theta} \ \mathbf{V} \mathbf{X} \\ & \begin{array}{c} \mathrm{e} \ \mathbf{\Psi} \ \mathbf{\Theta} \ \mathrm{id} \\ \end{array} \\ & \begin{array}{c} \mathrm{e} \ \mathbf{\Theta} \ \mathbf{V} \\ \end{array} \\ & \begin{array}{c} \mathrm{e} \ \mathbf{\Theta} \ \mathbf{V} \\ \mathbf{X} \\ \end{array} \end{array}$$

Then the assignment

$$\begin{bmatrix} - & X \rightarrow & X \\ & f \rightarrow & f \end{bmatrix}$$

defines a cofunctor $\underline{C} \rightarrow \underline{C}$.

[Note: The specific form of $^{\vee}f$ depends on the choices of $^{\vee}X$ and $^{\vee}Y$.]

5.8 <u>REMARK</u> If <u>C</u> is left autonomous and if $X, Y \in Ob \underline{C}$, then $^{\vee}(X \otimes Y)$ is isomorphic to $^{\vee}Y \otimes ^{\vee}X$.

[We have

 $Mor(^{\vee}(X \otimes Y) \otimes W, Z) \approx Mor(W, (X \otimes Y) \otimes Z)$ $\approx Mor(W, X \otimes (Y \otimes Z))$ $\approx Mor(^{\vee}X \otimes W, Y \otimes Z)$ $\approx Mor(^{\vee}Y \otimes (^{\vee}X \otimes W), Z)$ $\approx Mor((^{\vee}Y \otimes ^{\vee}X) \otimes W, Z)$ => $^{\vee}(X \otimes Y) \approx ^{\vee}Y \otimes ^{\vee}X,]$

5.9 <u>LEMMA</u> Suppose that <u>C</u> is left autonomous — then <u>C</u> is left closed. <u>PROOF</u> In fact, $\forall X \in Ob C$,

$$1hcm(X, --) = -- \otimes ^{\vee}X.$$

One can also introduce the notion of a right dual X^V of X, where this time

$$\begin{bmatrix} \varepsilon_X : X \otimes X^{\vee} \to e \\ \\ \eta_X : e \to X^{\vee} \otimes X \end{bmatrix}$$

subject to the obvious commutativity conditions. Here the functor — $\otimes X^{\vee}$ is a left adjoint for the functor — $\otimes X$ and the functor $X^{\vee} \otimes$ — is a right adjoint for the functor $X \otimes$ —.

[Note: If X admits a left dual $^{\vee}X$ and a right dual X^{\vee} , then in general $^{\vee}X$ and X^{\vee} are not isomorphic. On the other hand, it is true that

$$(^{\vee}X)^{\vee} \approx X \approx ^{\vee}(X^{\vee}).$$

E.g.:

Mor
$$(Y \otimes (^{\vee}X)^{\vee}, Z) \approx Mor (Y, Z \otimes ^{\vee}X) \approx Mor (Y \otimes X, Z)$$

=>
 $(^{\vee}X)^{\vee} \approx X.$]

The definition of "right autonomous" is clear and we shall term \underline{C} <u>autonomous</u> if it is both left and right autonomous.

5.10 <u>LEMMA</u> Suppose that <u>C</u> is right autonomous — then <u>C</u> is right closed. <u>PROOF</u> In fact, $\forall X \in Ob C$,

$$rhom(X, --) = X^{\vee} \otimes --$$
.

5.11 <u>REMARK</u> If <u>C</u> is autonomous, then - Q — preserves colimits in both variables.

Suppose that $F:\underline{C} \rightarrow \underline{C}'$ is a monoidal functor. Assume: X^{\vee} is a right dual of X — then FX^{\vee} is a right dual of FX. Proof: Consider the arrows

$$\begin{bmatrix} & \Xi & & F^{\varepsilon} X & \xi^{-1} \\ FX & FX^{\vee} & \longrightarrow & F(X & X^{\vee}) & \longrightarrow & Fe & \longrightarrow & e^{-1} \\ & \xi & & F^{\tau} X & & \Xi^{-1} \\ & e^{\tau} & \longrightarrow & Fe & \longrightarrow & F(X^{\vee} & Q & X) & \xrightarrow{\Xi^{-1}} & FX^{\vee} & Q^{\tau} & FX. \end{bmatrix}$$

[Note: Assume that C, C' are right autonomous - then there is an isomorphism

$$\Delta_{X}: \mathbf{FX}^{\vee} \rightarrow (\mathbf{FX})^{\vee},$$

namely the composition

$$FX^{\vee} \xrightarrow{L} e' \otimes' FX^{\vee}$$

$$\xrightarrow{\eta \otimes id} ((FX)^{\vee} \otimes' FX) \otimes' FX^{\vee}$$

$$\xrightarrow{A^{-1}} (FX)^{\vee} \otimes' (FX \otimes' FX^{\vee})$$

$$\xrightarrow{id \otimes E} (FX)^{\vee} \otimes' F(X \otimes X^{\vee})$$

$$\xrightarrow{id \otimes F\epsilon} (FX)^{\vee} \otimes' Fe$$

$$\xrightarrow{id \otimes \xi^{-1}} (FX)^{\vee} \otimes' e'$$

$$\xrightarrow{R} (FX)^{\vee} \otimes' e'$$

and the diagram



commutes.]

N.B. One can, of course, work equally well with left duals.

5.12 LEMMA Let

be monoidal functors and let $\alpha: F \rightarrow G$ be a monoidal natural transformation. Assume: The source <u>C</u> of F and G is autonomous — then α is a monoidal natural isomorphism.

PROOF The claim is that $\forall X \in Ob \ \underline{C}$,

$$\alpha_{\mathbf{x}}: \mathbf{F}\mathbf{X} \to \mathbf{G}\mathbf{X}$$

is an isomorphism. From the above, FX^{\vee} (GX^{\vee}) is a right dual of FX (GX) or still, FX (GX) is a left dual of FX^{\vee} (GX^{\vee}). This said, form

$$\alpha_{X^{\vee}}:FX^{\vee} \to GX^{\vee}$$

and consider

$$(\alpha)$$
 :GX \rightarrow FX.

Suppose that <u>C</u> is symmetric monoidal and left autonomous — then <u>C</u> is right autonomous, hence <u>C</u> is autonomous. Proof: Given $X \in Ob \underline{C}$, take $X^{\vee} = {}^{\vee}X$ and define morphisms

 $\begin{bmatrix} & \varepsilon_{X} & \bullet & T \\ & & X, & X \end{bmatrix}$

by

5.13 EXAMPLE FOVECT is autonomous. In fact, FOVECT is symmetric monoidal, so it suffices to set up a left duality. Thus given X, let
$$^{\vee}X$$
 be its dual and define

by

$$\varepsilon_{\mathbf{X}}(\lambda, \mathbf{x}) = \lambda(\mathbf{x}).$$

On the other hand, there is a canonical isomorphism

$$\phi: \operatorname{Hom}(X, X) \rightarrow \operatorname{Hom}(k, {}^{V}X \otimes X)$$

and we let

$$\eta_{\mathbf{X}} = \phi(\mathrm{id}_{\mathbf{X}}).$$

[Note: An object X in $\underline{\text{VEC}}_k$ admits a left dual iff it is finite dimensional.]

5.14 <u>EXAMPLE</u> The full subcategory of $\underline{MOD}_{\underline{k}}$ whose objects are finitely generated projective is autonomous (cf. 5.2).

Assume still that C is symmetric monoidal and left autonomous.

5.15 IEMMA There is a monoidal natural isomorphism

$$\operatorname{id}_{\underline{C}} \neq \operatorname{vv}(--)$$
.

[To see this, consider the composition

$$\begin{array}{c}
\mathbf{R}^{-1} \\
\mathbf{X} & \xrightarrow{\mathbf{R}^{-1}} & \mathbf{X} \otimes \mathbf{e} \\
\xrightarrow{\mathrm{id} \otimes n} & \mathbf{X} \otimes \mathbf{e} \\
\xrightarrow{\mathbf{M}} & \mathbf{X} \otimes \mathbf{V}^{\mathsf{V}} \mathbf{X} \otimes \mathbf{V}^{\mathsf{V}} \mathbf{X} \\
\xrightarrow{\mathbf{A}} & \mathbf{X} \otimes \mathbf{V}^{\mathsf{V}} \mathbf{X} \otimes \mathbf{V}^{\mathsf{V}} \mathbf{X} \\
\xrightarrow{\mathbf{T} \otimes \mathrm{id}} & \mathbf{V}^{\mathsf{V}} \mathbf{X} \otimes \mathbf{X} \otimes \mathbf{V}^{\mathsf{V}} \mathbf{X} \\
\xrightarrow{\mathbf{C} \otimes \mathrm{id}} & \mathbf{e} \otimes \mathbf{V}^{\mathsf{V}} \mathbf{X} \\
\xrightarrow{\mathbf{L}} & \xrightarrow{\mathbf{V}^{\mathsf{V}}} \mathbf{X} \cdot \mathbf{I}
\end{array}$$

N.B. Let

be the arrow constructed above -- then

$$(\delta_{X})^{-1} = {}^{\vee}(\delta_{X})$$
 (cf. 5.12).

But here $X^{\vee} = {}^{\vee}X$, so

$$(\delta_{\mathbf{X}})^{-1} = {}^{\vee}(\delta_{\mathbf{V}_{\mathbf{X}}}).$$

[Note: To make sense of this, recall that

X is a left dual of
$$X^{\vee}$$

V'X is a left dual of $V^{\vee}(X^{\vee})$.

And

$$\delta_{X^{V}}: X^{V} \to {}^{VV}(X^{V})$$

=>

$$(\delta_{X^{\vee}}): \overset{\vee}{X} \to X.$$

§6. TWISTS

Let <u>C</u> be symmetric monoidal and left autonomous — then a <u>twist</u> Ω is a monoidal natural isomorphism of the identity functor id_{C} such that $\forall X \in \operatorname{Ob} C$,

$$(\Omega_X \otimes id_{\vee_X}) \circ \eta_X = (id_X \otimes \Omega_{\vee_X}) \circ \eta_X.$$

[Note: Tacitly, $id_{\underline{C}}$ is taken to be strict ($\xi = id$, $\Xi = id$), thus from the definitions

$$\Omega_{\mathbf{X} \otimes \mathbf{Y}} = \Omega_{\mathbf{X}} \otimes \Omega_{\mathbf{Y}} \text{ and } \Omega_{\mathbf{e}} = \mathrm{id}_{\mathbf{e}}.$$

To consolidate the terminology, a symmetric monoidal <u>C</u> which is left autonomous and has a twist Ω will be referred to as a ribbon category.

<u>N.B.</u> The choice $\Omega_X = id_X$ is permissible, in which case <u>C</u> is said to be even.

It was pointed out near the end of §5 that an even ribbon category is right autonomous. This fact is true in general. Proof: Given $X \in Ob \subseteq$, take $X^{\vee} = {}^{\vee}X$ and define morphisms

$$\begin{bmatrix} x \otimes x^{\vee} \\ e \rightarrow x^{\vee} \otimes x \end{bmatrix}$$

by

$$\begin{bmatrix} \varepsilon_{\mathbf{X}} \circ \mathbf{\tau} & \mathbf{x}, \mathbf{v}_{\mathbf{X}} \circ \mathbf{\Omega}_{\mathbf{X}} \otimes \mathbf{id}_{\mathbf{v}_{\mathbf{X}}} \\ \mathbf{id}_{\mathbf{v}_{\mathbf{X}}} \otimes \mathbf{\Omega}_{\mathbf{X}} \circ \mathbf{\tau} & \mathbf{v}_{\mathbf{v}} \circ \mathbf{n}_{\mathbf{X}} \\ \end{bmatrix}$$

6.1 LEMMA In the presence of a twist Ω ,

$$X \approx (^{\vee}X)$$
.

PROOF Consider the composition

$$X \xrightarrow{\mathbb{R}^{-1}} X \otimes e$$

$$\xrightarrow{\Omega_X \otimes \eta_{V_X}} X \otimes (^{V_X} \otimes ^{V} (^{V_X}))$$

$$\xrightarrow{\mathbb{A}} (X \otimes ^{V_X}) \otimes ^{V} (^{V_X})$$

$$\xrightarrow{T_{V_Y} \otimes id}$$

$$\xrightarrow{X, \forall X} (\forall X \otimes X) \otimes \forall (\forall X)$$

$$\stackrel{\varepsilon_{X} \otimes id}{\longrightarrow} e \otimes (^{\vee}X)$$

$$\stackrel{L}{\longrightarrow} (^{\vee}X).$$

E.g.:

$${}^{\vee}e \approx {}^{\vee}e \otimes e \approx {}^{\vee}e \otimes {}^{\vee}({}^{\vee}e) \approx {}^{\vee}({}^{\vee}e \otimes e) \approx {}^{\vee}({}^{\vee}e) \approx e.$$

6.2 <u>LEMMA</u> In the presence of a twist Ω , the left and right dual of every morphism $f:X \rightarrow Y$ agree: ${}^{\vee}f = f^{\vee}$.

Let \underline{C} be a ribbon category. Given $f \in Mor(X,X)$, define the trace of f by

$$\operatorname{tr}_{X}(f) = \varepsilon_{X} \circ \tau \circ \Omega_{X} \otimes \operatorname{id}_{V_{X}} \circ (f \otimes \operatorname{id}_{V_{X}}) \circ \eta_{X}$$

[Note:

$$tr_X(f) \in Mor(e,e) (= \underline{M}(\underline{C})).]$$

6.3 LEMMA We have

1.
$$\operatorname{tr}_{X}(f) = \operatorname{tr}_{V_{X}}({}^{V}f);$$

2. $\operatorname{tr}_{X}(g \circ f) = \operatorname{tr}_{Y}(f \circ g) \quad (f:X \to Y, g:Y \to X);$
3. $\operatorname{tr}_{X_{1}} \otimes X_{2} \quad (f_{1} \otimes f_{2}) = \operatorname{tr}_{X_{1}}(f_{1}) \quad \operatorname{tr}_{X_{2}}(f_{2}).$

 Put

$$\dim X = \operatorname{tr}_{X}(\operatorname{id}_{X}),$$

the dimension of X.

So, on the basis of 6.3,

$$\dim X = \dim X$$

and

$$\dim (X \otimes Y) = (\dim X) (\dim Y).$$

N.B. Take Ω = id -- then the <u>categorical dimension</u> of X is the arrow

$$e \xrightarrow{\eta_{X}} x \otimes {}^{\vee}x \xrightarrow{}^{\top}x, {}^{\vee}x \xrightarrow{}^{\vee}x \otimes x \xrightarrow{}^{\varepsilon}x = e.$$

6.4 <u>EXAMPLE</u> Consider <u>FDVEC</u> (viewed as an even ribbon category (cf. 5.13)) then the trace of $f:X \rightarrow X$ is the composition

Therefore the abstract definition of $tr_{\chi}(f)$ is the usual one. In particular:

$$\dim X = (\dim_{\underline{k}} X) \mathbf{1}_{\underline{k}}$$

E.g.:

$$\dim \underline{k}^n = nl_{\underline{k}'}$$

the distinction between $n \in N$ and nl_k being potentially essential if k has non-zero characteristic.

6.5 <u>REMARK</u> While evident, it is important to keep in mind that the definitions of trace and dimension depend on all the underlying assumptions, viz. that our monoidal <u>C</u> is symmetric, left autonomous, and has a twist Ω .

Suppose that <u>C</u>, <u>C'</u> are ribbon categories with respective twists Ω , Ω' -then a symmetric monoidal functor $F:\underline{C} \neq \underline{C}'$ is <u>twist preserving</u> if $\forall X \in Ob \underline{C}$,

$$\mathbf{F}\Omega_{\mathbf{X}} = \Omega_{\mathbf{F}\mathbf{X}}^{\mathbf{I}}$$

6.6 LEMMA If $F: C \rightarrow C'$ is twist preserving, then $\forall f \in Mor(X,X)$, the diagram



commutes.

Matters are invariably simpler if \underline{C} is a strict ribbon category, which will be the underlying supposition in 6.7 - 6.9 below.

6.7 LEMMA The arrows

$$= \eta_e : e^{+v_e}$$
$$= \varepsilon_e : e^{+v_e}$$

are mutually inverse isomorphisms.

PROOF Take X = e in the relation

$$(\operatorname{id}_X \otimes \varepsilon_X) \circ (\operatorname{n}_X \otimes \operatorname{id}_X) = \operatorname{id}_X$$

to see that

$$\varepsilon_{e} \circ \eta_{e} = id_{e}$$

Now fix an isomorphism $\varphi \! : \! e \ \! \to \ \! ^{\vee} \! e \ \! - \! then$

=>

$$(\phi^{-1} \circ n_{e}) \circ (\varepsilon_{e} \circ \phi) = (\varepsilon_{e} \circ \phi) \circ (\phi^{-1} \circ n_{e}) \quad (cf. 1.4)$$
$$= \varepsilon_{e} \circ n_{e} = id_{e}$$
$$\Longrightarrow$$
$$n_{e} \circ \varepsilon_{e} = id_{e}.$$

$$tr_e(s) = s.$$

PROOF In fact,

$$tr_{e}(s) = \varepsilon_{e} \circ \tau_{e} \circ \Omega_{e} \otimes id_{v} \circ (s \otimes id_{v}) \circ \eta_{e}$$

$$= \varepsilon_{e} \circ id_{v} \circ id_{v} \circ (s \otimes id_{v}) \circ \eta_{e}$$

$$= \varepsilon_{e} \circ (s \otimes id_{v}) \circ \eta_{e}$$

$$= (id_{e} \otimes \varepsilon_{e}) (s \otimes id_{e} \otimes id_{v}) (id_{e} \otimes \eta_{e})$$

$$= s \otimes (\varepsilon_{e} \circ \eta_{e})$$

$$= s \otimes id_{e}$$

$$= s.$$

[Note: Therefore

$$\dim e = \operatorname{tr}_{e}(\operatorname{id}_{e}) = \operatorname{id}_{e}.$$

6.9 LEMMA $\forall x \in Ob \underline{C}$,

$$\Omega_{V_{X}} = {}^{V}\Omega_{X}.$$

PROOF The compositions

$$\begin{bmatrix} & & & & & & & & & \\ & e & \longrightarrow & X & & ^{\vee}X & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & & \\ & & & & & & & & & \\ & & & & & & & \\ & & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & & & & & & & \\ & & & & & & & \\ & & &$$

are equal, thus the compositions

$$\begin{array}{c} \stackrel{id \ \ensuremath{\varpi}\ \ensuremath{\eta_X}\ \ensuremath{\sim}\ \ensuremath{u}\ \ensuremath{\sim}\ \ensuremath{u}\ \ensuremath{\sim}\ \ensuremath{u}\ \ensuremath{\omega}\ \ensuremath{u}\ \ensuremath{\omega}\ \ensuremath{\omega$$

are equal. Postcompose with $\epsilon_X \otimes id_{\bigvee X}$ — then the first line gives \bigvee_X^{\bigvee} , while the second line is

$$\overset{\varepsilon_{\mathbf{X}} \otimes \operatorname{id}_{\mathbf{V}} \circ \operatorname{id}_{\mathbf{V}} \otimes \operatorname{id}_{\mathbf{X}} \otimes \Omega_{\mathbf{V}} \circ \operatorname{id}_{\mathbf{V}} \otimes \eta_{\mathbf{X}} }_{\mathbf{X} \times \mathbf{X} \times \mathbf$$

or still,

or still,

$$(\operatorname{id}_{e} \mathfrak{Q} \Omega_{v}) \circ (\varepsilon_{X} \mathfrak{Q} \operatorname{id}_{v}) \circ \operatorname{id}_{v} \mathfrak{Q} \eta_{X}$$

or still,

$$\begin{array}{c} \Omega_{v} \circ id_{v} = \Omega_{v} \\ x & x \end{array}$$

6.10 <u>REMARK</u> Let <u>C</u> be a ribbon category — then this structure can be transferred to \underline{C}_{str} . That the symmetry τ passes to a symmetry τ_{str} of \underline{C}_{str} was noted already in 4.6. As for the left duality, a generic element of \underline{C}_{str} is a finite sequence (X_1, \ldots, X_n) and

$$(x_1, ..., x_n) = (x_n, ..., x_1),$$

where ε and η are defined in the obvious way. It is also clear that the twist

8.

on <u>C</u> can be brought over to a twist on <u>C</u>_{str}. Accordingly, $\gamma:\underline{C} \rightarrow \underline{C}_{str}$ is a symmetric monoidal equivalence which is twist preserving, i.e., $\gamma:\underline{C} \rightarrow \underline{C}_{str}$ is a ribbon equivalence.

§7. *-CATEGORIES

Let <u>k</u> be a commutative ring -- then a category <u>C</u> is <u>k</u>-enriched if $\forall X, Y \in Ob \underline{C}$, Mor(X,Y) is a <u>k</u>-module and if the composition of morphisms is <u>k</u>-bilinear. A functor F between <u>k</u>-enriched categories is <u>k</u>-linear if the induced maps

$$Mor(X,Y) \rightarrow Mor(FX,FY)$$

are homomorphisms of k-modules.

[Note: If <u>C</u> is <u>k</u>-enriched and monoidal, then <u>C</u> × <u>C</u> is <u>k</u>-enriched and the functor $\&: C \times C \rightarrow C$ is assumed to be <u>k</u>-bilinear.]

<u>N.B.</u> An object X in a <u>k</u>-enriched category <u>C</u> is <u>irreducible</u> if Mor(X,X) = kid_X.

7.1 EXAMPLE Suppose that C is Z-enriched and monoidal. Put

 $\underline{\mathbf{k}} = \underline{\mathbf{M}}(\underline{\mathbf{C}})$.

Then <u>k</u> is a unital commutative ring (cf. 1.4) and <u>C</u> is <u>k</u>-enriched as a monoidal category (cf. 1.5).

[Note: Suppose in addition that C is a ribbon category -- then $\forall X \in Ob C$,

$$\operatorname{tr}_{X}:\operatorname{Mor}(X,X) \to \underline{k}$$

is k-linear and $\forall X, Y \in Ob C$, the map

$$= \operatorname{Mor}(X,Y) \otimes_{\underline{k}} \operatorname{Mor}(Y,X) + \underline{k}$$
$$f \otimes g \to \operatorname{tr}_{X}(g \circ f)$$

is k-bilinear.]

A <u>*-category</u> is a pair (C,*), where C is a category enriched over the field of complex numbers and

is an involutive, identity on objects, positive cofunctor. Spelled out: $\forall X, Y \in Cb \ \underline{C}, Mor(X, Y)$ is a complex vector space, composition

$$Mor(X,Y) \times Mor(Y,Z) \rightarrow Mor(X,Z)$$

is complex bilinear,

$$*:Mor(X,Y) \rightarrow Mor(Y,X)$$

subject to

$$(zf + wg)^* = zf^* + wg^*$$

and $\$

Finally, the requirement that * be positive means:

$$f^* \circ f = 0 \Rightarrow f = 0$$
.

[Note: $\forall X \in Ob \subseteq$, we have

$$id_{X}^{*} = id_{X} \circ id_{X}^{*}$$
$$= id_{X}^{**} \circ id_{X}^{*}$$
$$= (id_{X} \circ id_{X}^{*})^{**}$$
$$= id_{X}^{***}$$
$$= id_{X}^{**}$$
N.B. A monoidal *-category is a *-category which is monoidal with

$$(f \ \ g)^* = f^* \ \ g^*$$

for all f,g.

[Note: A symmetric monoidal *-category is a monoidal *-category such that $\forall X, Y \in Ob C$,

is unitary (see below).]

7.2 EXAMPLE FDHILB is a symmetric monoidal *-category.

[Note: For the record, <u>FDHILB</u> is a construct. As such, it is amnestic and transportable, thus there is no loss of generality in assuming that its monoidal structure is strict (cf. 3.12).]

7.3 <u>REMARK</u> Let A be a complex *-algebra -- then the involution is positive if $A^* \circ A = 0 \Rightarrow A = 0$ ($A \in A$). To illustrate, take $A = M_2(C)$ and consider the involutions

A ¹ =	- a11	^a 12		a ₂₁
	_ ^a 21	^a 22 _	 ^a _12	a ₂₂
*2 =	a 11	^a 12	- a ₂₂	a12
	^a 21	^a 22	= ^a 21	

Then $*_1$ is positive but $*_2$ is not positive since

0	1 *2	0	1	o	0
0	o _	_ 0	0 _	0	o _ .

[Note: It is wellknown that if A is finite dimensional and if the involution is positive, then A is a semisimple algebra, hence "is" a multimatrix algebra.]

Let $f:X \rightarrow Y$ be a morphism in a *-category <u>C</u> -- then f is an <u>isometry</u> if $f^* \circ f = id_X$ and f is <u>unitary</u> if both f and f^* are isometries.

Let F be a C-linear functor between *-categories -- then F is \star -preserving if \forall f, F(f*) = (Ff)*.

<u>N.B.</u> Suppose that F is a *-preserving monoidal functor between monoidal *-categories -- then F is unitary if the isomorphisms $\xi:e^{+} \rightarrow Fe$ and

$$E_{X,Y}:FX \ Q' \ FY \rightarrow F(X \ Q \ Y)$$

are unitary.

Let $p:X \rightarrow X$ be a morphism in a *-category <u>C</u> -- then p is a projection if $p = p^*$ and $p \circ p = p$.

[Note: If $g: Y \rightarrow X$ is an isometry, then $g \circ g^*: X \rightarrow X$ is a projection.]

Let <u>C</u> be a *-category and let $X, Y \in Ob \subseteq$ -- then X is a <u>subobject</u> of Y if \exists an isometry $f \in Mor(X, Y)$.

Definition: \underline{C} has subobjects if for any $Y \in Ob \underline{C}$ and any projection $q \in Mor(Y,Y), \exists X \in Ob \underline{C}$ and an isometry $f \in Mor(X,Y)$ such that $f \circ f^* = q$. Definition: \underline{C} has direct sums if for all $X,Y \in Ob \underline{C}$, $\exists Z \in Ob \underline{C}$ and isometries $f \in Mor(X,Z), g \in Mor(Y,Z)$ such that $f \circ f^* + g \circ g^* = id_Z$. E.g.: FDHILB has subobjects and direct sums.

7.4 <u>RAPPEL</u> A category <u>C</u> is <u>essentially small</u> if <u>C</u> is equivalent to a small category.

Suppose that <u>C</u> is a *-category which is essentially small -- then <u>C</u> is semisimple if the following conditions are met:

 $\underline{SS_1}: \forall X, Y \in Ob \underline{C},$

dim Mor $(X,Y) < \infty$.

 SS_2 : <u>C</u> has subobjects and direct sums.

 $SS_3: \underline{C}$ has a zero object.

N.B. A monoidal *-category is <u>semisimple</u> if it is semisimple as a *-category and if in addition, e is irreducible.

7.5 EXAMPLE FDHILB is a semisimple strict monoidal *-category (cf. 7.2).

7.6 <u>LEMMA</u> Suppose that C is a semisimple *-category -- then every nonzero object in C is a finite direct sum of irreducible objects.

 $[\forall X \in Ob \underline{C}, Mor(X,X) \text{ is a finite dimensional complex *-algebra and the involution *:Mor(X,X) + Mor(X,X) is positive (cf. 7.3).]$

[Note: Conventionally, zero objects are not irreducible.]

Therefore a semisimple *-category is abelian.

Given a semisimple *-category C, denote its set of isomorphism classes of irreducible objects by I_{C} and let $\{X_{i}: i \in I_{C}\}$ be a set of representatives -- then

$$i \neq j \Rightarrow Mor(X_i, X_j) = \{0\}$$

and $\forall X \in Ob C$, \exists a finite number of i such that

$$Mor(X_i, X) \neq \{0\},\$$

thereby defining $I_X \subset I_C$.

7.7 <u>REMARK</u> $\forall i \in I_X$, Mor (X_i, X) is a finite dimensional Hilbert space with inner product

$$\langle \phi, \psi \rangle$$
 id_x = $\phi^* \circ \psi$.

7.8 LEMMA Let C, C' be semisimple *-categories and suppose that $F:C \rightarrow C'$ is C-linear -- then F is faithful if FX is nonzero for every irreducible X.

<u>PROOF</u> Consider an $f \in Mor(X,Y)$: Ff = 0, the claim being that f = 0. Fix orthonormal bases

$$\begin{bmatrix} s_{ik} \in Mor(X_i, X) & (k = 1, ..., \dim Mor(X_i, X)) \\ t_{j\ell} \in Mor(Y_j, Y) & (\ell = 1, ..., \dim Mor(Y_j, Y)) \end{bmatrix}$$

such that

$$\begin{bmatrix} \Sigma & s_{ik} \circ s_{ik}^* = id_X \\ ik & ik \end{bmatrix}$$
$$\begin{bmatrix} \Sigma & t_{j\ell} \circ t_{j\ell}^* = id_Y \\ j\ell & j\ell \end{bmatrix}$$

Write

$$f = id_{y} \circ f \circ id_{x}$$

$$= \sum_{\substack{ik,jl \\ ik,jl \\ ik}} t_{jl} \circ t_{jl} \circ f \circ s_{ik} \circ s_{ik}^{*}$$
$$= \sum_{\substack{ikl \\ ikl \\ ik}} c_{ikl} t_{il} \circ s_{ik}^{*} (\exists c_{ikl} \in C).$$

Then for indices m, μ, ν ,

$$0 = F(t_{m_{v}}^{*}) \circ Ff \circ F(s_{m\mu})$$

$$= \sum_{ik\ell} c_{ik\ell} F(t_{m_{v}}^{*} \circ t_{i\ell} \circ s_{ik}^{*} \circ s_{m\mu})$$

$$= \sum_{k\ell} c_{mk\ell} F(t_{m_{v}}^{*} \circ t_{m\ell} \circ s_{mk}^{*} \circ s_{m\mu})$$

$$= \sum_{k\ell} c_{mk\ell} F(id_{x_{m}} \circ id_{x_{m}})$$

$$= c_{m\mu_{v}} F(id_{x_{m}})$$

$$= c_{m\mu_{v}} id_{Fx_{m}}.$$

But by assumption, $\operatorname{id}_{FX}_{m} \neq 0$, thus the $\operatorname{c}_{m\mu\nu}$ vanish, so f = 0.

7.9 <u>LEMMA</u> Let <u>C</u>, <u>C'</u> be semisimple *-categories and suppose that $F:\underline{C} \neq \underline{C'}$ is C-linear and faithful -- then F is full iff (a) $X \in Ob \underline{C}$ irreducible => $FX \in Ob \underline{C'}$ irreducible and (b) $X, Y \in Ob \underline{C}$ irreducible and nonisomorphic => $FX, FY \in Ob \underline{C'}$ irreducible and nonisomorphic.

§8. NATURAL TRANSFORMATIONS

Let \underline{C} , \underline{C}' be *-categories and let $F:\underline{C} \rightarrow \underline{C}'$ be a *-preserving functor.

8.1 LEMMA Nat(F,F) is a unital *-algebra under the following operations:

$$(a\alpha + b\beta)_{X} = a\alpha_{X} + b\beta_{X}$$
$$(\alpha \circ \beta)_{X} = \alpha_{X} \circ \beta_{X}$$
$$(\alpha^{*})_{X} = (\alpha_{X})^{*}$$
$$\mathbf{l}_{X} = \mathbf{id}_{FX}.$$

[To check the *-condition, observe that $\forall f \in Mor(X, Y)$,

Ff

•
$$(\alpha^*)_X = Ff \circ (\alpha_X)^*$$

= $(Ff^*)^* \circ (\alpha_X)^*$
= $(\alpha_X \circ Ff^*)^*$
= $(Ff^* \circ \alpha_Y)^*$
= $(\alpha_Y)^* \circ (Ff^*)^*$
= $(\alpha^*)_Y \circ Ff.]$

8.2 EXAMPLE Take $\underline{C}' = \underline{FDHILB}$, put $\operatorname{Nat}_{F} = \operatorname{Nat}(F,F)$, and let $\underline{\operatorname{Rep}}_{fd}$ Nat_{F} be the category whose objects are the finite dimensional *-representations of

Nat and whose morphisms are the intertwining operators. Define a \star -preserving functor

$$\Phi: \underline{C} \rightarrow \underline{\text{Rep}}_{fd} \text{Nat}_{F}$$

as follows:

$$\Phi X = (\pi_X, FX) \quad (X \in Ob \ \underline{C})$$
$$\Phi f = Ff \quad (f \quad Mor(X, Y)).$$

Here

$$\pi_{X}(\alpha) = \alpha_{X},$$

thus the diagram



commutes, so Ff is an intertwining operator.

[Note: If

$$U: \underline{\operatorname{Rep}}_{fd} \operatorname{Nat} F \to \underline{FDHILB}$$

is the forgetful functor, i.e., $U(\pi, H) = H$, then $U \circ \Phi = F$.]

8.3 <u>THEOREM</u> Let C, C' be *-categories and let $F: C \rightarrow C'$ be a *-preserving functor. Assume: C is semisimple -- then there is an isomorphism

$$\Psi_{\mathbf{F}}: \mathsf{Nat}(\mathbf{F}, \mathbf{F}) \rightarrow \underset{i \in \mathbf{I}_{\underline{C}}}{\uparrow \uparrow} \mathsf{Mor}(\mathbf{F} \mathbf{X}_{i}, \mathbf{F} \mathbf{X}_{i})$$

of unital *-algebras.

 \underline{PROOF} The definition of $\Psi_{_{\mathbf{F}}}$ is the obvious one:

$$\Psi_{\mathbf{F}}(\alpha) = \prod_{\mathbf{i}\in\mathbf{I}_{\underline{C}}} \alpha_{\mathbf{X}_{\mathbf{i}}}.$$

 $\Psi_{\mathbf{F}}$ is injective:

$$\alpha_{X_{i}} = 0 \forall i \in I_{\underline{C}} \Rightarrow \alpha_{X} = 0 \forall X \in Ob \underline{C}.$$

To see this, choose the $s_{\underline{i}k}\in Mor\left(X_{\underline{i}},X\right)$ as in the proof of 7.8 -- then

$$\alpha_{\mathbf{X}} = \alpha_{\mathbf{X}} \circ \operatorname{Fid}_{\mathbf{X}}$$
$$= \sum_{\mathbf{i}\mathbf{k}} \alpha_{\mathbf{X}} \circ \operatorname{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}}^{*})$$
$$= \sum_{\mathbf{i}\mathbf{k}} \alpha_{\mathbf{X}} \circ \operatorname{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}) \circ \operatorname{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}^{*})$$

But the diagram



commutes, hence

$$\alpha_{\mathbf{X}} = \sum_{\mathbf{i}k} \mathbf{F}(\mathbf{s}_{\mathbf{i}k}) \circ \alpha_{\mathbf{X}} \circ \mathbf{F}(\mathbf{s}_{\mathbf{i}k}^{*})$$

= 0.

 $\frac{\Psi_{F}}{F}$ is surjective:

$$\forall \{\alpha_{i} \in Mor(FX_{i}, FX_{i}) : i \in I_{\underline{C}}\}, \exists \alpha \in Nat(F, F) : \Psi_{F}(\alpha) = \prod_{i \in I_{\underline{C}}} \alpha_{i}$$

Thus define $\boldsymbol{\alpha}_{X} \in \operatorname{Mor}\left(\mathsf{FX},\mathsf{FX}\right)$ by

$$\alpha_{X} = \sum_{ik} F(s_{ik}) \circ \alpha_{i} \circ F(s_{ik}^{*})$$

and define $\boldsymbol{\alpha}_{\boldsymbol{Y}} \in Mor(F\boldsymbol{Y},F\boldsymbol{Y})$ by

$$\alpha_{\mathbf{Y}} = \sum_{j\ell} \mathbf{F}(\mathbf{t}_{j\ell}) \circ \alpha_{j} \circ \mathbf{F}(\mathbf{t}_{j\ell}^{*}).$$

Then $\forall f \in Mor(X,Y)$,

$$\begin{aligned} \mathrm{Ff} \circ \alpha_{\mathbf{X}} &= \sum_{\mathbf{i}\mathbf{k}} \mathrm{F}(\mathbf{f} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}}) \circ \alpha_{\mathbf{i}} \circ \mathrm{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}^{*}) \\ &= \sum_{\mathbf{i}\mathbf{k}, \mathbf{j}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}} \circ \mathbf{t}_{\mathbf{j}\mathbf{\ell}}^{*} \circ \mathbf{f} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}}) \circ \alpha_{\mathbf{i}} \circ \mathrm{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}^{*}) \\ &= \sum_{\mathbf{i}\mathbf{k}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{i}\mathbf{\ell}} \circ (\mathbf{t}_{\mathbf{i}\mathbf{\ell}}^{*} \circ \mathbf{f} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}})) \circ \alpha_{\mathbf{i}} \circ \mathrm{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}^{*}) \\ &= \sum_{\mathbf{i}\mathbf{k}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{i}\mathbf{\ell}}) \circ \mathrm{F}(\mathbf{t}_{\mathbf{i}\mathbf{\ell}}^{*} \circ \mathbf{f} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}}) \circ \alpha_{\mathbf{i}} \circ \mathrm{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}^{*}) \\ &= \sum_{\mathbf{i}\mathbf{k}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{i}\mathbf{\ell}}) \circ \alpha_{\mathbf{i}} \circ \mathrm{F}(\mathbf{t}_{\mathbf{i}\mathbf{\ell}}^{*} \circ \mathbf{f} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}}) \circ \mathrm{F}(\mathbf{s}_{\mathbf{i}\mathbf{k}}^{*}) \\ &= \sum_{\mathbf{i}\mathbf{k}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}}) \circ \alpha_{\mathbf{j}} \circ \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}}^{*} \circ \mathbf{f} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}} \circ \mathbf{s}_{\mathbf{i}\mathbf{k}}^{*}) \\ &= \sum_{\mathbf{i}\mathbf{k}, \mathbf{j}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}}) \circ \alpha_{\mathbf{j}} \circ \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}}^{*} \circ \mathbf{f}) \\ &= \sum_{\mathbf{j}\mathbf{\ell}} \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}}) \circ \alpha_{\mathbf{j}} \circ \mathrm{F}(\mathbf{t}_{\mathbf{j}\mathbf{\ell}}^{*} \circ \mathbf{f}) \\ &= \alpha_{\mathbf{y}} \circ \mathrm{F}\mathbf{f}. \end{aligned}$$

Accordingly, the diagram



commutes, meaning that $\alpha \in Nat(F,F)$. And, by construction, $\alpha_{X_{\underline{i}}} = \alpha_{\underline{i}}$, so

$$\Psi_{\mathbf{F}}(\alpha) = \prod_{\mathbf{i}\in\mathbf{I}_{\mathbf{C}}} \alpha_{\mathbf{i}}.$$

8.4 <u>EXAMPLE</u> Take $\underline{C}' = \underline{C}$ and let $F = id_{\underline{C}}$ (the identity functor) — then

$$\operatorname{Nat}(\operatorname{id}_{\underline{C}},\operatorname{id}_{\underline{C}}) \approx \prod_{\underline{i} \in \underline{I}_{\underline{C}}} C.$$

8.5 <u>EXAMPLE</u> Suppose that <u>C</u> is a semisimple monoidal *-category -- then <u>C</u> × <u>C</u> is a semisimple *-category with

$$\mathbf{I}_{\underline{\mathbf{C}}} \times \underline{\mathbf{C}} = \mathbf{I}_{\underline{\mathbf{C}}} \times \mathbf{I}_{\underline{\mathbf{C}}}.$$

And

$$x_{i} \otimes x_{j} \approx \bigoplus_{k \in I_{\underline{C}}} N_{ij}^{k} x_{k'}$$

where

$$N_{ij}^{k} = \dim Mor(X_{k}, X_{i} \otimes X_{j}),$$

SO

$$\operatorname{Mor}(X_{i} \otimes X_{j}, X_{i} \otimes X_{j}) \approx \bigoplus_{k \in I_{\underline{C}}} \operatorname{M}_{ij}^{k} (C).$$

This said, let

 $\mathbf{F} = \mathbf{Q}: \mathbf{C} \times \mathbf{C} \to \mathbf{C}.$

Then

$$\begin{aligned} \mathsf{Nat}(\mathfrak{Q},\mathfrak{Q}) &\approx \prod_{i,j \in \mathbf{I}_{C}} \mathsf{Mor}(\mathfrak{Q}(\mathsf{X}_{i},\mathsf{X}_{j}), \mathfrak{Q}(\mathsf{X}_{i},\mathsf{X}_{j})) \end{aligned}$$

$$\approx \prod_{i,j \in I_{\underline{C}}} \operatorname{Mor}(X_{i} \otimes X_{j}, X_{i} \otimes X_{j})$$
$$\approx \prod_{i,j,k \in I_{\underline{C}}} \operatorname{M}_{k}(C).$$
$$\underset{N_{ij}^{k} \neq 0}{\operatorname{N}_{ij}^{k} \neq 0}$$

Suppose that C is a semisimple *-category, let $F:C \rightarrow \underline{FDHILB}$ be *-preserving and put

$$A_{\mathbf{F}}^{\prime} = \bigoplus_{\mathbf{i} \in \mathbf{I}_{\underline{C}}} \mathbf{B}(\mathbf{F}\mathbf{X}_{\mathbf{i}})$$

which, of course, can be embedded in

$$\prod_{i \in I_C} B(FX_i) \ (\approx Nat(F,F)).$$

Needless to say, A_F is a *-algebra, unital iff $I_{\underline{C}}$ is finite. The projections $p_i:A_F \rightarrow B(FX_i)$ are finite dimensional irreducible *-representations. Moreover, any finite dimensional nondegenerate *-representation of A_F is a direct sum of finite dimensional irreducible *-representations and every finite dimensional irreducible *-representation is unitarily equivalent to a p_i .

Define now a *-preserving functor

$$\phi: \underline{C} \to \underline{\operatorname{Rep}}_{fd} A_F$$

as in 8.2 -- then Φ is an equivalence of categories iff F is faithful. In fact, since Φ and F agree on morphisms, it is clear that

Assume therefore that F is faithful. From the definitions, $\pi_{X_i} = p_i$ (or still, $\forall \alpha \in A, \alpha_{X_i} = p_i(\alpha)$), which is a finite dimensional irreducible *-representation of A_F . Given an irreducible $X \in Ob \subseteq$, $\exists i \in I_{\underline{C}}$ and an isomorphism $\phi_i: X_i \to X$. Since the diagram



commutes, π_X is also a finite dimensional irreducible *-representation of A_F . If i \neq j, then

$$Mor(p_{i}, p_{j}) = \{0\},\$$

so if $X, Y \in Ob \subseteq$ are irreducible and nonzero, then

$$Mor(\pi_{X'}\pi_{Y}) = \{0\}.$$

Because Φ is faithful (and $\underline{\operatorname{Rep}}_{\operatorname{fd}} \stackrel{A_{\operatorname{F}}}{\operatorname{F}}$ is a semisimple *-category), the foregoing considerations imply that Φ is full (cf. 7.9). Finally, Φ has a representative image. Indeed, as mentioned above, every finite dimensional irreducible *-representation of A_{F} is unitarily equivalent to a p_i .

To recapitulate:

8.6 <u>THEOREM</u> Let <u>C</u> be a semisimple *-category and let $F: \underline{C} \rightarrow \underline{FDHILB}$ be a *-preserving functor. Put

$$A_{\mathbf{F}} = \bigoplus_{i \in \mathbf{I}_{\underline{C}}} B(FX_i)$$

and define

 $\Phi: \underline{C} \to \underline{\operatorname{Rep}}_{\mathrm{fd}} \overset{A}{F}$

by

$$\Phi X = (\pi_X, FX) \quad (X \in Ob C)$$

$$\Phi f = Ff \quad (f \in Mor(X, Y)).$$

Then Φ is an equivalence of categories iff F is faithful.

Let \underline{C} be a semisimple strict monoidal *-category. Definition: An <u>embedding functor</u> (for \underline{C}) is a faithful unitary functor

 $F:C \rightarrow FDHILB.$

[Note: Recall from §7 that in this context, "unitary" means that F is a *-preserving monoidal functor for which the isomorphisms $\xi: e \rightarrow Fe$ and

$$\Xi_{X,Y}: FX \ \underline{\textcircled{O}} \ FY \ \rightarrow \ F(X \ \underline{\textcircled{O}} \ Y)$$

are unitary (\underline{e} = standard unit in <u>FDHILB</u>, \underline{o} = strict monoidal structure in <u>FDHILB</u> (cf. 7.5)).]

8.7 LEMMA There is an isomorphism

$$\Psi_{\mathbf{F}}: \mathsf{Nat}(\mathbf{F}, \mathbf{F}) \rightarrow \prod_{i \in \mathbf{I}} \mathsf{B}(\mathsf{FX}_{i})$$

of unital *-algebras (cf. 8.3).

8.8 LEMMA The map

$$\varepsilon_{\mathbf{F}}: \operatorname{Nat}(\mathbf{F}, \mathbf{F}) \rightarrow \operatorname{Mor}(\mathbf{F}\mathbf{e}, \mathbf{F}\mathbf{e}) \approx \mathbb{C}$$

that sends

$$\alpha = \{\alpha_X\} \text{ to } \alpha_e$$

is a unital *-homomorphism.

8.9 SCHOLIUM The map

$$\tilde{\varepsilon}_{F}: \prod_{i \in I_{C}} B(FX_{i}) \rightarrow C$$

that sends

T to
$$\varepsilon_{\mathbf{F}} \circ \Psi_{\mathbf{F}}^{-1}(\mathbf{T})$$

is a unital *-homomorphism.

Let

$$\varepsilon = \overline{\varepsilon}_{\mathbf{F}} | \mathbf{A}_{\mathbf{F}}.$$

Then ε is a unital *-homomorphism, the counit.

8.10 LEMMA There is an isomorphism

$${}^{\Psi}_{\mathbf{F}} \circ \mathfrak{A}^{:\operatorname{Nat}(\mathbf{F}} \circ \mathfrak{A}, \mathbf{F} \circ \mathfrak{A}) \longrightarrow \prod_{i, j \in \mathbf{I}_{\underline{C}}} {}^{\operatorname{B}(\mathbf{F}X_{i})} \mathfrak{A}_{\mathsf{C}} {}^{\operatorname{B}(\mathbf{F}X_{j})}$$

of unital *-algebras.

PROOF In fact,

$$\approx \prod_{i,j \in I_{\underline{C}}} B(F(X_{i} \otimes X_{j})) \quad (cf. 8.3)$$
$$\approx \prod_{i,j \in I_{\underline{C}}} B(FX_{i} \otimes' FX_{j})$$
$$i, j \in I_{\underline{C}}$$

$$\prod_{i,j\in I} B(FX_i) \otimes_{C} B(FX_j).$$

8.11 LEMMA The map

$$\Delta_{\mathbf{F}}^{\mathsf{:Nat}(\mathbf{F},\mathbf{F})} \rightarrow \mathsf{Nat}(\mathbf{F} \circ \mathbf{Q},\mathbf{F} \circ \mathbf{Q})$$

that sends

$$\alpha = \{\alpha_X\} \text{ to } \{\alpha_X \otimes Y\}$$

is a unital *-homomorphism.

8.12 SCHOLIUM The map

$$\overline{\Delta}_{F}: \prod_{i \in I_{\underline{C}}} B(FX_{\underline{i}}) \rightarrow \prod_{i, j \in I_{\underline{C}}} B(FX_{\underline{i}}) \otimes_{\underline{C}} B(FX_{\underline{j}})$$

that sends

$$T \operatorname{to} \Psi_{\mathbf{F}} \circ \mathfrak{Q} \circ \vartriangle_{\mathbf{F}} \circ \Psi_{\mathbf{F}}^{-1} \langle \mathbf{T} \rangle$$

is a unital *-homomorphism.

 \mathbf{Iet}

$$\Delta = \overline{\Delta}_{\mathbf{F}} | \mathbf{A}_{\mathbf{F}}.$$

Then \triangle is a unital *-homomorphism, the <u>coproduct</u>.

Let

$$\pi_{1}: A_{\mathbf{F}} \rightarrow B(H_{1})$$
$$\pi_{2}: A_{\mathbf{F}} \rightarrow B(H_{2})$$

be nondegenerate *-representations of $\boldsymbol{A}_{\mathbf{F}}$ on finite dimensional Hilbert spaces

10.

 $\begin{bmatrix} H_1 \\ (\text{the zero representation is a possibility}) -- \text{ then we can form} \\ H_2 \end{bmatrix}$

$$\pi_1 \ \underline{\otimes} \ \pi_2 : A_F \ \underline{\otimes}_C \ A_F \ \Rightarrow \ B(H_1) \ \underline{\otimes}_C \ B(H_2) \ \approx \ B(H_1 \ \underline{\otimes} \ H_2).$$

Since

$$A_{F} \stackrel{\text{a}_{C}}{=} f^{F} \stackrel{\text{a}_{C}}{=} f^{F} \stackrel{\text{b}_{C}}{=} f^{F} \stackrel$$

it follows that $\pi_1 \otimes \pi_2$ admits a unique extension to a unital *-homomorphism

$$\overline{\pi_1 \otimes \pi_2}: \prod_{i,j \in \mathbf{I}_{\underline{C}}} \mathbb{B}(FX_i) \otimes_{\mathbf{C}} \mathbb{B}(FX_j) \rightarrow \mathbb{B}(H_1 \otimes H_2).$$

This being so, put

$$\pi_1 \times \pi_2 = \overline{\pi_1 \otimes \pi_2} \circ \Delta.$$

Then $\pi_1 \times \pi_2$ is a nondegenerate *-representation of A_F on the finite dimensional Hilbert space $H_1 \ \underline{a} \ H_2$.

8.13 <u>LEMMA</u> The data (×, ε ,...) is a monoidal structure on <u>Rep_{fd}</u> A_F.

Therefore $\underline{\text{Rep}}_{fd} \stackrel{A}{}_{F}$ is a semisimple monoidal *-category (the counit ε is the irreducible unit).

8.14 THEOREM Let C be a semisimple strict monoidal *-category and let

$$F:\underline{C} \rightarrow \underline{FDHILB}$$

be an embedding functor. Put

$$A_{F} = \bigoplus_{i \in I_{C}} B(FX_{i})$$

and define

$$\Phi: \underline{C} \rightarrow \underline{\operatorname{Rep}}_{fd} A_F$$

by

$$\Phi X = (\pi_X, FX) \quad (X \in Ob \underline{C})$$
$$\Phi f = Ff \quad (f \in Mor(X, Y)).$$

Then Φ is a monoidal equivalence.

<u>PROOF</u> By hypothesis, F is faithful, hence Φ is an equivalence of categories (cf. 8.6). So, in view of 2.8, it suffices to show that Φ is monoidal. There are two points. First

$$\Phi e = (\pi_e, Fe)$$

and $\forall \alpha \in A_{\mathbf{F}}$, the diagram



commutes, i.e., ξ intertwines ε and $\pi_e.$ Next, given $X,Y\in Ob\ \underline{C},$ consider

$$\Phi X \times \Phi Y = (\pi_X \times \pi_Y, FX \oplus FY)$$
$$\Phi (X \oplus Y) = (\pi_X \oplus Y', F(X \oplus Y)).$$

Then

$$E_{X,Y}:FX \ \underline{\otimes} \ FY \ \neq \ F(X \ \underline{\otimes} \ Y)$$

is an intertwining operator: $\forall \ \alpha \in A_{\mathbf{F}}$,

$$\Xi_{\mathbf{X},\mathbf{Y}} \circ (\pi_{\mathbf{X}} \times \pi_{\mathbf{Y}}) (\alpha) = \pi_{\mathbf{X} \otimes \mathbf{Y}} (\alpha) \circ \Xi_{\mathbf{X},\mathbf{Y}}.$$

The interchange $\sigma: A_F \otimes_C A_F \to A_F \otimes_C A_F$ ($\sigma(\alpha \otimes \beta) = \beta \otimes \alpha$) is a nondegenerate *-homomorphism, thus has a unique extension to an involutive *-automorphism

$$\vec{\sigma}: \prod_{i,j \in I} B(FX_i) \otimes_{C} B(FX_j) \rightarrow \prod_{i,j \in I} B(FX_i) \otimes_{C} B(FX_j).$$

Let

 $\Delta^{\rm OP} = \overline{\sigma} \circ \Delta.$

Then A_F is said to be <u>cocommutative</u> if $\Delta = \Delta^{OP}$.

8.15 LEMMA Suppose that $A_{\rm F}$ is cocommutative -- then

$$\forall \begin{vmatrix} (\pi_1, H_1) \\ & \in \text{Ob } \underline{\text{Rep}}_{\text{fd}} A_{\text{F}}, \\ (\pi_2, H_2) \end{vmatrix}$$

the diagram

$$\begin{array}{c} H_{1} \underline{\otimes} H_{2} & \xrightarrow{H_{1}, H_{2}} & H_{2} \underline{\otimes} H_{1} \\ \pi_{1} \times \pi_{2} \downarrow & \downarrow & & & & \downarrow \\ H_{1} \underline{\otimes} H_{2} & \xrightarrow{T_{H_{1}, H_{2}}} & H_{2} \underline{\otimes} H_{1} \end{array}$$

commutes.

PROOF Abbreviate
$$T_{H_1,H_2}$$
 to $T_{1,2}$ and note that

$$\forall \ \mathtt{T} \in \prod_{i, j \in \mathtt{I}_{\underline{C}}} \mathtt{B}(\mathtt{FX}_{i}) \ \mathtt{B}_{\mathtt{C}} \ \mathtt{B}(\mathtt{FX}_{j}),$$

we have

$$\overline{(\pi_1 \ \underline{\otimes} \ \pi_2)} \overline{\sigma} (\mathbf{T}) = \tau_{2,1} \overline{(\pi_2 \ \underline{\otimes} \ \pi_1)} (\mathbf{T}) \tau_{1,2}.$$

So, $\forall \alpha \in A_{\mathbf{F}}$,

$$\tau_{1,2}(\pi_{1} \times \pi_{2})(\alpha)$$

$$= \tau_{1,2}(\overline{\pi_{1} \otimes \pi_{2}})(\Delta(\alpha))$$

$$= \tau_{1,2}(\overline{\pi_{1} \otimes \pi_{2}})(\Delta^{OP}(\alpha))$$

$$= \tau_{1,2}(\overline{\pi_{1} \otimes \pi_{2}})(\overline{\sigma}(\Delta(\alpha)))$$

$$= \tau_{1,2}^{T}(\overline{\pi_{2} \otimes \pi_{1}})(\Delta(\alpha))\tau_{1,2}$$

$$= (\pi_{2} \times \pi_{1})(\alpha)\tau_{1,2}.$$

Thus, if A_F is cocommutative, then $\underline{\operatorname{Rep}}_{fd} \ A_F$ is a semisimple symmetric monoidal *-category.

8.16 <u>REMARK</u> Assume further that the category <u>C</u> is symmetric and that the embedding functor

$$F:\underline{C} \rightarrow \underline{FDHILB}$$

is symmetric monoidal -- then A_F is cocommutative and $\Phi: \underline{C} \to \underline{\text{Rep}}_{fd} A_F$ is a symmetric monoidal equivalence.

§9. CONJUGATES

Suppose that <u>C</u> is a strict monoidal *-category which is left autonomous. Put $X^{\sf V}$ = ${}^{\sf V}X$ -- then

$$\begin{bmatrix} - \varepsilon_X : & X \otimes X \rightarrow e \\ & n_X : e \rightarrow X \otimes & X \end{bmatrix}$$

≈>

$$\begin{bmatrix} \varepsilon_X^*: e \to V_X \otimes X \\ X^*: e \to V_X \otimes X \\ \varepsilon_X^*: X \otimes V_X \to e \end{bmatrix}$$

=>

$$\begin{array}{c} & \Pi_X^*: X \ \otimes \ X^{\vee} \ \to \ e \\ & & \\ & \varepsilon_X^*: e \ \to \ X^{\vee} \ \otimes \ X. \end{array}$$

And

=>

$$\begin{array}{c} (n_X^* \otimes \operatorname{id}_X) \circ (\operatorname{id}_X \otimes \varepsilon_X^*) = \operatorname{id}_X \\ (\operatorname{id}_X \otimes n_X^*) \circ (\varepsilon_X^* \otimes \operatorname{id}_Y) = \operatorname{id}_Y \\ X^{\vee} & X^{\vee} & X^{\vee} \end{array}$$

I.e.: The left duality $({}^{v}X,\epsilon_{X},n_{X})$ automatically leads to a right duality $(X^{v},n_{X}^{\star},\epsilon_{X}^{\star})$.

$$(x^{\vee}, \varepsilon_{X} \circ \tau, v_{X}, x, v_{X} \circ \eta_{X}).$$

9.1 COHERENCY HYPOTHESIS $\forall x \in Ob C$,

$$\varepsilon_X^* = \tau \circ \eta_X^*$$

[Note: The asymmetry is only apparent. For

$$n_{X} = \tau^{-1} \circ \epsilon_{X}^{*}$$
$$= \tau_{v_{X,X}} \circ \epsilon_{X}^{*}$$

=>

$$\eta_{X}^{*} = \varepsilon_{X} \circ \tau_{V_{X,X}}^{*}$$
$$= \varepsilon_{X} \circ \tau_{V_{X,X}}^{-1}$$
$$= \varepsilon_{X} \circ \tau_{X,X}^{-1}$$

In the presence of 9.1, let

$$\overline{\mathbf{x}} = \mathbf{v} \mathbf{x} (= \mathbf{x}^{\mathbf{v}})$$
$$\mathbf{r}_{\mathbf{x}} = \mathbf{\varepsilon}_{\mathbf{x}}^{\mathbf{\star}}$$
$$\overline{\mathbf{r}}_{\mathbf{x}} = \mathbf{\tau}_{\mathbf{x}} \circ \mathbf{r}_{\mathbf{x}'}$$

thus

$$\begin{bmatrix} \mathbf{r}_{\mathbf{X}} : \mathbf{e} \to \mathbf{\bar{X}} & \mathbf{a} \\ \mathbf{\bar{r}}_{\mathbf{X}} : \mathbf{e} \to \mathbf{X} & \mathbf{a} \\ \mathbf{\bar{X}} \end{bmatrix}$$

anđ

$$(\bar{X}, r_{X}^{*}, \bar{r}_{X})$$
 is a left duality
 $(\bar{X}, \bar{r}_{X}^{*}, r_{X})$ is a right duality.

Therefore

$$\begin{bmatrix} (\mathrm{id}_X \otimes \mathbf{r}_X^*) \circ (\bar{\mathbf{r}}_X \otimes \mathrm{id}_X) = \mathrm{id}_X \\ (\mathbf{r}_X^* \otimes \mathrm{id}_X) \circ (\mathrm{id}_X \otimes \bar{\mathbf{r}}_X) = \mathrm{id}_X \\ \bar{\mathbf{X}} & \bar{\mathbf{X}} & \bar{\mathbf{X}} & \bar{\mathbf{X}} \end{bmatrix} = \begin{bmatrix} (\bar{\mathbf{r}}_X^* \otimes \mathrm{id}_X) \circ (\mathrm{id}_X \otimes \mathbf{r}_X) = \mathrm{id}_X \\ (\mathrm{id}_X \otimes \bar{\mathbf{r}}_X) \circ (\mathrm{id}_X \otimes \mathbf{r}_X) = \mathrm{id}_X \\ \bar{\mathbf{X}} & \bar{\mathbf{X}} & \bar{\mathbf{X}} & \bar{\mathbf{X}} \end{bmatrix}$$

The relations

$$\begin{array}{c} (\operatorname{id}_{X} \otimes \mathbf{r}_{X}^{*}) \circ (\overline{\mathbf{r}}_{X} \otimes \operatorname{id}_{X}) = \operatorname{id}_{X} \\ (\operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_{X}^{*}) \circ (\mathbf{r}_{X} \otimes \operatorname{id}_{\overline{X}}) = \operatorname{id}_{\overline{X}} \\ - & \overline{X} & \overline{X} & \overline{X} \end{array}$$

are called the <u>conjugate equations</u>, the triple $(\bar{X}, r_{\bar{X}}, \bar{r}_{\bar{X}})$ being a <u>conjugate</u> for X. <u>N.B.</u> The conjugate equations imply that

$$(\overline{r}_{X}^{*} \otimes id_{X}) \circ (id_{X} \otimes r_{X}) = id_{X}$$
$$(r_{X}^{*} \otimes id_{X}) \circ (id_{\overline{X}} \otimes \overline{r}_{X}) = id_{\overline{X}}.$$

Having made these points, matters can be turned around. So start with a symmetric strict monoidal *-category \underline{C} -- then \underline{C} has conjugates if one can assign to each $X \in Ob \ \underline{C}$ an object \overline{X} and a morphism

such that the triple $(\bar{x}, r_X, \bar{r}_X)$ satisfies the conjugate equations (here, of course, $\bar{r}_X = \tau \circ r_X$).

E.g.: FDHILB has conjugates.

9.2 <u>REMARK</u> If <u>C</u> has conjugates, then <u>C</u> is left autonomous (consider $(\bar{X}, r_X^*, \bar{r}_X)$) and right autonomous (consider $(\bar{X}, \bar{r}_X^*, r_X)$). Moreover, the coherency hypothesis is in force: $(r_X^*)^* = r_X^*$, while

$$\overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}, \overline{\mathbf{x}}$$

9.3 LEMMA Suppose that C has conjugates.

• Under the identification

$$Mor(X \otimes Y, Z) \approx Mor(Y, \overline{X} \otimes Z),$$

the arrows

$$\begin{array}{c} \overline{f} \rightarrow \operatorname{id}_{X} & \otimes f \circ r_{X} & \operatorname{id}_{Y} \\ g \rightarrow \overline{r}_{X}^{*} & \operatorname{id}_{Z} & \operatorname{id}_{X} & \operatorname{id}_{Y} \end{array}$$

are mutually inverse.

• Under the identification

 $Mor(Y \otimes X, Z) \approx Mor(Y, Z \otimes \overline{X}),$

the arrows

$$\begin{array}{c} f \neq f \otimes id \circ id_{Y} \otimes \overline{r}_{X} \\ \overline{X} \\ g \neq id_{Z} \otimes r_{X}^{*} \circ g \otimes id_{X} \end{array}$$

are mutually inverse.

E.g.:
$$\forall X \in Ob \underline{C}$$
,

Mor(X,X)
$$\approx$$
 Mor(e, $\overline{X} \otimes X$).

9.4 LEMMA If

$$= (\bar{\mathbf{x}}, \mathbf{r}_{\mathbf{X}}, \bar{\mathbf{r}}_{\mathbf{X}})$$
$$= (\bar{\mathbf{x}}, \mathbf{r}_{\mathbf{X}}, \bar{\mathbf{r}}_{\mathbf{X}}, \bar{\mathbf{r}}_{\mathbf{X}})$$

are conjugates for X, then

$$\mathbf{r}_{X}^{\star} \overset{\text{a}}{=} \operatorname{id}_{\overline{X}'} \circ \operatorname{id}_{\overline{X}} \overset{\text{a}}{=} \mathbf{r}_{X}^{\prime} \in \operatorname{Mor}(\overline{X}, \overline{X}^{\prime})$$

is unitary.

PROOF Put

$$U = r_X^* \otimes id \otimes id \otimes \bar{r}'_X$$
$$(= id \otimes \bar{r}_X^* \circ r_X' \otimes id \dots).$$

Then the claim is that

$$\begin{bmatrix} U \circ U^* = id \\ X^* \end{bmatrix}$$
$$U^* \circ U = id .$$
$$X$$

And for this, it will be enough to consider U \circ U*. So write

U

•
$$\mathbf{U}^* = \mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^* \circ \mathbf{U}^*$$

= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ (\operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^* \circ \mathbf{U}^* \otimes \operatorname{id}_{\mathbf{e}})$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ \mathbf{U}^* \otimes \overline{\mathbf{r}}_X^*$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ (\mathbf{U}^* \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*)$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ (\mathbf{U}^* \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*)$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*})$
 $\circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \circ \overline{\mathbf{r}}_X \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*} \circ \operatorname{id}_{\overline{X} \otimes \overline{X}^*})$
 $\circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*} \circ \operatorname{id}_{\overline{X} \otimes \overline{X}^*})$
 $\circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*} \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*})$
 $\circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*} \circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*})$
 $\circ \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X^*$
= $\mathbf{r}_X^* \otimes \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}} \otimes \overline{\mathbf{r}}_X \otimes \operatorname{id}_{\overline{X} \otimes \overline{X}^*}$

$$\circ \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X} \, \otimes \, \operatorname{id}_{X \, \otimes \, \overline{X}}, \, \circ \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \\ = \operatorname{id}_{e} \, \otimes \, (r_{X}^{*} \, \otimes \, \operatorname{id}_{\overline{X}}) \, \circ \, r_{X}'^{*} \, \otimes \, \operatorname{id}_{\overline{X} \, \otimes \, X \, \otimes \, \overline{X}}, \\ \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X} \, \otimes \, \operatorname{id}_{X \, \otimes \, \overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, r_{X}^{*} \, \otimes \, \operatorname{id}_{\overline{X}}, \\ \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X} \, \otimes \, \operatorname{id}_{\overline{X}} \, \otimes \, x_{X} \, \otimes \, \overline{x}, \, \circ \, \operatorname{id}_{\overline{X}} \, \otimes \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}} \, \otimes \, x_{X} \, \otimes \, x_{X}^{*} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}} \, \otimes \, x_{X} \, \otimes \, x_{X}^{*} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \\ \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, \operatorname{id}_{\overline{X}}, \\ \circ \, \left(\operatorname{id}_{\overline{X}}, \, \otimes \, (\operatorname{id}_{X} \, \otimes \, r_{X}^{*}) \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, (\overline{r}_{X} \, \otimes \, \operatorname{id}_{X}) \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}}, \, \circ \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \otimes \, (\operatorname{id}_{\overline{X}} \, \otimes \, r_{X}^{*}) \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}}, \, \otimes \, (\overline{r}_{\overline{X}} \, \otimes \, \operatorname{id}_{\overline{X}}) \, \otimes \, \operatorname{id}_{\overline{X}}, \, \circ \, \operatorname{id}_{\overline{X}}, \, \circ \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \\ = r_{X}^{**} \, \otimes \, \operatorname{id}_{\overline{X}}, \, \otimes \, \overline{r}_{X}' \, \otimes \, \operatorname{id}_{\overline{X}}, \, \otimes \, \operatorname{$$

$$\circ (\operatorname{id}_{\overline{X}} \circ \operatorname{id}_{\overline{X}} \otimes ((\operatorname{id}_{\overline{X}} \otimes r_{\overline{X}}^{*}) \otimes \operatorname{id}_{\overline{X}}, \circ (\overline{r}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}}) \otimes \operatorname{id}_{\overline{X}}))$$

$$\circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{\overline{X}}^{*}$$

$$= r_{X}^{**} \otimes \operatorname{id}_{\overline{X}},$$

$$\circ (\operatorname{id}_{\overline{X}} \otimes ((\operatorname{id}_{\overline{X}} \otimes r_{\overline{X}}^{*}) \circ (\overline{r}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}})) \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}},$$

$$\circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{\overline{X}}^{*}$$

$$= r_{X}^{**} \otimes \operatorname{id}_{\overline{X}}, \circ (\operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}}) \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{\overline{X}}^{*}$$

$$= r_{X}^{**} \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{\overline{X}}^{*}$$

$$= r_{X}^{**} \otimes \operatorname{id}_{\overline{X}}, \circ \operatorname{id}_{\overline{X}} \otimes \overline{r}_{\overline{X}}^{*}$$

[Note: Evidently,

$$\begin{bmatrix} \mathbf{r}_{\mathbf{X}}^{\mathsf{I}} = (\mathbf{U} \otimes \mathrm{id}_{\mathbf{X}}) \circ \mathbf{r}_{\mathbf{X}} \\ \\ \mathbf{\bar{r}}_{\mathbf{X}}^{\mathsf{I}} = (\mathrm{id}_{\mathbf{X}} \otimes \mathbf{U}) \circ \mathbf{\bar{r}}_{\mathbf{X}^{\mathsf{I}}} \end{bmatrix}$$

Conjugates are therefore determined up to "unitary equivalence". Put

$$\Omega_{X} = r_{X}^{*} \otimes id_{X} \circ id_{X} \otimes \tau_{X,X} \circ r_{X} \otimes id_{X}.$$

Then

$$\Omega_{\mathbf{X}} \in Mor(\mathbf{X}, \mathbf{X})$$

is unitary and it can be verified by computation that the assignment $X + \Omega_X$ defines a twist Ω . This fact, however, is a trivial consequence of the following result.

$$\Omega_{\mathbf{X}} = \mathrm{id}_{\mathbf{X}}$$

PROOF We have

$$(\mathbf{\tilde{r}}_{\mathbf{X}}^{\star} \otimes \mathbf{id}_{\mathbf{X}}) \circ (\mathbf{id}_{\mathbf{X}} \otimes \mathbf{r}_{\mathbf{X}}) = \mathbf{id}_{\mathbf{X}}.$$

On the other hand, there is a commutative diagram

$$\begin{array}{c|c} x = e \otimes x & \xrightarrow{T_{e,X}} & x \otimes e = x \\ r_{x} \otimes id_{x} & \downarrow & & \downarrow id_{x} \otimes r_{x} \\ \hline \overline{x} \otimes x \otimes x & \xrightarrow{T_{x}} & x \otimes \overline{x} \otimes x, \end{array}$$

so

$$(\vec{r}_{X}^{*} \otimes id_{X}) \circ (id_{X} \otimes r_{X})$$

$$= \vec{r}_{X}^{*} \otimes id_{X} \circ \tau \circ r_{X} \otimes id_{X}$$

$$= \vec{r}_{X}^{*} \otimes id_{X} \circ \tau \otimes id_{X} \circ id_{X} \circ r_{X} \otimes id_{X}.$$

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$$\bar{r}_{X}^{*} \otimes id_{X} \circ \tau \otimes id_{X}$$

$$= (\tau \circ r_{X})^{*} \otimes id_{X} \circ \tau \otimes id_{X}$$

$$= r_{X}^{*} \circ \tau \otimes id_{X} \circ \tau \otimes id_{X}$$

$$= r_{X}^{*} \circ \tau \otimes id_{X} \circ \tau \otimes id_{X}$$

$$= r_{X}^{*} \circ \tau \otimes \tau \otimes id_{X} \circ id_{X}$$

$$= r_{X}^{*} \otimes id_{X}.$$

[Note: Therefore, in the terminology of §6, C is an even ribbon category.]

9.6 REMARK $\forall f \in Mor(X,X)$, the diagram

commutes.



10.

Maintaining the supposition that <u>C</u> has conjugates, recall that <u>C</u> is left autonomous with left duality $(\bar{X}, r_X^*, \bar{r}_X)$ (cf. 9.2), thus by definition the categorical dimension of X is the arrow

$$e \xrightarrow{\bar{r}_{X}} X \otimes \bar{X} \xrightarrow{T} X, \bar{X} \xrightarrow{T} X \otimes X \xrightarrow{T_{X}} e \quad (cf. \S6).$$

But

$$\bar{r}_{X} = \tau \circ r_{X'}$$

so the categorical dimension of X is the composition

$$\mathbf{r}_{\mathbf{X}}^{\star} \circ \mathbf{r}_{\mathbf{X}, \mathbf{X}}^{\star} \circ \mathbf{r}_{\mathbf{X}, \mathbf{X}}^{\star} \circ \mathbf{r}_{\mathbf{X}}^{\star}$$
$$= \mathbf{r}_{\mathbf{X}}^{\star} \circ \mathbf{r}_{\mathbf{X}} \in \operatorname{Mor}(\mathbf{e}, \mathbf{e})$$
$$\equiv \dim \mathbf{X}.$$

[Note: Since $\Omega = id$, $\forall f \in Mor(X,X)$,

$$tr_{X}(f) = r_{X}^{\star} \circ \tau \circ \Omega_{X} \otimes id \circ (f \otimes id) \circ \tilde{r}_{X}$$

$$= r_{X}^{\star} \circ \tau \circ id \circ (f \otimes id) \circ \tau \circ r_{X}$$

$$= r_{X}^{\star} \circ \tau \circ id \otimes \tilde{X} \otimes \tilde$$

N.B. dim X does not depend on the choice of a conjugate for X. Indeed, if $U:\bar{X} \, \div \, \bar{X}^{1}$ is unitary, then

$$((U \otimes id_X) \circ r_X)^* \circ ((U \otimes id_X) \circ r_X)$$
$$= r_X^* \circ U^* \otimes id_X \circ U \otimes id_X \circ r_X$$

$$= r_{X}^{*} \circ id_{X} \otimes id_{X} \circ r_{X}$$
$$= r_{X}^{*} \circ id_{X} \otimes x$$
$$= r_{X}^{*} \circ r_{X}.$$

9.7 LEMMA If

$$(\bar{X}, r_{X}, \bar{r}_{X}) \text{ is a conjugate for } X$$
$$(\bar{Y}, r_{Y}, \bar{r}_{Y}) \text{ is a conjugate for } Y,$$

then

is a conjugate for X @ Y, where

$$\vec{r}_{X \otimes Y} = \operatorname{id}_{\overline{Y}} \otimes r_{X} \otimes \operatorname{id}_{Y} \circ r_{Y}$$
$$\vec{r}_{X \otimes Y} = \operatorname{id}_{X} \otimes \vec{r}_{Y} \otimes \operatorname{id}_{\overline{X}} \circ \vec{r}_{X}.$$

[The proof that

$$\begin{bmatrix} (\mathrm{id}_{X \otimes Y} \otimes r_{X \otimes Y}^{*}) \circ (\overline{r}_{X \otimes Y} \otimes \mathrm{id}_{X \otimes Y}) = \mathrm{id}_{X \otimes Y} \\ (\mathrm{id}_{\overline{Y} \otimes \overline{X}} \otimes \overline{r}_{X \otimes Y}^{*}) \circ (r_{X \otimes Y} \otimes \mathrm{id}_{\overline{Y} \otimes \overline{X}}) = \mathrm{id}_{\overline{Y} \otimes \overline{X}} \\ \underline{\overline{Y} \otimes \overline{X}} \otimes \overline{r}_{X \otimes Y}^{*}) \circ (r_{X \otimes Y} \otimes \mathrm{id}_{\overline{Y} \otimes \overline{X}}) = \mathrm{id}_{\overline{Y} \otimes \overline{X}} \\ \end{bmatrix}$$

will be left to the reader but we shall provide the verification that

$$\overline{\mathbf{r}}_{\mathbf{X} \otimes \mathbf{Y}} = \overline{\mathbf{r}}_{\mathbf{\overline{Y}} \otimes \mathbf{\overline{X}}, \mathbf{X} \otimes \mathbf{Y}} \circ \mathbf{r}_{\mathbf{X} \otimes \mathbf{Y}}$$

Thus write

$$= \tau_{\overline{X}, X, Q, Y, Q, \overline{Y}} \circ \tau_{\overline{Y}, \overline{X}, Q, X, Q, Y} \circ r_{X, Q, Y} \circ r_{X, Q, Y} \circ r_{\overline{X}, \overline{X}, Q, X, Q, Y} \circ r_{\overline{Y}, \overline{X}, Q, X} \circ r_{X, Q, Y}$$

$$= id_{\overline{X}, Q, X} \circ \tau_{\overline{Y}, Y} \circ \tau_{\overline{Y}, \overline{X}, Q, X} \circ r_{X, Q, Y} \circ r_{\overline{Y}, \overline{Y}, \overline{Y}, \overline{X}, Q, X} \circ r_{\overline{Y}, Q, \overline{Y}} \circ r_{\overline{Y}, \overline{Y}, \overline{Y$$

Therefore

$$\overline{\overline{Y}} \otimes \overline{X}, X \otimes Y \xrightarrow{\circ} \overline{X} \otimes Y$$

$$= \overline{\overline{X}}, X \otimes Y \otimes \overline{\overline{Y}} \xrightarrow{\circ} \operatorname{id}_{\overline{X}} \otimes X \xrightarrow{\otimes} \overline{\overline{Y}} \xrightarrow{\circ} \overline{T}_{X}$$

$$= \overline{\overline{X}}, X \otimes Y \otimes \overline{\overline{Y}} \xrightarrow{\circ} \operatorname{id}_{\overline{X}} \otimes \operatorname{id}_{X} \otimes \overline{\overline{T}}_{Y} \xrightarrow{\circ} T_{X}$$

$$= \operatorname{id}_{X} \otimes \overline{\overline{T}}_{Y} \otimes \operatorname{id}_{\overline{X}} \xrightarrow{\circ} \overline{\overline{X}}, X \xrightarrow{\circ} T_{X}$$

13.

Then

$$= \operatorname{id}_X \, \mathfrak{Q} \, \overline{r}_Y \, \mathfrak{Q} \, \operatorname{id}_{\overline{X}} \circ \overline{r}_X = \overline{r}_X \, \mathfrak{Q} \, Y^{*}$$

For all $X, Y \in Ob \ \underline{C}$, the map

$$Mor(X,Y) \rightarrow Mor(\overline{Y},\overline{X})$$

that sends f to $^{\vee}f$ is a linear bijection.

N.B. Here, as will be recalled from \$5,

$${}^{\mathsf{v}} \mathbf{f} = \mathbf{r}_{\mathbf{Y}}^{\star} \otimes \operatorname{id}_{\mathbf{X}} \circ \operatorname{id}_{\mathbf{Y}} \otimes \mathbf{f} \otimes \operatorname{id}_{\mathbf{X}} \circ \operatorname{id}_{\mathbf{X}} \otimes \mathbf{\bar{r}}_{\mathbf{X}}.$$

Now put

 $f^{+} = (^{v}f) *,$

thus

$$f^{+} = id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ id_{\overline{Y}} \otimes f^{*} \otimes id_{\overline{X}} \circ r_{Y} \otimes id_{\overline{X}}$$

and

$$f^+ \in Mor(\tilde{X}, \tilde{Y})$$
.

Properties:

1.
$$id_X^+ = id_{\overline{X}};$$

2. $(f^+)^* = (f^*)^+;$
3. $(f \circ g)^+ = f^+ \circ g^+.$

9.8 LEMMA Given $f \in Mor(X,Y)$, we have

$$f^+ \otimes id_X \circ r_X = id_{\overline{Y}} \otimes f^* \circ r_Y.$$

PROOF Start with the LHS and write

$$f^{+} \otimes id_{X} \circ r_{X}$$

$$= (id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ id_{\overline{Y}} \otimes f^{*} \otimes id_{\overline{X}} \circ r_{Y} \otimes id_{\overline{X}}) \otimes id_{X} \circ r_{X}$$

$$= (id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ id_{\overline{Y}} \otimes f^{*} \otimes id_{\overline{X}} \circ r_{Y} \otimes id_{\overline{X}}) \otimes (id_{X} \circ id_{X} \circ id_{X}) \circ r_{X}$$

$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \otimes id_{X} \circ id_{\overline{Y}} \otimes f^{*} \otimes id_{\overline{X}} \otimes id_{X} \circ r_{Y} \otimes id_{\overline{X}} \otimes id_{X} \circ r_{X}.$$

But

$$r_{Y} \bigotimes id_{\overline{X}} \bigotimes id_{X} \circ r_{X}$$

$$= r_{Y} \bigotimes id_{\overline{X}} \otimes id_{X} \circ id_{e} \bigotimes r_{X}$$

$$= id_{\overline{Y}} \bigotimes r_{X} \circ r_{Y} \bigotimes id_{e}$$

$$= id_{\overline{Y}} \bigotimes id_{Y} \bigotimes r_{X} \circ r_{Y}$$

=>

$$id_{\overline{Y}} \otimes f^* \otimes id_{\overline{X}} \otimes id_{\overline{X}} \circ id_{\overline{Y}} \otimes id_{\overline{Y}} \otimes r_{\overline{X}} \circ r_{\overline{Y}}$$
$$= id_{\overline{Y}} \otimes id_{\overline{X}} \otimes r_{\overline{X}} \circ id_{\overline{Y}} \otimes f^* \otimes id_{\overline{e}} \circ r_{\overline{Y}}$$

=>

$$f^{+} \otimes id_{X} \circ r_{X}$$

$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \otimes id_{X} \circ id_{\overline{Y}} \otimes id_{X} \otimes r_{X} \circ id_{\overline{Y}} \otimes f^{*} \circ r_{Y}$$

$$= id_{\overline{Y}} \otimes (\overline{r}_{X}^{*} \otimes id_{X} \circ id_{X} \otimes r_{X} \circ f^{*}) \circ r_{Y}$$

$$= \operatorname{id}_{\overline{Y}} \otimes \operatorname{id}_{\overline{X}} \circ f^* \circ r_{\overline{Y}}$$
$$= \operatorname{id}_{\overline{Y}} \otimes f^* \circ r_{\overline{Y}}.$$

9.9 <u>REMARK</u> Suppose that $T \in Mor(\overline{X}, \overline{Y})$ satisfies the equation

$$\mathbf{T} \boldsymbol{\otimes} \operatorname{id}_{X} \circ \mathbf{r}_{X} = \operatorname{id}_{\overline{Y}} \boldsymbol{\otimes} f^{*} \circ \mathbf{r}_{Y}.$$

Then

 $\mathbf{T} = \mathbf{f}^+$.

Proof:

$$f^{+} = id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ id_{\overline{Y}} \otimes f^{*} \otimes id_{\overline{X}} \circ r_{Y} \otimes id_{\overline{X}}$$
$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ id_{\overline{Y}} \otimes f^{*} \circ r_{Y} \otimes id_{\overline{X}} \circ id_{\overline{X}}$$
$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{*} \circ T \otimes id_{X} \circ r_{X} \otimes id_{\overline{X}}.$$

On the other hand,

$$T = T \circ id_{\overline{X}}$$

$$= T \circ id_{\overline{X}} \otimes \overline{r}_{X}^{\star} \circ r_{X} \otimes id_{\overline{X}}$$

$$= T \otimes id_{e} \circ id_{\overline{X}} \otimes \overline{r}_{X}^{\star} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{\star} \circ T \otimes id_{X} \otimes \overline{x} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{\star} \circ T \otimes id_{X} \otimes id_{\overline{X}} \circ r_{X} \otimes id_{\overline{X}}$$

$$= id_{\overline{Y}} \otimes \overline{r}_{X}^{\star} \circ T \otimes id_{X} \otimes id_{\overline{X}} \circ r_{X} \otimes id_{\overline{X}}$$
$$= \operatorname{id}_{\widetilde{Y}} \otimes \widetilde{r}_{X}^{*} \circ T \otimes \operatorname{id}_{X} \circ r_{X} \otimes \operatorname{id}_{\widetilde{X}}.$$

[Note: It is thus a corollary that if

$$id_{\overline{Y}} \otimes f^* \circ r_{\overline{Y}} = 0,$$

then $f^+ = 0$, so

$$(f^{\vee})^{\star} = 0 \Longrightarrow (f^{\vee})^{\star \star} = 0 \Longrightarrow f^{\vee} = 0 \Longrightarrow f = 0.1$$

9.10 <u>SCHOLIUM</u> f^+ is the unique element of Mor (\bar{X}, \bar{Y}) such that

$$f^+ \otimes id_X \circ r_X = id_{\overline{Y}} \otimes f^* \circ r_Y$$

[Note: $^{\vee}f$ is the unique element of Mor($\overline{Y}, \overline{X}$) such that

$$\operatorname{id}_{Y} \otimes {}^{\vee} f \circ \overline{r}_{Y} = f \otimes \operatorname{id}_{\overline{X}} \circ \overline{r}_{X},$$

so f^+ is the unique element of Mor $(\overline{X},\overline{Y})$ such that

$$\vec{r}_{Y}^{*} \circ id_{Y} \otimes f^{+} = \vec{r}_{X}^{*} \circ f^{*} \otimes id_{.}$$

9.11 LEMMA Suppose that

$$F: C \rightarrow FDHILB$$

is symmetric and unitary. Given $X \in Ob \ \underline{C}, \ put$

$$\begin{bmatrix} \mathbf{r}_{FX} = (\widehat{\Xi})^{-1} \circ Fr_{X} \circ \xi \\ \overline{X}, X \\ \overline{r}_{FX} = \tau \circ \mathbf{r}_{FX} \\ FX & F\overline{X}, FX \end{bmatrix}$$

Then the triple $(F\bar{X}, r_{FX}, \bar{r}_{FX})$ is a conjugate for FX.

PROOF What we know is that

$$\begin{bmatrix} (\mathrm{id}_X \otimes \mathbf{r}_X^*) \circ (\bar{\mathbf{r}}_X \otimes \mathrm{id}_X) = \mathrm{id}_X \\ (\mathrm{id}_X \otimes \bar{\mathbf{r}}_X^*) \circ (\mathbf{r}_X \otimes \mathrm{id}_X) = \mathrm{id}_X \\ \bar{\mathbf{X}} & \bar{\mathbf{X}} & \bar{\mathbf{X}} \end{bmatrix}$$

hence

$$F(\operatorname{id}_X \otimes r_X^*) \circ F(\overline{r}_X \otimes \operatorname{id}_X) = \operatorname{id}_{FX}$$

$$F(\operatorname{id}_X \otimes \overline{r}_X^*) \circ F(r_X \otimes \operatorname{id}_X) = \operatorname{id}_{FX},$$

$$- \overline{X} \quad F(\overline{r}_X \otimes \operatorname{id}_X) = \operatorname{id}_{FX},$$

and what we want to prove is that

$$[\begin{array}{c} (\mathrm{id}_{\mathrm{FX}} \ \underline{\otimes} \ \mathbf{r}_{\mathrm{FX}}^{\star}) \circ (\overline{\mathbf{r}}_{\mathrm{FX}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}}) = \mathrm{id}_{\mathrm{FX}} \\ (\mathrm{id}_{\mathrm{FX}} \ \underline{\otimes} \ \overline{\mathbf{r}}_{\mathrm{FX}}^{\star}) \circ (\mathbf{r}_{\mathrm{FX}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}}) = \mathrm{id}_{\mathrm{FX}} \\ \\ - & \mathbf{F} \overline{\mathbf{x}} \ \underline{} \ \overline{\mathbf{r}}_{\mathrm{FX}}^{\star}) \circ (\mathbf{r}_{\mathrm{FX}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}}) = \mathrm{id}_{\mathrm{FX}}.$$

The LHS of the first of these is the composition

$$id_{FX} \stackrel{a}{=} \xi^{-1} \circ Fr_{X}^{*} \circ \Xi_{X,X}$$
$$\circ \tau \qquad \circ (\Xi_{X,X})^{-1} \circ Fr_{X} \circ \xi \stackrel{a}{=} id_{FX}'$$
$$F_{X,FX} \qquad X_{X}$$

F being unitary. Write

$$\begin{array}{c} \left(\Xi \right)^{-1} \circ \operatorname{Fr}_{X} \circ \xi \underline{\Theta} \operatorname{id}_{FX} \\ & \operatorname{FX}, \operatorname{FX} & \overline{X}, X \end{array}^{2} \circ \operatorname{Id}_{FX} \circ \operatorname{id}_{FX} \\ & = \ldots \underline{\Theta} \operatorname{id}_{FX} \circ \operatorname{id}_{FX} \circ \operatorname{id}_{FX} \circ \operatorname{id}_{FX} \\ & = \operatorname{FX}, \operatorname{FX} \underline{\Theta} \operatorname{id}_{FX} \circ (\Xi)^{-1} \underline{\Theta} \operatorname{id}_{FX} \circ \operatorname{Fr}_{X} \underline{\Theta} \operatorname{id}_{FX} \circ \xi \underline{\Theta} \operatorname{id}_{FX}. \end{array}$$



Taking into account the commutative diagrams

F(X & X) & FX -----

Ξ

we have

$$Fr_{X} \stackrel{a}{\cong} id_{FX} \circ \xi \stackrel{a}{\cong} id_{FX}$$

$$= (E)^{-1} \circ F(r_{X} \stackrel{a}{\cong} id_{X}) \circ E_{e,X} \circ \xi \stackrel{a}{\cong} id_{FX}$$

$$= (E)^{-1} \circ F(r_{X} \stackrel{a}{\cong} id_{X}) \circ id_{FX}.$$

This leaves

$$\begin{array}{c} T \\ F\overline{X}, FX \end{array} \overset{{}_{\bullet}}{\stackrel{}_{FX}} \circ \begin{array}{c} (\Xi \\ \overline{X}, X \end{array} \end{array} \right)^{-1} \overset{{}_{\bullet}}{\stackrel{}_{FX}} \circ \begin{array}{c} (\Xi \\ \overline{X} \\$$

 \longrightarrow F($\overline{X} \otimes X \otimes X$),

Next

$$F(\bar{r}_X \otimes id_X) = F(\tau \circ r_X \otimes id_X)$$
$$= F(\tau \otimes id_X) \circ F(r_X \otimes id_X).$$
$$\bar{X}_X$$

Since F is symmetric, there is a commutative diagram



$$= (\Xi_{X,\overline{X}})^{-1} \underbrace{\otimes} \operatorname{id}_{FX} \circ (\Xi_{X,\overline{X}})^{-1} \circ F(\tau_{\overline{X},X} \otimes \operatorname{id}_{X})$$

$$=>$$

$$\tau_{F\overline{X},FX} \underbrace{\otimes} \operatorname{id}_{FX} \circ (\Xi_{X,X})^{-1} \underbrace{\otimes} \operatorname{id}_{FX} \circ (\Xi_{\overline{X},X})^{-1} \circ F(r_{X} \otimes \operatorname{id}_{X}) \circ \operatorname{id}_{FX}$$

$$= (\Xi_{X,\overline{X}})^{-1} \underbrace{\otimes} \operatorname{id}_{FX} \circ (\Xi_{X,\overline{X}})^{-1} \circ F(\tau_{\overline{X},X} \otimes \operatorname{id}_{X}) \circ F(r_{X} \otimes \operatorname{id}_{X}) \circ \operatorname{id}_{FX}$$

$$= (\Xi_{X,\overline{X}})^{-1} \underbrace{\otimes} \operatorname{id}_{FX} \circ (\Xi_{X,\overline{X},X})^{-1} \circ F(\tau_{\overline{X},X} \otimes \operatorname{id}_{X}) \circ \operatorname{id}_{FX}$$

$$= (\Xi_{X,\overline{X}})^{-1} \underbrace{\otimes} \operatorname{id}_{FX} \circ (\Xi_{X,\overline{X},X})^{-1} \circ F(\overline{r}_{X} \otimes \operatorname{id}_{X}) \circ \operatorname{id}_{FX}.$$

Analogously,

$$id_{FX} \stackrel{Q}{=} \xi^{-1} \circ Fr_{X}^{*} \circ \Xi_{\overline{X},X}$$

= $id_{FX} \circ F(id_{X} \otimes r_{X}^{*}) \circ \Xi_{X,\overline{X}} \otimes X \circ id_{FX} \stackrel{Q}{=} \Xi_{\overline{X},X}$

So, in summary,

$$(\mathrm{id}_{\mathrm{FX}} \ \underline{\otimes} \ r_{\mathrm{FX}}^{*}) \circ (\bar{r}_{\mathrm{FX}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}})$$

$$= \mathrm{id}_{\mathrm{FX}} \circ \mathrm{F}(\mathrm{id}_{\mathrm{X}} \ \underline{\otimes} \ r_{\mathrm{X}}^{*})$$

$$\circ \ \underline{\mathrm{E}}_{\mathrm{X}, \overline{\mathrm{X}}} \circ \mathrm{id}_{\mathrm{FX}} \ \underline{\otimes} \ \underline{\mathrm{E}}_{\mathrm{X}, \mathrm{X}} \circ (\underline{\mathrm{E}}_{\mathrm{X}, \mathrm{X}})^{-1} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{FX}} \circ (\underline{\mathrm{E}}_{\mathrm{X}, \mathrm{X}})^{-1}$$

$$\circ \ \mathrm{F}(\bar{r}_{\mathrm{X}} \ \underline{\otimes} \ \mathrm{id}_{\mathrm{X}}) \circ \mathrm{id}_{\mathrm{FX}},$$

thus to finish, it need only be shown that

$$\begin{array}{c} \Xi & \circ \operatorname{id}_{\mathrm{FX}} \underline{\otimes} \Xi & \circ (\Xi)^{-1} \underline{\otimes} \operatorname{id}_{\mathrm{FX}} \circ (\Xi)^{-1} \\ \mathrm{X}, \overline{\mathrm{X}} \underline{\otimes} \mathrm{X} & \mathrm{X}, \mathrm{X} & \mathrm{X}, \mathrm{X} \\ \end{array}$$
$$= \operatorname{id}_{\mathrm{F}(\mathrm{X} \underline{\otimes} \mathrm{X} \underline{\otimes} \mathrm{X})} \cdot$$

This, however, follows from the commutative diagram



9.12 REMARK We have

$$\overline{\mathbf{r}}_{\mathbf{F}\mathbf{X}} = (\Xi)^{-1} \circ \overline{\mathbf{Fr}}_{\mathbf{X}} \circ \xi.$$

In fact, the RHS equals

$$\begin{array}{c} \left(\Xi\right)^{-1} \circ F_{\mathsf{T}} \circ F_{\mathsf{X}} \circ F_{\mathsf{X}} \\ \mathsf{X}, \tilde{\mathsf{X}} & \tilde{\mathsf{X}}, \mathsf{X} \end{array}$$

and there is a commutative diagram



\$10. TANNAKIAN CATEGORIES

Let <u>C</u> be a symmetric strict monoidal *-category which is essentially small \leftarrow then <u>C</u> is said to be tannakian if the following conditions are met:

 $T_1: \forall X, Y \in Ob \underline{C},$

dim Mor $(X, Y) < \infty$.

 T_2 : <u>C</u> has subobjects, direct sums, and conjugates.

- $T_3: \underline{C}$ has a zero object.
- T_4 : e is irreducible.

10.1 REMARK A tannakian category is necessarily semisimple, hence is abelian.

10.2 <u>EXAMPLE</u> Let <u>CPTGRP</u> be the category whose objects are the compact Hausdorff topological groups (in brief, the "compact groups") and whose morphisms are the continuous homomorphisms. Given an object G in this category, let <u>Rep</u> G be the category whose objects are the finite dimensional continuous unitary representations of G and whose morphisms are the intertwining operators — then <u>Rep</u> G is tannakian (define r and \vec{r} by

$$\vec{r} \lambda = \lambda \Sigma \vec{e}_{i} \boldsymbol{\Omega} \vec{e}_{i}$$

$$(\lambda \in C \ (=e)),$$

$$\vec{r} \lambda = \lambda \Sigma \vec{e}_{i} \boldsymbol{\Omega} \vec{e}_{i}$$

where $\{e_i\} \in {\it H} \mbox{ is an orthonormal basis for the representation space and <math display="inline">\{\bar{e}_i\} \in \bar{\it H} \mbox{ is }$

its conjugate). In particular: FDHILB is tannakian (take $G = \{*\}$).

[Note: The construct <u>Rep</u> G is amnestic and transportable, so we can and will assume that its monoidal structure is strict (cf. 3.12).]

10.3 <u>RAPPEL</u> An additive functor $F:A \rightarrow B$ between abelian categories A and B is exact if it preserves finite limits and finite colimits.

Accordingly, since a tannakian category is not only abelian but also autonomous, $\forall X \in Ob C$, the functors

$$\begin{bmatrix} - & & x, & lhom(x, --) \\ & & x & & --, & rhom(x, --) \end{bmatrix}$$

are exact.

If \underline{C} is tannakian, then e is irreducible and

dim:Ob $\underline{C} \rightarrow Mor(e,e)$

has the following properties.

- 1. dim X = dim \overline{X} .
- 2. $\dim(X \otimes Y) = (\dim X) (\dim Y)$.
- 3. $\dim(X \oplus Y) = \dim X + \dim Y$.
- 4. dim e = 1, dim 0 = 0.

10.4 <u>LEMMA</u> If X is not a zero object, then dim X (= $r_X^* \circ r_X$) ≥ 1 .

<u>PROOF</u> First, from the positivity of the involution, dim X > 0. But $X \otimes \overline{X}$ contains e as a direct summand, thus

$$(\dim X)^2 \ge 1 \Rightarrow \dim X \ge 1.$$

[Note: If dim X = 1, then $X \otimes \overline{X} \approx e \approx \overline{X} \otimes X$.]

Given $X \neq 0$ in Ob C, define

$$\Pi_n^X: \mathfrak{S}_n \to \operatorname{Aut} X^{\mathfrak{A}n}$$

as in 4.5.

<u>N.B.</u> Π_n^X is a homomorphism from \mathscr{G}_n to the unitary group of Mor(X,X). Put

$$x^{(0)} = e,$$

$$\begin{bmatrix} - & Sym_0^X = id_e \\ & \\ & Alt_0^X = id_e, \end{bmatrix}$$

and for $n \in \underline{N}$, put

Sym_n^X =
$$\frac{1}{n!} \sum_{\sigma \in \mathfrak{F}_n} \Pi_n^X(\sigma)$$

Alt_n^X = $\frac{1}{n!} \sum_{\sigma \in \mathfrak{F}_n} (\operatorname{sgn} \sigma) \Pi_n^X(\sigma)$.

Then

are projections.

10.5 LEMMA We have

$$tr_{X^{\otimes n}}(Alt_{n}^{X})$$
$$= \frac{1}{n!} (\dim X) (\dim X - 1) \dots (\dim X - n + 1)$$

PROOF The key preliminary is the observation that

$$\operatorname{tr}_{X^{\operatorname{\mathfrak{S}n}}}(\Pi_{n}^{X}(\sigma)) = (\dim X)^{\#\sigma},$$

where $\#\sigma$ is the number of cycles into which σ decomposes, thus

$$\operatorname{tr}_{X^{\underline{\otimes}n}}(\operatorname{Alt}_{n}^{X}) = \frac{1}{n!} \sum_{\sigma \in \mathfrak{F}_{n}} (\operatorname{sgn} \sigma) (\dim X)^{\#\sigma}.$$

But for every complex number z,

$$\sum_{\sigma \in \mathfrak{S}_n} (\operatorname{sgn} \sigma) z^{\#\sigma} = z(z-1) \dots (z-n+1).$$

10.6 <u>THEOREM</u> \forall nonzero X in Ob C,

dim
$$X \in N$$
.

<u>PROOF</u> Let $A_n(X)$ be the subobject of $X^{\otimes n}$ corresponding to Alt_n^X . Fix an isometry $f:A_n(X) \to X^{\otimes n}$ such that $f \circ f^* = Alt_n^X$ -- then $tr_{X^{\otimes n}}(Alt_n^X)$ $= tr_{X^{\otimes n}}(f \circ f^*)$ $= tr_{A_n}(X)(f^* \circ f) \quad (cf. 6.3)$

$$= \operatorname{tr}_{A_{n}(X)} (\operatorname{id}_{A_{n}(X)})$$

= dim $A_{n}(X) \ge 1$ (cf. 10.4).

On the other hand, thanks to 10.5,

$$\operatorname{tr}_{X^{\otimes n}}(\operatorname{Alt}_{n}^{X})$$

is negative for some $n \in N$ unless dim $X \in N$.

10.7 <u>LEMMA</u> Let $d = \dim X$ -- then

dim
$$A_d(X) = tr_{X^{\otimes d}}Alt_d^X = \frac{d!}{d!} = 1.$$

The isomorphism class of $A_d(X)$ is called the <u>determinant</u> of X (written det (X)). Properties:

- 1. det(\overline{X}) $\approx \overline{\det(X)}$;
- 2. det(X \oplus Y) \approx det(X) \otimes det(Y);
- 3. det (X $\oplus \overline{X}$) \approx e.

§11. FIBER FUNCTORS

Let \underline{C} be a tannakian category -- then a symmetric embedding functor

$\mathcal{F}: \underline{\mathsf{C}} \to \underline{\mathsf{FDHILB}}$

is called a fiber functor.

E.g.: Take $\underline{C} = \underline{Rep} G$ (cf. 10.2) — then the forgetful functor

$U: \underline{Rep} \ G \rightarrow \underline{FDHILB}$

is a fiber functor.

<u>N.B.</u> It is a nontrivial result that every tannakian category admits a fiber functor (proof omitted).

11.1 REMARK Let

$$\mathcal{F}:C \rightarrow FDHILB$$

be a fiber functor. Consider

$$A_{F} = \bigoplus_{i \in \mathbf{I}_{\underline{C}}} B(FX_{\underline{i}}),$$

viewed as a subset of $\underline{Nat}(F,F)$ -- then the <u>coinverse</u> is the map $S:A_F \to A_F$ defined by

$$\mathrm{S}(\alpha)_{X} = \mathcal{F}(\mathrm{id}_{X} \otimes r_{X}^{*}) \circ \mathrm{id}_{\mathcal{F}X} \cong \alpha_{\overline{X}} \cong \mathrm{id}_{\mathcal{F}X} \circ \mathcal{F}(\overline{r}_{X} \otimes \mathrm{id}_{X}),$$

matters being slightly imprecise in that the identification

$$F(X \otimes \overline{X} \otimes X) \approx FX \otimes F\overline{X} \otimes FX$$

has been suppressed. It is not difficult to see that the equation defining $S(\alpha)_X$ is independent of the choice $(\bar{X}, r_X, \bar{r}_X)$ of a conjugate for X and $\forall f \in Mor(X, Y)$, the diagram



commutes. Algebraically, S is linear and antimultiplicative. Moreover,

$$S \circ * \circ S \circ * = id_{A_{F}}$$

hence S is invertible.

[Note: There are various relations among Δ, ε, S which, however, need not be detailed. Still, despite appearances, in general $(A_{\mathbf{f}}, \Delta, \varepsilon, S)$ is not a Hopf *-algebra but rather in the jargon of the trade is a "cocommutative discrete algebraic quantum group".]

Write $ff(\underline{C})$ for the full subcategory of

whose objects are the fiber functors -- then ff(C) is a groupoid (cf. 5.12).

11.2 <u>THEOREM</u> $ff(\underline{C})$ is a transitive groupoid, i.e., if F_1, F_2 are fiber functors, then F_1, F_2 are isomorphic.

Definition: Given fiber functors F_1, F_2 , a <u>unitary</u> monoidal natural transformation $\alpha: F_1 \to F_2$ is a monoidal natural transformation such that $\forall X \in Ob C$,

$$\alpha_{X}:\mathcal{F}_{1}X \to \mathcal{F}_{2}X$$

is unitary.

Write $ff^*(\underline{C})$ for the category whose objects are the fiber functors and whose morphisms are the unitary monoidal natural transformations — then $ff^*(\underline{C})$ is a subcategory of $ff(\underline{C})$.

11.3 <u>THEOREM</u> ff*(<u>C</u>) is a transitive groupoid, i.e., if F_1, F_2 are fiber functors, then F_1, F_2 are unitarily isomorphic.

Obviously,

As for the proof of 11.3, there will be three steps.

<u>Step 1</u>: Construct a commutative unital *-algebra $A(F_1,F_2)$ whose dual space is in a one-to-one correspondence with the natural transformations $F_1 \rightarrow F_2$, to wit:

$$\operatorname{Nat}(\mathcal{F}_1, \mathcal{F}_2) \iff \operatorname{A}(\mathcal{F}_1, \mathcal{F}_2)^*$$

<u>Step 2</u>: Under this bijection, prove that the monoidal natural transformations correspond to the nonzero multiplicative linear functionals on $A(\mathcal{F}_1, \mathcal{F}_2)$ and the unitary monoidal natural transformations correspond to the *-preserving multiplicative linear functionals on $A(\mathcal{F}_1, \mathcal{F}_2)$.

Step 3: Establish that $A(F_1,F_2)$ admits a C*-norm, thus is a pre-C*-algebra. Therefore, since the structure space $\Delta(\vec{A}(F_1,F_2))$ of the C*-completion $\vec{A}(F_1,F_2)$ of $A(F_1,F_2)$ is not empty, it follows that $Mor(F_1,F_2)$ is also not empty, from which 11.3. [Note: Here, of course, Mor is computed in $ff^*(\underline{C})$.]

To fix notation, bear in mind that there are isomorphisms

$$\begin{bmatrix} \xi^{1} : \underline{e} \to \mathcal{F}_{1} e \\ \xi^{2} : \underline{e} \to \mathcal{F}_{2} e \end{bmatrix}, \begin{bmatrix} \Xi^{1}_{X,Y} : \mathcal{F}_{1} X \stackrel{\otimes}{\cong} \mathcal{F}_{1} Y \to \mathcal{F}_{1} (X \stackrel{\otimes}{\boxtimes} Y) \\ \Xi^{2}_{X,Y} : \mathcal{F}_{2} X \stackrel{\otimes}{\cong} \mathcal{F}_{2} Y \to \mathcal{F}_{2} (X \stackrel{\otimes}{\boxtimes} Y) \end{bmatrix}$$

subject to the compatibility conditions enumerated in §2.

Let $\mathsf{A}_0(\mathcal{F}_1,\mathcal{F}_2)$ be the complex vector space

Given $X \in Ob \subseteq$ and $\phi \in Mor(\mathcal{F}_2 X, \mathcal{F}_1 X)$, write $[X, \phi]_0$ for the element of $A_0(\mathcal{F}_1, \mathcal{F}_2)$ that is ϕ at X and is zero elsewhere -- then $A_0(\mathcal{F}_1, \mathcal{F}_2)$ is simply the complex linear span of the $[X, \phi]_0$. Define a product in $A_0(\mathcal{F}_1, \mathcal{F}_2)$ by stipulating that

$$[\mathbf{X},\phi]_{0} \cdot [\mathbf{Y},\psi]_{0} = [\mathbf{X} \otimes \mathbf{Y},\mathbf{u}]_{0},$$

where u is the composition

$$F_{2}(X \otimes Y) \xrightarrow{(\Xi_{X,Y}^{2})^{-1}} F_{2}X \boxtimes F_{2}Y$$

$$\xrightarrow{\varphi \boxtimes \psi} F_{1}X \boxtimes F_{1}Y$$

$$\xrightarrow{\Xi_{X,Y}^{1}} F_{1}(X \otimes Y).$$

11.4 <u>LEMMA</u> $A_0(\mathcal{F}_1, \mathcal{F}_2)$ is associative.

11.5 <u>**LEMMA**</u> $A_0(F_1,F_2)$ is unital.

PROOF Let

$$l_{A_0} = [e, \xi^1 \circ (\xi^2)^{-1}]_0.$$

Then l_{A_0} is the unit. E.g.: Consider

$$[x,\phi]_0 \cdot [e,\xi^1 \circ (\xi^2)^{-1}]_0 = [x,u]_0,$$

the claim being that the composite

$$F_{2}X = F_{2}(X \otimes e) \xrightarrow{(\Xi_{X,e}^{2})^{-1}} F_{2}X \otimes F_{2}e$$

$$\xrightarrow{\phi \otimes (\xi^{1} \circ (\xi^{2})^{-1})} F_{1}X \otimes F_{1}e$$

$$\xrightarrow{\Xi_{X,e}^{1}} F_{1}(X \otimes e) = F_{1}X$$

reduces to $\boldsymbol{\phi}$ itself. To see this, recall that the composition

is the identity morphism of $\mathcal{F}_1^{} \mathbf{X}$ and the composition

$$F_{2}x = F_{2}x \stackrel{\mathbb{Q}}{\cong} e \xrightarrow{f_{2}x} \stackrel{\mathbb{Q}}{\longrightarrow} F_{2}x \stackrel{\mathbb{Q}}{\cong} F_{2}e \xrightarrow{\Xi_{x,e}^{2}} F_{2}(x \otimes e) = F_{2}x$$

is the identity morphism of ${\rm F}_2 {\rm X}.$ Now write

$$\begin{array}{c} \Xi_{X,e}^{1} \circ \phi \ \underline{\Theta} \ (\xi^{1} \circ (\xi^{2})^{-1}) \circ (\Xi_{X,e}^{2})^{-1} \\ = \operatorname{id}_{\mathcal{F}_{1}X} \circ \operatorname{id}_{\mathcal{F}_{1}X} \ \underline{\Theta} \ (\xi^{1})^{-1} \circ \phi \ \underline{\Theta} \ (\xi^{1} \circ (\xi^{2})^{-1}) \circ \operatorname{id}_{\mathcal{F}_{2}X} \ \underline{\Theta} \ \xi^{2} \circ \operatorname{id}_{\mathcal{F}_{2}X} \end{array}$$

$$= \operatorname{id}_{\mathcal{F}_{1}X} \circ (\operatorname{id}_{\mathcal{F}_{1}X} \circ \phi \circ \operatorname{id}_{\mathcal{F}_{2}X} \underline{Q} (\xi^{-1})^{-1} \circ (\xi^{1} \circ (\xi^{2})^{-1} \circ \xi^{2}) \circ \operatorname{id}_{\mathcal{F}_{2}X}$$
$$= \operatorname{id}_{\mathcal{F}_{1}X} \circ \phi \underline{Q} \operatorname{id}_{\underline{P}_{2}X} = \operatorname{id}_{\mathcal{F}_{1}X} \circ \phi \circ \operatorname{id}_{\mathcal{F}_{2}X} = \phi.$$

Let $\mathcal{I}_0(\mathcal{F}_1, \mathcal{F}_2)$ be the complex linear span of the

$$[x, a \circ F_2 f]_0 - [Y, F_1 f \circ a]_0,$$

where

$$f \in Mor(X,Y)$$
, $a \in Mor(\mathcal{F}_2Y,\mathcal{F}_1X)$.

Then $I_0(\mathcal{F}_1,\mathcal{F}_2)$ is an ideal in $A_0(\mathcal{F}_1,\mathcal{F}_2)$.

Denote by $A(F_1,F_2)$ the quotient algebra

$$A_0(F_1,F_2)/I_0(F_1,F_2),$$

let

$$\mathrm{pr}\!:\!\mathsf{A}_0(\mathcal{F}_1,\mathcal{F}_2) \twoheadrightarrow \mathsf{A}(\mathcal{F}_1,\mathcal{F}_2)$$

be the projection, and put

$$[\mathbf{X},\phi] = \operatorname{pr}[\mathbf{X},\phi]_{0}.$$

11.6 EXAMPLE Let $f: X \rightarrow X$ be an isomorphism — then

$$[\mathbf{X}, \phi] = [\mathbf{X}, \mathcal{F}_{1} \mathbf{f} \circ \mathcal{F}_{1} \mathbf{f}^{-1} \circ \phi]$$
$$= [\mathbf{X}, \mathcal{F}_{1} \mathbf{f}^{-1} \circ \phi \circ \mathcal{F}_{2} \mathbf{f}].$$

11.7 EXAMPLE Let

$$\Phi \in \operatorname{Mor}(\mathcal{F}_{2}(\overline{X} \otimes X), \mathcal{F}_{1}(\overline{X} \otimes X)).$$

Then

$$[\vec{X} \otimes X, \mathcal{F}_{1}r_{X} \circ \mathcal{F}_{1}r_{X}^{*} \circ \Phi]$$
$$= [e, \mathcal{F}_{1}r_{X}^{*} \circ \Phi \circ \mathcal{F}_{2}r_{X}].$$

[Note: We also have

$$[\overline{X} \otimes X, \mathcal{F}_{1}(\mathbf{r}_{X} \circ \mathbf{r}_{X}^{*}) \circ \Phi]$$

= $[\overline{X} \otimes X, \Phi \circ \mathcal{F}_{2}(\mathbf{r}_{X} \circ \mathbf{r}_{X}^{*})].]$

11.8 <u>REMARK</u> Every $A \in A(F_1, F_2)$ can be written as $[X, \phi]$ for a suitable choice of X and ϕ . Thus suppose that $A = \sum_{i} [X_i, \phi_i]$, put $X = \bigoplus_{i} X_i$, and choose isometries $v_i: X_i \to X$ such that $\sum_{i} v_i \circ v_i^* = id_X$ — then

$$\mathbf{a}_{\mathbf{i}} = \boldsymbol{\phi}_{\mathbf{i}} \circ \boldsymbol{F}_{2} \mathbf{v}_{\mathbf{i}}^{\star} \in \operatorname{Mor}(\boldsymbol{F}_{2} \mathbf{X}, \boldsymbol{F}_{1} \mathbf{X}_{\mathbf{i}})$$

=>

$$A = \sum_{i} [X_{i}, \phi_{i}]$$

$$= \sum_{i} [X_{i}, \phi_{i} \circ id_{\mathcal{F}_{2}X_{i}}]$$

$$= \sum_{i} [X_{i}, \phi_{i} \circ \mathcal{F}_{2}(v_{i}^{*} \circ v_{i})]$$

$$= \sum_{i} [X_{i}, \phi_{i} \circ \mathcal{F}_{2}v_{i}^{*} \circ \mathcal{F}_{2}v_{i}]$$

$$= \sum_{i} [X_{i}, a_{i} \circ \mathcal{F}_{2}v_{i}]$$

$$= \sum_{i} [X, \mathcal{F}_{1}v_{i} \circ a_{i}]$$

$$= \sum_{i} [X, \mathcal{F}_{1} v_{i} \circ \phi_{i} \circ \mathcal{F}_{2} v_{i}^{*}]$$

$$= [X, \sum_{i} \mathcal{F}_{1} v_{i} \circ \phi_{i} \circ \mathcal{F}_{2} v_{i}^{*}]$$

$$= [X, \phi],$$

where

$$\phi = \sum_{i} \mathcal{F}_{1} \mathbf{v}_{i} \circ \phi_{i} \circ \mathcal{F}_{2} \mathbf{v}_{i}^{*} \in \operatorname{Mor}\left(\mathcal{F}_{2} \mathbf{X}, \mathcal{F}_{1} \mathbf{X}\right).$$

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11.9 LEMMA
$$A(F_1, F_2)$$
 is commutative.

 \underline{PROOF} Let

$$\begin{bmatrix} [X,\phi]_0 & (\phi:\mathcal{F}_2 X \to \mathcal{F}_1 X) \\ [Y,\psi]_0 & (\psi:\mathcal{F}_2 Y \to \mathcal{F}_1 Y) \end{bmatrix}$$

be elements of $A_0(\mathcal{F}_1,\mathcal{F}_2)$ — then

$$[x,\phi]_{0} \cdot [Y,\psi]_{0} = [x \otimes Y,\Xi_{X,Y}^{1} \circ \phi \underline{\otimes} \psi \circ (\Xi_{X,Y}^{2})^{-1}].$$

On the other hand,

____ . . .

$$[\mathbf{Y}, \psi]_{0} \cdot [\mathbf{X}, \phi]_{0} = [\mathbf{Y} \otimes \mathbf{X}, \Xi^{1}_{\mathbf{Y}, \mathbf{X}} \circ \psi \otimes \phi \circ (\Xi^{2}_{\mathbf{Y}, \mathbf{X}})^{-1}]$$

and there is a commutative diagram

$$\begin{array}{c} & \stackrel{^{\mathsf{T}} \mathcal{F}_{2} Y, \mathcal{F}_{2} X}{} & \stackrel{^{\mathsf{T}} \mathcal{F}_{2} Y, \mathcal{F}_{2} X}{} & \rightarrow & \mathcal{F}_{2} X \underline{\otimes} \mathcal{F}_{2} Y \\ \psi \underline{\otimes} \phi \\ \downarrow \\ & \downarrow$$

Thus

$$\begin{split} & \Xi_{\mathbf{Y},\mathbf{X}}^{1} \circ \psi \, \underline{\boldsymbol{\Theta}} \, \phi \, \circ \, (\Xi_{\mathbf{Y},\mathbf{X}}^{2})^{-1} \\ & = \, \Xi_{\mathbf{Y},\mathbf{X}}^{1} \circ \, \tau_{\mathcal{F}_{1}\mathbf{X},\mathcal{F}_{1}\mathbf{Y}} \circ \phi \, \underline{\boldsymbol{\Theta}} \, \psi \, \circ \, \tau_{\mathcal{F}_{2}\mathbf{Y},\mathcal{F}_{2}\mathbf{X}} \circ \, (\Xi_{\mathbf{Y},\mathbf{X}}^{2})^{-1}. \end{split}$$

But there are also commutative diagrams



and



Thus

$$\begin{split} & \Xi_{\mathbf{Y},\mathbf{X}}^{1} \circ \tau_{\mathcal{F}_{1}\mathbf{X},\mathcal{F}_{1}\mathbf{Y}} \circ \phi \underline{\otimes} \psi \circ \tau_{\mathcal{F}_{2}\mathbf{Y},\mathcal{F}_{2}\mathbf{X}} \circ (\Xi_{\mathbf{Y},\mathbf{X}}^{2})^{-1} \\ & = \mathcal{F}_{1}\tau_{\mathbf{X},\mathbf{Y}} \circ \Xi_{\mathbf{X},\mathbf{Y}}^{1} \circ \phi \underline{\otimes} \psi \circ (\Xi_{\mathbf{X},\mathbf{Y}}^{2})^{-1} \circ \mathcal{F}_{2}\tau_{\mathbf{Y},\mathbf{X}}. \end{split}$$

Let
$$f = \tau_{X,Y}$$
 and put
 $a = \Xi_{X,Y}^1 \circ \phi \ \mathfrak{Q} \ \psi \circ (\Xi_{X,Y}^2)^{-1} \circ \mathcal{F}_2 \tau_{Y,X}$.
Then
 $f \in Mor(X \ \mathfrak{Q} \ Y, Y \ \mathfrak{Q} \ X)$

and

Then

$$a \in Mor(F_2(Y \ \& \ X), F_1(X \ \& \ Y)).$$

Moreover

$$[\mathbf{Y}, \boldsymbol{\psi}]_0 \cdot [\mathbf{X}, \boldsymbol{\phi}]_0 = [\mathbf{Y} \otimes \mathbf{X}, \mathbf{F}_1 \mathbf{f} \circ \mathbf{a}]_0.$$

Meanwhile

$$F_{2}^{T}Y, X \circ F_{2}^{T}X, Y$$

$$= F_{2}(T_{Y,X} \circ T_{X,Y})$$

$$= F_{2}(id_{X \otimes Y})$$

$$= id_{F_{2}}(X \otimes Y)'$$

so

._____

$$[\mathbf{X},\phi]_{0} \cdot [\mathbf{Y},\psi]_{0} = [\mathbf{X} \boldsymbol{\Theta} \mathbf{Y},\mathbf{a} \circ \mathcal{F}_{2}\mathbf{f}]_{0}$$

Therefore

$$\begin{split} \left[\mathbf{X}, \phi \right]_{0} \cdot \left[\mathbf{Y}, \psi \right]_{0} &- \left[\mathbf{Y}, \psi \right]_{0} \cdot \left[\mathbf{X}, \phi \right]_{0} \in I_{0}(\mathcal{F}_{1}, \mathcal{F}_{2}) \\ &=> \\ \left[A_{0}(\mathcal{F}_{1}, \mathcal{F}_{2}), A_{0}(\mathcal{F}_{1}, \mathcal{F}_{2}) \right] &\subset I_{0}(\mathcal{F}_{1}, \mathcal{F}_{2}) \end{split}$$

And this implies that $A(\mathcal{F}_1, \mathcal{F}_2)$ is commutative.

Given $[X,\phi]_0$, choose a conjugate (\bar{X},r_X,\bar{r}_X) for X and let

 $[\mathbf{X},\phi]_0^\star = \ [\bar{\mathbf{X}},\bar{\phi}]_0,$

where $\overline{\phi}$ is the composition

$$F_{2}\overline{\mathbf{x}} = \underline{\mathbf{e}} \ \underline{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{F_{1}\mathbf{r}_{\mathbf{X}} \ \underline{\mathbf{e}} \ \mathrm{id}}_{\mathbf{F}_{1}\mathbf{r}_{\mathbf{X}} \ \underline{\mathbf{e}} \ \mathrm{id}} \longrightarrow F_{1}\mathbf{e} \ \underline{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{F_{2}\overline{\mathbf{x}}}_{\mathbf{F}_{2}\overline{\mathbf{x}}} \qquad \underbrace{F_{1}\mathbf{r}_{\mathbf{X}} \ \underline{\mathbf{e}} \ \mathrm{id}}_{\mathbf{X},\mathbf{X}} \longrightarrow F_{1}(\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ \mathbf{x}) \ \underline{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{(\underline{\mathbf{e}}^{1} \)^{-1} \ \underline{\mathbf{e}} \ \mathrm{id}}_{\mathbf{X},\mathbf{X}} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \xrightarrow{\mathbf{e}}_{\mathbf{X},\mathbf{x}} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{2}\mathbf{x} \ \underline{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \ \underline{\mathbf{e}}^{2}_{\mathbf{X},\mathbf{x}} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{2}\mathbf{x} \ \underline{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{\mathbf{e}} \ F_{2}\overline{\mathbf{x}} \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \ \underline{\mathbf{e}}^{2}_{\mathbf{X},\mathbf{x}} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{2}(\mathbf{x} \ \underline{\mathbf{e}} \ \overline{\mathbf{x}}) \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \ \underline{\mathbf{e}}^{2}_{\mathbf{x}} \xrightarrow{\mathbf{e}}_{\mathbf{X},\mathbf{x}} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{2}\mathbf{e} \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \ (\xi^{2})^{-1} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ F_{2}\mathbf{e} \qquad \underbrace{\mathbf{id}}_{\mathbf{E}} \ (\xi^{2})^{-1} \longrightarrow F_{1}\overline{\mathbf{x}} \ \underline{\mathbf{e}} \ \mathbf{e}} = F_{1}\overline{\mathbf{x}}.$$

N.B. We have

$$(\Xi^{1})^{-1} \bigotimes \operatorname{id}_{\mathcal{F}_{2}} \tilde{x} \circ \mathcal{F}_{1} r_{X} \bigotimes \operatorname{id}_{\mathcal{F}_{2}} \tilde{x} \circ \xi^{1} \bigotimes \operatorname{id}_{\mathcal{F}_{2}} \tilde{x}$$

$$= (\Xi^{1})^{-1} \circ \mathcal{F}_{1} r_{X} \circ \xi^{1} \bigotimes \operatorname{id}_{\mathcal{F}_{2}} \tilde{x}$$

$$= r_{\mathcal{F}_{1}} x \bigotimes \operatorname{id}_{\mathcal{F}_{2}} \tilde{x}$$

$$(cf. 9.10)$$

anđ

$$\stackrel{\text{id}}{\stackrel{\mathcal{F}_{1}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{1}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{1}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{1}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathcal{F}_{2}\bar{x}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=}} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{\stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{G}}{=} \stackrel{\mathfrak{$$

Therefore

 $\bar{\phi} = \operatorname{id}_{\mathcal{F}_1 \bar{X}} \underbrace{\mathfrak{Q}}_{\mathcal{F}_2 X} \overset{\mathbf{r}}{\mathfrak{F}_2 X} \circ \operatorname{id}_{\mathcal{F}_1 \bar{X}} \underbrace{\mathfrak{Q}}_{\mathcal{P}} \phi^* \otimes \operatorname{id}_{\mathcal{F}_2 \bar{X}} \circ \overset{\mathbf{r}}{\mathfrak{F}_2 \bar{X}} \underbrace{\mathfrak{Q}}_{\mathcal{P}_2 \bar{X}} \overset{\mathrm{id}}{\mathfrak{F}_2 \bar{X}}.$

[Note: By definition, $\phi \in Mor(\mathcal{F}_2 X, \mathcal{F}_1 X)$, so $^{\vee}\phi \in Mor(\mathcal{F}_1 \overline{X}, \mathcal{F}_2 \overline{X})$, where, as in

§5,

$${}^{\vee}\phi = \varepsilon_{\mathcal{F}_1} X \stackrel{\texttt{Q}}{=} \overset{\texttt{id}}{\stackrel{\mathcal{F}_2}{\times}} \circ \stackrel{\texttt{id}}{=} \frac{\texttt{Q}}{\stackrel{\mathcal{F}_1}{\times}} \stackrel{\texttt{Q}}{=} \overset{\texttt{id}}{\stackrel{\mathcal{F}_2}{\times}} \circ \stackrel{\texttt{id}}{=} \frac{\texttt{Q}}{\stackrel{\mathcal{F}_2}{\times}} \circ \stackrel{\texttt{id}}{=} \frac{\texttt{Q}}{\xrightarrow{\mathcal{F}_2}} \circ \stackrel{\texttt{id}}{=} \stackrel{\texttt{$$

or still,

$$\overset{\vee}{\phi} = r_{\mathcal{F}_{1}X}^{\star} \overset{\mathbb{Q}}{=} \overset{\mathrm{id}}{\stackrel{\mathcal{F}_{2}}{\times}} \circ \overset{\mathrm{id}}{=} \overset{\mathbb{Q}}{=} \overset{\mathrm{id}}{\stackrel{\mathcal{Q}}{=}} \overset{\mathrm{id}}{=} \overset{\mathrm{o}}{=} \overset{\mathrm{id}}{\stackrel{\mathcal{F}_{2}}{\times}} \circ \overset{\mathrm{id}}{=} \overset{\mathbb{Q}}{=} \overset{\mathrm{id}}{\stackrel{\mathcal{F}_{2}}{\times}} \circ \overset{\mathrm{id}}{=} \overset{\mathbb{Q}}{=} \overset{\mathrm{id}}{=} \overset{\mathrm{o}}{=} \overset{\mathrm{o}}{=} \overset{\mathrm{id}}{=} \overset{\mathrm{o}}{=} \overset{\mathrm{o}}{=} \overset{\mathrm{o}}{=} \overset{\mathrm{id}}{=} \overset{\mathrm{o}}{=} \overset$$

Therefore

 $\overline{\phi} = (^{\vee}\phi) \star .]$

Replacing $(\bar{X}, r_X, \bar{r}_X)$ by $(\bar{X}^*, r_X^1, \bar{r}_X^1)$ and using 9.4, one finds that

$$[\bar{\mathbf{x}},\bar{\boldsymbol{\phi}}]_0 - [\bar{\mathbf{x}}',\bar{\boldsymbol{\phi}}']_0 \in I_0(\mathcal{F}_1,\mathcal{F}_2).$$

Therefore the image of $[X,\phi]_0^*$ in $A(\mathcal{F}_1,\mathcal{F}_2)$ is independent of the choice of a conjugate for X.

11.10 LEMMA $I_0(F_1,F_2)$ is *-invariant.

 $\texttt{Consequently, } \star: \texttt{A}_0(\texttt{F}_1,\texttt{F}_2) \to \texttt{A}_0(\texttt{F}_1,\texttt{F}_2) \texttt{ induces a map } \star: \texttt{A}(\texttt{F}_1,\texttt{F}_2) \to \texttt{A}(\texttt{F}_1,\texttt{F}_2).$

11.11 <u>LEMMA</u> $A(F_1,F_2)$ is a *-algebra.

Summary: $A(F_1, F_2)$ is a commutative unital *-algebra.

Accordingly, to complete Step 1, it remains to construct an isomorphism between $A(F_1,F_2)^*$ and $Nat(F_1,F_2)$.

On general grounds,

$$A_0(\mathcal{F}_1,\mathcal{F}_2)^* = \prod_{X \in Ob \ \underline{C}} \operatorname{Mor} (\mathcal{F}_2 X, \mathcal{F}_1 X)^*.$$

But the pairing

$$\mathsf{Mor}(\mathcal{F}_{2}X,\mathcal{F}_{1}X) \times \mathsf{Mor}(\mathcal{F}_{1}X,\mathcal{F}_{2}X) \to \mathsf{C}$$

that sends $\phi \times \psi$ to tr($\phi \circ \psi$) is nondegenerate, thus

$$A_0(\mathcal{F}_1,\mathcal{F}_2) * \approx \prod_{X \in Ob \ C} \operatorname{Mor}(\mathcal{F}_1X,\mathcal{F}_2X).$$

On the other hand, $\operatorname{Nat}(\mathcal{F}_1,\mathcal{F}_2)$ consists of those elements

$$\alpha \in \prod_{X \in Ob \ \underline{C}} Mor(\mathcal{F}_1 X, \mathcal{F}_2 X)$$

such that $\forall f \in Mor(X,Y)$,

$$\mathcal{F}_{2}f \circ \alpha_{X} = \alpha_{Y} \circ \mathcal{F}_{1}f,$$

and the dual of $A(F_1,F_2)$ is the subspace of $A_0(F_1,F_2)^*$ comprised of those elements that vanish identically on $I_0(F_1,F_2)$. To characterize the latter, take an

$$\alpha \in \prod_{X \in Ob \ \underline{C}} Mor(\mathcal{F}_1 X, \mathcal{F}_2 X)$$

and suppose that $\forall \ \mathtt{A} \in \mathtt{I}_0(\mathtt{F}_1, \mathtt{F}_2)$,

 $< A, \alpha > = 0$

or still,

$$\langle [X,a \circ \mathcal{F}_2 f]_0 - [Y,\mathcal{F}_1 f \circ a]_0, \alpha \rangle = 0$$

for all

$$f \in Mor(X,Y)$$
, $a \in Mor(\mathcal{F}_2Y,\mathcal{F}_1X)$.

I.e.:

$$\operatorname{tr}_{\mathcal{F}_{1}X}(a \circ \mathcal{F}_{2}f \circ \alpha_{X}) = \operatorname{tr}_{\mathcal{F}_{1}Y}(\mathcal{F}_{1}f \circ a \circ \alpha_{Y}).$$

From the nondegeneracy of the trace, it then follows that

$$\mathcal{F}_2 \mathbf{f} \circ \alpha_{\mathbf{X}} = \alpha_{\mathbf{Y}} \circ \mathcal{F}_1 \mathbf{f},$$

implying thereby that

$$\alpha \in \operatorname{Nat}(F_1,F_2)$$
.

11.12 LEMMA Under the bijection

$$\operatorname{Nat}(\mathcal{F}_1, \mathcal{F}_2) \iff \operatorname{A}(\mathcal{F}_1, \mathcal{F}_2)^*,$$

the monoidal natural transformations correspond to the nonzero multiplicative linear functionals on $A(F_1,F_2)$.

<u>PROOF</u> To say that a linear functional on $A(F_1,F_2)$ corresponding to an $\alpha \in \operatorname{Nat}(F_1,F_2)$ is multiplicative amounts to saying that

$$= \langle [X,\phi],\alpha \rangle \cdot \langle [Y,\psi],\alpha \rangle$$

for all

$$[\mathbf{X},\phi], [\mathbf{Y},\psi] \in \mathsf{A}(\mathbf{F}_1,\mathbf{F}_2).$$

Since < ---, α > is null on $I_0(F_1,F_2)$, it suffices to work upstairs, hence explicated we have

$$\begin{aligned} \operatorname{tr}_{\mathcal{F}_{1}}(X \otimes Y) \stackrel{(\Xi_{X,Y}^{1} \circ \phi \otimes \psi \circ (\Xi_{X,Y}^{2})^{-1} \circ \alpha_{X \otimes Y})}{= \operatorname{tr}_{\mathcal{F}_{1}X}(\phi \circ \alpha_{X}) \operatorname{tr}_{\mathcal{F}_{1}Y}(\psi \circ \alpha_{Y})} \\ = \operatorname{tr}_{\mathcal{F}_{1}X \otimes \mathcal{F}_{1}Y}(\phi \circ \alpha_{X}) \otimes (\psi \circ \alpha_{Y})) \\ = \operatorname{tr}_{\mathcal{F}_{1}X \otimes \mathcal{F}_{1}Y}(\phi \otimes \psi \circ \alpha_{X} \otimes \alpha_{Y}). \end{aligned}$$

Therefore

$$\alpha_{\mathbf{X} \ \mathbf{Q} \ \mathbf{Y}} = \Xi_{\mathbf{X},\mathbf{Y}}^{2} \circ \alpha_{\mathbf{X}} \ \mathbf{\underline{Q}} \ \alpha_{\mathbf{Y}} \circ (\Xi_{\mathbf{X},\mathbf{Y}}^{1})^{-1},$$

the condition that α be monoidal.

[Note: Tacitly,

$$<1_{A_0}, \alpha > = 1$$

$$< [e, \xi^1 \circ (\xi^2)^{-1}]_0, \alpha > = 1$$

or still,

$$tr_{F_1e}(\xi^1 \circ (\xi^2)^{-1} \circ \alpha_e) = 1,$$

from which the commutativity of the diagram



11.13 LEMMA Under the bijection

$$\operatorname{Nat}(\mathcal{F}_1, \mathcal{F}_2) \iff \operatorname{A}(\mathcal{F}_1, \mathcal{F}_2)^*,$$

the unitary monoidal natural transformations correspond to the *-preserving non-zero multiplicative linear functionals on $A(F_1,F_2)$.

<u>PROOF</u> Given $[X, \varphi] \in A(\mathcal{F}_1, \mathcal{F}_2)$, the claim is that

$$\langle [X,\phi]^*,\alpha\rangle = \overline{\langle [X,\phi],\alpha\rangle} (= \langle [X,\phi],\alpha\rangle^* \dots)$$

iff $\alpha_X^{\star} = \alpha_X^{-1}$.

From the definitions,

$$< [\mathbf{X}, \phi], \alpha > = \operatorname{tr}_{\mathcal{F}_1} \mathbf{X}^{(\phi \circ \alpha_X)}$$

=>

$$\overline{\langle [X,\phi],\alpha\rangle} = \operatorname{tr}_{\mathcal{F}_{1}X}(\phi^{\star} \circ \alpha_{X}^{\star}).$$

In the other direction,

$$< [X, \phi]^{*}, \alpha > = < [X, \phi], \alpha >$$

$$= \operatorname{tr}_{\mathcal{F}_{1}} (\overline{\phi} \circ \alpha)$$

$$\mathcal{F}_{1} \overline{X} \qquad \overline{X}$$

$$= \operatorname{r}^{*}_{\mathcal{F}_{1}} \circ \operatorname{id}_{\mathcal{F}_{1}} X \stackrel{\textcircled{0}}{=} (\overline{\phi} \circ \alpha) \circ \operatorname{r}_{\mathcal{F}_{1}} \overline{X}$$

$$= \overline{r}_{F_{1}X}^{*} \circ \operatorname{id}_{F_{1}X} \underline{\mathfrak{Q}} (\overline{\phi} \circ \alpha_{\overline{X}}) \circ \overline{r}_{F_{1}X}.$$

$$\overline{r}_{F_{2}X}^{*} \circ \phi^{*} \underline{\mathfrak{Q}} \alpha_{\overline{X}} \circ \overline{r}_{F_{1}X}$$

$$\overline{r}_{F_{2}X}^{*} \circ \phi^{*} \underline{\mathfrak{Q}} \operatorname{id}_{F_{2}\overline{X}} \circ \operatorname{id}_{F_{1}X} \underline{\mathfrak{Q}} \alpha_{\overline{X}} \circ \overline{r}_{F_{1}X}$$

$$r_{F_{1}X}^{*} \circ \operatorname{id}_{F_{1}X} \underline{\mathfrak{Q}} \overline{\phi} \circ \operatorname{id}_{F_{1}X} \underline{\mathfrak{Q}} \alpha_{\overline{X}} \circ \overline{r}_{F_{1}X} \quad (cf. 9.10)$$

$$= \overline{r}_{F_{1}X}^{*} \circ \operatorname{id}_{F_{1}X} \underline{\mathfrak{Q}} (\overline{\phi} \circ \alpha_{\overline{X}}) \circ \overline{r}_{F_{1}X}.$$

I.e.:

But

$$\operatorname{tr}_{\mathcal{F}_{1}\overline{X}} \stackrel{(\overline{\phi} \circ \alpha)}{=} = \overline{r}_{\mathcal{F}_{2}X} \circ \phi^{*} \underline{\mathfrak{B}} \alpha \circ \overline{r}_{\mathcal{F}_{1}X}.$$

Proceeding, write

=

.....

$$\vec{r}_{\vec{F}_{2}X} \circ \phi^{\star} \underline{\Theta} \alpha_{\overline{X}} \circ \vec{r}_{\vec{F}_{1}X}$$

$$= \vec{r}_{\vec{F}_{2}X} \circ ((\alpha_{X} \circ \alpha_{X}^{-1} \circ \phi^{\star}) \underline{\Theta} \alpha_{\overline{X}} \circ id_{\vec{F}_{1}\overline{X}}) \circ \vec{r}_{\vec{F}_{1}X}$$

$$= \vec{r}_{\vec{F}_{2}X} \circ \alpha_{X} \underline{\Theta} \alpha_{\overline{X}} \circ \alpha_{\overline{X}}^{-1} \circ \phi^{\star} \underline{\Theta} id_{\vec{F}_{1}\overline{X}} \circ \vec{r}_{\vec{F}_{1}X}.$$

We then claim that

$$\bar{r}_{\mathcal{F}_{2}X}^{\star} \circ \alpha_{X} \stackrel{\boldsymbol{a}}{=} \alpha_{\bar{X}} = \bar{r}_{\mathcal{F}_{1}X}^{\star}$$

implying thereby that

$$\operatorname{tr}_{\mathcal{F}_{1}\bar{X}}(\bar{\phi} \circ \alpha) = \operatorname{tr}_{\mathcal{F}_{1}X}(\alpha_{X}^{-1} \circ \phi^{\star})$$

which, when combined with the initial observation, renders the contention of the

lemma manifest. From the commutative diagram

we see that

$$\alpha_{\mathbf{X}} \stackrel{\boldsymbol{\Omega}}{=} \alpha_{\mathbf{\overline{X}}} = (\Xi^2)^{-1} \circ \alpha \qquad \circ \Xi^1$$

and from the commutative diagram



we see that

$$F_2 \overline{r}_X^{\star} \circ \alpha = \alpha_e \circ F_1 \overline{r}_X^{\star}.$$

Recalling now that

$$\overline{\mathbf{r}}_{\mathcal{F}_{1}\mathbf{X}} = (\Xi_{\mathbf{X},\mathbf{\overline{X}}}^{1})^{-1} \circ \mathcal{F}_{1}\overline{\mathbf{r}}_{\mathbf{X}} \circ \xi^{1}$$
(cf. 9.12)
$$\overline{\mathbf{r}}_{\mathcal{F}_{2}\mathbf{X}} = (\Xi_{\mathbf{X},\mathbf{\overline{X}}}^{2})^{-1} \circ \mathcal{F}_{2}\overline{\mathbf{r}}_{\mathbf{X}} \circ \xi^{2}$$

we have

$$\begin{split} \bar{r}_{F_{2}X}^{\star} \circ \alpha_{X} & \stackrel{\text{de}}{=} \alpha_{\overline{X}} \\ &= (\xi^{2})^{-1} \circ \mathcal{F}_{2} \bar{r}_{X}^{\star} \circ \Xi^{2}_{X,\overline{X}} \circ (\Xi^{2}_{X,\overline{X}})^{-1} \circ \alpha_{X} \otimes \overline{X} \circ \Xi^{1}_{X,\overline{X}} \\ &= (\xi^{2})^{-1} \circ \mathcal{F}_{2} \bar{r}_{X}^{\star} \circ \alpha_{X} \otimes \overline{X} \circ \Xi^{1}_{X,\overline{X}} \\ &= (\xi^{2})^{-1} \circ \mathcal{F}_{2} \bar{r}_{X}^{\star} \circ \alpha_{X} \otimes \overline{X} \circ \Xi^{1}_{X,\overline{X}} \\ &= (\xi^{2})^{-1} \circ \alpha_{e} \circ \mathcal{F}_{1} \bar{r}_{X}^{\star} \circ \Xi^{1}_{X,\overline{X}} \\ &= (\xi^{1})^{-1} \circ \mathcal{F}_{1} \bar{r}_{X}^{\star} \circ \Xi^{1}_{X,\overline{X}} \\ &= \bar{r}_{F_{1}X}^{\star}, \end{split}$$

as claimed.

The results embodied in 11.12 and 11.13 finish Step 2 of the program, which leaves Step 3 to be dealt with.

Put

$$A_{\mathcal{F}_1,\mathcal{F}_2} = \bigoplus_{i \in I_{\underline{C}}} \operatorname{Mor}(\mathcal{F}_2 X_i, \mathcal{F}_1 X_i).$$

11.14 LEMMA The linear map

$$\mathsf{Y}:\mathsf{A}_{\mathcal{F}_1,\mathcal{F}_2} \to \mathsf{A}(\mathcal{F}_1,\mathcal{F}_2)$$

that sends

$$\phi_i \in Mor(\mathcal{F}_2 X_i, \mathcal{F}_1 X_i)$$

to $[X_{i},\phi_{i}]$ is an isomorphism of vector spaces.

<u>PROOF</u> Every $A \in A(F_1, F_2)$ is an $[X, \phi]$ (cf. 11.8) and every $[X, \phi]$ is a sum of elements $[X_i, \phi_i]$ with X_i irreducible. Therefore 4 is surjective. That 4 is injective is a consequence of the fact that

$$i \neq j \Rightarrow Mor(X_i, X_j) = \{0\}.$$

Put

$$A_{i} = \Psi(Mor(\mathcal{F}_{2}X_{i},\mathcal{F}_{1}X_{i})).$$

Then there is a direct sum decomposition

$$A(\mathcal{F}_{1},\mathcal{F}_{2}) = \bigoplus_{i \in I_{\underline{C}}} A_{i}.$$

Define a linear functional

$$\omega: \mathbb{A}(\mathcal{F}_1, \mathcal{F}_2) \ \rightarrow \ \mathbb{C}$$

by taking it to be zero on ${\rm A}_{\rm i}$ if i does not correspond to e but on ${\rm A}_{\rm e},$ let

$$\omega([e,\phi]) = (\xi^1)^{-1} \circ \phi \circ \xi^2 \in C.$$

11.15 <u>LEMMA</u> $\forall A \neq 0, \omega(A^*A) > 0.$

PROOF Write

$$A = \sum_{i} [X_{i}, \phi_{i}],$$

where the $\boldsymbol{X}_{\underline{i}}$ are irreducible and distinct -- then

$$i \neq j \Rightarrow \omega([X_i,\phi_i] * \cdot [X_j,\phi_j]) = 0.$$

In fact,

$$Mor(e, \overline{X}_{i} \otimes X_{j}) \approx Mor(X_{i}, X_{j}) = \{0\} \quad (cf. 9.3),$$

so e is not a subobject of $\bar{X}_{j} \otimes X_{j}$. One can therefore assume that $A = [X, \phi] \neq 0$ with X irreducible. Recall now that

$$r_X^* \circ r_X = \dim X = n_X id_e \quad (n_X \in \mathbb{N}).$$

This said, let

$$\mathbf{p}_{\mathbf{X}} = \frac{\mathbf{r}_{\mathbf{X}} \circ \mathbf{r}_{\mathbf{X}}^{\star}}{\mathbf{n}_{\mathbf{X}}} \, .$$

Then $p_X^{\star} = p_X^{\bullet}$ and

$$p_X \circ p_X = \frac{r_X \circ r_X^*}{n_X} \circ \frac{r_X \circ r_X^*}{n_X}$$
$$= \frac{1}{n_X^2} (r_X \circ r_X^* \circ r_X \circ r_X^*)$$
$$= \frac{1}{n_X^2} r_X \circ n_X id_e \circ r_X^*$$
$$= \frac{r_X \circ r_X^*}{n_X} = p_X.$$

I.e.:

$$p_{X} \in Mor(\bar{X} \boxtimes X, \bar{X} \boxtimes X)$$

is a projection. Write

$$A^*A = [X,\phi]^* \cdot [X,\phi]$$

$$= [\overline{X},\overline{\phi}] \cdot [X,\phi]$$

$$= [\overline{X} \otimes X,\Xi^{1} \circ \overline{\phi} \otimes \phi \circ (\Xi^{2})^{-1}]$$

$$= [\overline{X} \otimes X,\mathcal{F}_{1}(p_{X}) \circ \Xi^{1} \circ \overline{\phi} \otimes \phi \circ (\Xi^{2})^{-1}]$$

$$+ [\overline{X} \ \underline{\otimes} \ X, \mathcal{F}_{1}(\operatorname{id}_{\overline{X} \ \underline{\otimes} \ X}^{-} - \mathbf{p}_{X}^{-}) \circ \overline{\mathbb{I}}_{\overline{X}, X}^{1} \circ \overline{\Phi} \ \underline{\otimes} \ \phi \circ (\overline{\mathbb{I}}_{\overline{X}, X}^{2})^{-1}]$$

$$= [\overline{X} \ \underline{\otimes} \ X, \mathcal{F}_{1}(\mathbf{p}_{X}^{-}) \circ \overline{\mathbb{I}}_{\overline{X}, X}^{1} \circ \overline{\Phi} \ \underline{\otimes} \ \phi \circ (\overline{\mathbb{I}}_{\overline{X}, X}^{2})^{-1}]$$

$$= \frac{1}{n_{X}} [\mathbf{e}, \mathcal{F}_{1} \mathbf{r}_{X}^{*} \circ (\overline{\mathbb{I}}_{\overline{X}, X}^{1} \circ \overline{\Phi} \ \underline{\otimes} \ \phi \circ (\overline{\mathbb{I}}_{\overline{X}, X}^{2})^{-1} \circ \mathcal{F}_{2} \mathbf{r}_{X}^{-1}]$$

$$= \frac{1}{n_{X}} ((\xi^{1})^{-1} \circ \mathcal{F}_{1} \mathbf{r}_{X}^{*} \circ (\overline{\mathbb{I}}_{\overline{X}, X}^{1} \circ \overline{\Phi} \ \underline{\otimes} \ \phi \circ (\overline{\mathbb{I}}_{\overline{X}, X}^{2})^{-1}) \circ \mathcal{F}_{2} \mathbf{r}_{X} \circ \xi^{2}) [\mathbf{e}, \xi^{1} \circ (\xi^{2})^{-1}]$$

$$= \frac{1}{n_{X}} (\mathbf{r}_{\overline{\mathcal{F}}_{1} X}^{*} \circ \overline{\Phi} \ \underline{\otimes} \ \phi \circ \mathbf{r}_{\overline{\mathcal{F}}_{2} X}) [\mathbf{e}, \xi^{1} \circ (\xi^{2})^{-1}]$$

$$= \frac{1}{n_{X}} (\mathbf{r}_{\overline{\mathcal{F}}_{1} X}^{*} \circ \operatorname{id}_{\mathcal{F}_{1} \overline{X}} \ \underline{\otimes} \ (\phi \circ \phi^{*}) \circ \mathbf{r}_{\overline{\mathcal{F}}_{1} X}) [\mathbf{e}, \xi^{1} \circ (\xi^{2})^{-1}]$$

$$= \frac{1}{n_{X}} (\phi^{*} \circ \phi) [\mathbf{e}, \xi^{1} \circ (\xi^{2})^{-1}],$$

where

$$\Phi = \operatorname{id}_{\mathcal{F}_1 \overline{X}} \underline{\otimes} \Phi^* \circ r_{\mathcal{F}_1 X}.$$

Then

when viewed as a constant, is nonnegative. But $\phi \neq 0 \Rightarrow \phi \neq 0$. Proof: $\overline{\phi}$ is the unique element of Mor $(\mathcal{F}_2 \overline{X}, \mathcal{F}_1 \overline{X})$ such that

$$\overline{\phi} \stackrel{\text{a}}{=} \operatorname{id}_{\mathcal{F}_2 X} \circ r_{\mathcal{F}_2 X} = \operatorname{id}_{\mathcal{F}_1 \overline{X}} \stackrel{\text{a}}{=} \phi^* \circ r_{\mathcal{F}_1 X} \quad (\text{cf. 9.10}),$$

so $\phi = 0 \Rightarrow \overline{\phi} = 0$

$$=> \ ({}^{\vee} \phi) \star = 0 \ => \ ({}^{\vee} \phi) \star \star = 0 \ => \ {}^{\vee} \phi = 0 \ => \ \phi = 0.$$

[Note: To justify the equation

 $\overline{\phi} \stackrel{\otimes}{=} \phi \circ r_{\mathcal{F}_{2}X} \stackrel{=}{=} \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \stackrel{\otimes}{=} (\phi \circ \phi^{\star}) \circ r_{\mathcal{F}_{1}X'}$

write

$$\overline{\phi} \stackrel{\otimes}{=} \phi = \operatorname{id}_{\mathcal{F}_1 \overline{X}} \stackrel{\otimes}{=} \phi \circ \overline{\phi} \stackrel{\otimes}{=} \operatorname{id}_{\mathcal{F}_2 X}.$$

Then

$$\overline{\phi} \ \underline{\Theta} \ \phi \ \circ \ r_{\mathcal{F}_{2}X} = \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\Theta} \ \phi \ \circ \ \overline{\phi} \ \underline{\Theta} \ \operatorname{id}_{\mathcal{F}_{2}X} \ \circ \ r_{\mathcal{F}_{2}X}$$

$$= \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\Theta} \ \phi \ \circ \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\Theta} \ \phi^{*} \ \circ \ r_{\mathcal{F}_{1}X}$$

$$= \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \circ \ \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\Theta} \ (\phi \ \circ \ \phi^{*}) \ \circ \ r_{\mathcal{F}_{1}X}$$

$$= \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\Theta} \ (\phi \ \circ \ \phi^{*}) \ \circ \ r_{\mathcal{F}_{1}X}$$

$$= \operatorname{id}_{\mathcal{F}_{1}\overline{X}} \ \underline{\Theta} \ (\phi \ \circ \ \phi^{*}) \ \circ \ r_{\mathcal{F}_{1}X}$$

Given A, B $\in A(\mathcal{F}_1,\mathcal{F}_2)$, let

$$\langle A,B\rangle = \omega(A^*B)$$
.

Then < , > equips $A(F_1, F_2)$ with the structure of a pre-Hilbert space w.r.t. which the left multiplication operators

$$\mathbb{A}(\mathcal{F}_1,\mathcal{F}_2) \to \mathbb{A}(\mathcal{F}_1,\mathcal{F}_2)$$

are continuous. Denoting by $H(F_1,F_2)$ the Hilbert space completion of $A(F_1,F_2)$, it thus follows that $A(F_1,F_2)$ admits a faithful *-representation

$$L:A(\mathcal{F}_1,\mathcal{F}_2) \rightarrow B(\mathcal{H}(\mathcal{F}_1,\mathcal{F}_2)),$$

hence $A(\mathcal{F}_1, \mathcal{F}_2)$ admits a C*-norm as claimed in Step 3.

§12. THE INTRINSIC GROUP

Let <u>C</u> be a tannakian category and suppose that

$$\mathcal{F}: \underline{C} \rightarrow FDHILB$$

is a fiber functor -- then its intrinsic group G_F is the group of unitary monoidal natural transformations $\alpha:F \rightarrow F$, i.e., in the notation of §11,

$$G_{\mathbf{F}} = Mor(\mathcal{F}, \mathcal{F}),$$

where Mor(F,F) is computed in $ff^*(\underline{C})$.

So

$$G_{\mathcal{F}} \subset \prod_{X \in Ob \subseteq} u(\mathcal{F}X),$$

 $\mathcal{U}(FX)$ the compact group of unitary operators $FX \to FX.$ And G_F is closed if

is equipped with the product topology, thus $\boldsymbol{G}_{\boldsymbol{F}}$ is a compact group.

N.B. Define

$$\pi_{X}^{*G}_{F} \rightarrow U(FX)$$

by $\pi_X(\alpha) = \alpha_X$ -- then

$$(\pi_X, FX) \in Ob \underline{\operatorname{Rep}} G_F.$$

12.1 LEMMA \exists a faithful symmetric monoidal *-preserving functor $\Phi: \underline{C} \neq \underline{\text{Rep}} \ G_{F}$ such that $U \circ \Phi = F$, where

$$U: \operatorname{Rep} G_{F} \rightarrow \operatorname{FDHILB}$$

is the forgetful functor.

PROOF Define Φ on objects by

$$\Phi X = (\pi_v, FX)$$

and on morphisms $f:X \to Y$ by $\Phi f = Ff$ (cf. 8.2) and take for ξ , Ξ the corresponding entities per F. To see that this makes sense for Ξ say, one must check that $\Xi_{X,Y}$ is a morphism in <u>Rep</u> G_F , viz.:

$$\Xi_{X,Y} \circ (\pi_X(\alpha) \stackrel{\text{\tiny def}}{=} \pi_Y(\alpha)) = \pi_X \stackrel{\text{\tiny def}}{=} Y^{(\alpha)} \circ \Xi_{X,Y}$$

But this is obvious since the diagram



commutes. That Φ is symmetric is equally clear.

More is true: Φ is an equivalence of categories. Because Φ is faithful, it remains to establish that Φ is full and has a representative image (details below).

12.2 <u>REMARK</u> The category $\underline{\operatorname{Rep}}_{\mathrm{fd}} \stackrel{A}{}_{\mathcal{F}}$ is a semisimple symmetric monoidal *-category which can be shown to have conjugates, thus $\underline{\operatorname{Rep}}_{\mathrm{fd}} \stackrel{A}{}_{\mathcal{F}}$ is "almost" tannakian. Specializing 8.14, it was pointed out in 8.16 that the " Φ " defined there is a symmetric monoidal equivalence $\underline{C} \rightarrow \underline{\operatorname{Rep}}_{\mathrm{fd}} \stackrel{A}{}_{\mathcal{F}}$. Denote now by $\underline{\operatorname{Rep}}_{\mathrm{fd}} \stackrel{G}{}_{\mathcal{F}}$ the category whose objects are the finite dimensional continuous representations
of $G_{\mathcal{F}}$ and whose morphisms are the intertwining operators — then the inclusion functor

$$\underline{\operatorname{Rep}} \; \operatorname{G}_{F} \rightarrow \underline{\operatorname{Rep}}_{\mathrm{fd}} \; \operatorname{G}_{F}$$

is an equivalence. On the other hand, there is a canonical functor

$$\underline{\operatorname{Rep}}_{\mathrm{fd}} \overset{A}{\mathcal{F}} \stackrel{*}{\rightarrow} \underline{\operatorname{Rep}}_{\mathrm{fd}} \overset{G}{\mathcal{F}}$$

and it too is an equivalence (a nontrivial fact).

12.3 <u>LEMMA</u> If $X \in Ob \subseteq$ is irreducible, then the complex linear span of the $\pi_X(\alpha)$ ($\alpha \in G_F$) is dense in B(FX).

12.4 <u>LEMMA</u> If X,Y \in Ob <u>C</u> are irreducible and nonisomorphic, then the complex linear span of the $\pi_X(\alpha) \oplus \pi_Y(\alpha)$ ($\alpha \in G_F$) is dense in B(FX) \oplus B(FY).

12.5 <u>REMARK</u> If X_1, \ldots, X_n are distinct elements of $I_{\underline{C}}$, then the complex linear span of the

$$\pi_{X_{1}}^{(\alpha)} \oplus \cdots \oplus \pi_{X_{n}}^{(\alpha)} \quad (\alpha \in G_{\mathcal{F}})$$

is dense in

$$B(FX_1) \oplus \cdots \oplus B(FX_n).$$

To prove that ϕ is full, we shall appeal to 7.9.

(a) X irreducible => ϕ X irreducible. In fact, thanks to 12.3, the only T \in B(FX) that intertwine the $\pi_X(\alpha)$ ($\alpha \in G_F$) are the scalar multiples of the identity.

(b) X,Y irreducible and nonisomorphic => $\Phi X, \Phi Y$ irreducible and nonisomorphic.

For suppose that $T:FX \to FY$ intertwines π_X and π_Y , thus $T\pi_X(\alpha) = \pi_Y(\alpha)T$ ($\alpha \in G_F$). But then Tu = vT for all $u \in B(FX)$, $v \in B(FY)$ (cf. 12.4). Now take u = 0, v = 1 to conclude that T = 0, hence $\Phi X, \Phi Y$ are nonisomorphic.

The final claim is that Φ has a representative image. To see this, consider the map

$$\gamma_{\mathcal{F}}: \mathbf{I}_{\underline{\mathbf{C}}} \to \mathbf{I}_{\underline{\operatorname{Rep}}} \operatorname{G}_{\mathcal{F}}$$

defined by the rule

$$\gamma_{\mathcal{F}}(X_{\mathbf{i}}) = (\pi_{X_{\mathbf{i}}}, \mathcal{F}X_{\mathbf{i}}).$$

Then $\gamma_{\mathbf{r}}$ is injective.

12.6 LEMMA γ_F is surjective.

<u>PROOF</u> The complex linear span of the matrix elements of the π_{X_1} as i ranges over $I_{\underline{C}}$ is a unital *-subalgebra of $C(G_{\underline{F}})$ which separates the points of $G_{\underline{F}}$, thus is dense in $C(G_{\underline{F}})$. Accordingly, there can be no irreducible object in <u>Rep</u> $G_{\underline{F}}$ which is not unitarily equivalent to a π_{X_1} for some i, so $\gamma_{\underline{F}}$ is surjective.

Therefore γ_{F} is bijective and Φ has a representative image.

12.7 REMARK Suppose that

$$F_1: \underline{C} \to \underline{FDHILB}$$

$$F_2: \underline{C} \to \underline{FDHILB}$$

are fiber functors — then as objects of $ff^*(\underline{C}), \mathcal{F}_1, \mathcal{F}_2$ are isomorphic (cf. 11.3),

so G_{F_1}, G_{F_2} are isomorphic (in the category <u>CPTGRP</u>).

Let G be a compact group -- then the forgetful functor

$$U:\underline{Rep} \ G \rightarrow \underline{FDHILB}$$

is a fiber functor. Define a map $\Gamma\colon G \twoheadrightarrow G_U$ by sending $\sigma \in G$ to the string

$$\{\pi(\sigma): (\pi, \mathcal{H}_{\pi}) \in Ob \ \underline{\operatorname{Rep}} \ G\}.$$

That this is meaningful follows upon noting that if

$$(\pi_{1}, \mathcal{H}_{\pi_{1}}) \in Ob \underline{\text{Rep }} G,$$
$$(\pi_{2}, \mathcal{H}_{\pi_{2}})$$

then

$$\forall T \in Mor((\pi_1, H_{\pi_1}), (\pi_2, H_{\pi_2}))$$

there is a commutative diagram



thus the string

$$\{\pi(\sigma):(\pi, \mathcal{H}_{\pi}) \in \mathsf{Ob} \ \underline{\mathsf{Rep}} \ \mathsf{G}\}$$

defines an element

$$\alpha(\sigma) \in Mor(U,U)$$

where technically

$$\alpha(\sigma)_{(\pi,H_{\pi})} = \pi(\sigma),$$

12.8 LEMMA f is a continuous injective homomorphism. [This is immediate from the definitions.]

In fact, Γ is surjective, hence G and G_{II} are isomorphic.

[If I were not surjective, replace G by FG and think of G as a proper closed subgroup of G_U — then there would be an irreducible representation of G_U that contains a nonzero vector invariant under G but not under G_U . This, however, is impossible:

$$\gamma_{U}: I_{\underline{\operatorname{Rep}}} G \xrightarrow{\rightarrow} I_{\underline{\operatorname{Rep}}} G_{U}$$

is bijective.]

12.9 <u>THEOREM</u> Up to isomorphism in <u>CPTGRP</u>, G is the "intrinsic group" of <u>Rep</u> G.

[If

$F: \operatorname{Rep} G \rightarrow \operatorname{FDHILB}$

is a fiber functor, then $G_{f} \approx G_{U}$ (cf. 12.7).]

12.10 <u>REMARK</u> Compact groups G,G' are said to be <u>isocategorical</u> if <u>Rep</u> G, <u>Rep</u> G' are equivalent as monoidal categories. In general, this does not mean that <u>Rep</u> G, <u>Rep</u> G' are equivalent as symmetric monoidal categories and G,G' may very well be isocategorical but not isomorphic.

§13. CLASSICAL THEORY

A <u>character</u> of a commutative unital C*-algebra A is a nonzero homomorphism $\omega: A \rightarrow C$ of algebras. The set of all characters of A is called the <u>structure space</u> of A and is denoted by $\Delta(A)$.

N.B. We have

$$\Delta(A) = \emptyset \quad (A = \{0\})$$
$$\Delta(A) \neq \emptyset \quad (A \neq \{0\}).$$

13.1 LEMMA Let $\omega \in \Delta(A)$ — then ω is necessarily bounded. In fact,

$$||\omega|| = 1 = \omega(1_A).$$

<u>N.B.</u> The elements of $\Delta(A)$ are the pure states of A, hence, in particular, are *-homomorphisms: $\forall A \in A$,

$$\omega(\mathbf{A^*}) = \overline{\omega(\mathbf{A})}.$$

Given $A \in A$, define

$$A: \Delta(A) \rightarrow C$$

by

$$\widehat{A}(\omega) = \omega(A)$$

Equip $\Delta(A)$ with the initial topology determined by the A, i.e., equip $\Delta(A)$ with the relativised weak* topology.

13.2 LEMMA $\Delta(A)$ is a compact Hausdorff space.

If X is a compact Hausdorff space, then C(X) equipped with the supremum norm

$$||\mathbf{f}|| = \sup_{\mathbf{x} \in \mathbf{X}} |\mathbf{f}(\mathbf{x})|$$

and involution

$$f^*(x) = \overline{f(x)}$$

is a commutative unital C*-algebra. Moreover, $\forall \ x \in X$, the Dirac measure $\delta_x \in \Delta(C(X)) \text{ and the arrow}$

$$\begin{bmatrix} X + \Delta(C(X)) \\ x + \delta_{x} \end{bmatrix}$$

is a homeomorphism.

13.3 LEMMA $\hat{A} \in C(\Delta(A))$ and the arrow

$$\begin{bmatrix} A \rightarrow C(\Delta(A)) \\ A \rightarrow \hat{A} \end{bmatrix}$$

is a unital *-isomorphism.

<u>N.B.</u> If $A = \{0\}$, then $\Delta(A) = \emptyset$ and there is exactly one map $\emptyset \neq C$, namely the empty function $(\emptyset = \emptyset \times C)$, which we shall take to be 0.

Notation: Let <u>CPTSP</u> be the category whose objects are the compact Hausdorff spaces and whose morphisms are the continuous functions.

Notation: Let <u>COMUNC*ALG</u> be the category whose objects are the commutative unital C*-algebras and whose morphisms are the unital *-homomorphisms.

Let X and Y be compact Hausdorff spaces. Suppose that $\phi: X \rightarrow Y$ is a continuous function -- then ϕ induces a unital *-homomorphism

$$\phi^*: C(Y) \rightarrow C(X),$$

3.

viz. $\phi^*(f) = f \circ \phi$. Therefore the association that sends X to C(X) defines a cofunctor

$$C:\underline{CPTSP} \rightarrow \underline{COMUNC*ALG}.$$

Let A and B be commutative unital C*-algebras. Suppose that $\Phi: A \rightarrow B$ is a unital *-homomorphism --- then Φ induces a continuous function

$$\Phi^*: \Delta(\mathcal{B}) \to \Delta(\mathcal{A}),$$

viz. $\Phi^*(\omega) = \omega \circ \Phi$. Therefore the association that sends A to $\Delta(A)$ defines a cofunctor

$$\triangle$$
:COMUNC*ALG \rightarrow CPTSP.

13.4 <u>THEOREM</u> The category <u>CPTSP</u> is coequivalent to the category <u>COMUNC*ALG</u>. PROOF Define

$$\Xi_X: X \to \Delta(C(X))$$

by the rule $\Xi_X(x)=\delta_x$ — then Ξ_X is a homeomorphism and there is a commutative diagram



Define

 $\Xi_A: A \ \neq \operatorname{C}(\bigtriangleup(A))$

by the rule $E_A(A) = A -$ then E_A is a unital *-isomorphism and there is a commutative

diagram



Therefore

 $\begin{bmatrix} - & \text{id} \approx \Delta \circ \mathbf{C} \\ & \text{id} \approx \mathbf{C} \circ \Delta. \end{bmatrix}$

The category <u>CPTSP</u> has finite products with final object {*}. Therefore the category <u>COMUNC*ALG</u> has finite coproducts with initial object C. To explicate the latter, invoke the nuclearity of the objects of COMUNC*ALG, thus

$$A \bigotimes_{\max} B = A \bigotimes_{\min} B,$$

call it A 🛛 B -- then

and there are arrows

$$\begin{bmatrix} A \rightarrow A & B \\ A \rightarrow A & B \end{bmatrix} \begin{bmatrix} B \rightarrow A & B \\ B \rightarrow 1_{A} & B \\ B \rightarrow 1_{A} & B \end{bmatrix}$$

13.5 EXAMPLE We have

$$\begin{bmatrix} - & C(\{*\}) \approx C & \text{and} & C(X \times Y) \approx C(X) & \underline{\Theta} & C(Y) \\ & \Delta(C) \approx \{*\} & \text{and} & \Delta(A & \underline{\Theta} & B) \approx \Delta(A) \times \Delta(B) \\ \end{bmatrix}$$

13.6 <u>REMARK</u> Let A be a commutative unital C*-algebra — then the algebraic tensor product A @ A can be viewed as an involutive subalgebra of A @ A. Another point is this: Since A @ A is the coproduct, there is a canonical arrow $A @ A \xrightarrow{m} A$ with m(A @ B) = AB, i.e., the restriction of m to A @ A is the multiplication in A.

[Note: If A_1, A_2, B are commutative unital C*-algebras and if

$$\begin{bmatrix} \Phi_1 : A_1 \to B \\ \Phi_2 : A_2 \to B \end{bmatrix}$$

are unital *-homomorphisms, then the diagram



admits a unique filler

$$\Phi_1 \stackrel{\otimes}{=} \Phi_2 : A_1 \stackrel{\otimes}{=} A_2 \stackrel{*}{\to} B$$

such that

$$(\Phi_1 \ \underline{\otimes} \ \Phi_2) (A_1 \ \underline{\otimes} \ A_2) = \Phi_1(A_1) \Phi_2(A_2) \quad (A_1 \in A_1, \ A_2 \in A_2).]$$

13.7 <u>RAPPEL</u> Let <u>C</u> be a category with finite products and final object T -then a group object in <u>C</u> consists of an object G and morphisms

 $\mu: G \times G \rightarrow G, \eta: T \rightarrow G, \iota: G \rightarrow G$

such that the following diagrams commute:





There are obvious definitions of internal group homomorphism $G \neq G'$, composition of internal group homomorphisms $G \neq G'$, $G' \neq G''$, and the identity internal group homomorphism $\mathrm{id}_{G}: G \neq G$. Accordingly, there is a category $\underline{\mathrm{GRP}}(\underline{C})$ whose objects are the group objects in \underline{C} and whose morphisms are the internal group homomorphisms.

[Note: If instead <u>C</u> is a category with finite coproducts and initial object I, then we put

$$\underline{\text{COGRP}}(\underline{C}) = \underline{\text{GRP}}(\underline{C}^{\text{OP}})^{\text{OP}}$$

and call the objects the cogroup objects in <u>C</u> and the morphisms the internal cogroup homomorphisms.]

13.8 EXAMPLE Take
$$\underline{C} = \underline{SET}$$
 — then

$$GRP(SET) = GRP$$
.

13.9 LEMMA We have

$$GRP(CPTSP) = CPTGRP$$
.

13.10 REMARK The forgetful functor

CPTGRP
$$\rightarrow$$
 SET

has a left adjoint. Proof: Given a set X, equip it with the discrete topology, form the associated free topological group $F_{\rm gr}(X)$, and consider its Bohr compact-ification.

A <u>commutative Hopf C*-algebra</u> is commutative unital C*-algebra H together with unital *-homomorphisms

$$\Delta: H \rightarrow H \otimes H, \epsilon: H \rightarrow C, S: H \rightarrow H$$

for which the following diagrams commute:





[Note: Such an H is not necessarily a Hopf algebra (in general, \triangle takes values in H $\underline{0}$ H rather than H $\underline{0}$ H).]

<u>N.B.</u> Consider, e.g., $(id_{H},S) \rightarrow then in terms of the coproduct diagram$

$$\begin{array}{c} \operatorname{in}_{1} & \operatorname{in}_{2} \\ \operatorname{H} \longrightarrow \operatorname{H} \underline{\otimes} \operatorname{H} \longleftarrow \underline{-} \operatorname{H}, \end{array}$$

the arrow

$$(\mathrm{id}_{\mathrm{H}}, \mathrm{S}) : \mathrm{H} \underline{\mathrm{Q}} \mathrm{H} \to \mathrm{H}$$

is characterized by the condition that

$$(id_{H},S) \circ in_{1} = id_{H}$$
$$(id_{H},S) \circ in_{2} = S.$$

On the other hand, there is an arrow

$$\operatorname{id}_{\operatorname{H}} \underline{\mathfrak{Q}} \operatorname{S:} \operatorname{H} \underline{\mathfrak{Q}} \operatorname{H} \to \operatorname{H} \underline{\mathfrak{Q}} \operatorname{H}$$

characterized by the condition that

$$id_{H} \otimes S \circ in_{1} = in_{1} \circ id_{H}$$
(cf. 13.6).
$$id_{H} \otimes S \circ in_{2} = in_{2} \circ S$$

And

$$m \circ id_H \underline{\Theta} S = (id_H, S).$$

Proof:

$$\begin{array}{|c|c|c|c|c|c|c|c|} \hline m & \circ & \operatorname{id}_{H} & \underline{\mathbf{G}} & S & \circ & \operatorname{in}_{1} & = & \operatorname{id}_{1} & \circ & \operatorname{id}_{H} & = & \operatorname{id}_{H} \\ \hline m & \circ & \operatorname{id}_{H} & \underline{\mathbf{G}} & S & \circ & \operatorname{in}_{2} & = & \operatorname{m} & \circ & \operatorname{in}_{2} & \circ & S & = & \operatorname{id}_{H} & \circ & S & = & S. \end{array}$$

Denote by <u>COMHOPFC*ALG</u> the category whose objects are the commutative Hopf C*-algebras and whose morphisms $f: H \to H'$ are the unital *-homomorphisms such that $f \otimes f \circ \Delta = \Delta' \circ f, \epsilon = \epsilon' \circ f, f \circ S = S' \circ f.$

13.11 LEMMA We have

____

COGRP(COMUNC*ALG) = COMHOPFC*ALG.

Let G be a compact group — then the group operations in G induce operations Δ , ε , S in C(G) w.r.t. which C(G) acquires the structure of a commutative Hopf C*-algebra. And the association that sends G to C(G) defines a cofunctor

C:CPTGRP \rightarrow COMHOPFC*ALG.

Let H be a commutative Hopf C*-algebra -- then the cogroup operations in H induce operations μ , η , ι in $\Delta(H)$ w.r.t. which $\Delta(H)$ acquires the structure of a compact group. And the association that sends H to $\Delta(H)$ defines a cofunctor

\triangle :COMHOPFC*ALG \rightarrow CPTGRP.

13.12 <u>THEOREM</u> The category <u>CPTGRP</u> is coequivalent to the category COMHOPFC*ALG (cf. 13.4).

13.13 <u>RAPPEL</u> Given a compact group G, let A(G) be its set of representative functions -- then A(G) is a unital *-subalgebra of C(G) and when endowed with the

restrictions of Δ , ϵ , S forms a commutative Hopf *-algebra.

[Note: Recall that A(G) is dense in C(G).]

* Let $\Delta\left(A\left(G\right)\right)$ be the set of nonzero multiplicative linear functionals on $A\left(G\right)$.

• Let $\Delta \star (A(G))$ be the set of $\star\text{-preserving}$ nonzero multiplicative linear functionals on A(G) .

Then

$$\Delta^{\star}(A(G)) \subset \Delta(A(G))$$

and the containment is proper in general.

Equip $\Delta(A(G))$ (and hence $\Delta^*(A(G))$) with the topology of pointwise convergence and introduce the following operations:

(i)
$$(\omega_1 \cdot \omega_2) = (\omega_1 \ \underline{Q} \ \omega_2) \circ \Delta;$$
 (ii) $1_{A(G)} = \varepsilon;$ (iii) $\omega^{-1} = \omega \circ S.$

Then $\Delta(A(G))$ is a group containing $\Delta^*(A(G))$ as a subgroup (in this connection, note that $\Delta(f^*) = \Delta(f)^*$ and $S(f^*) = S(f)^*$).

13.14 LEMMA $\Delta^*(A(G))$ is a compact group.

13.15 THEOREM Define

by

$$ev(\sigma) = \delta_{\sigma} (\delta_{\sigma}(f) = f(\sigma)).$$

Then ev is an isomorphism in CPTGRP.

Let

U:Rep G
$$\rightarrow$$
 FDHILB

be the forgetful functor.

13.16 LEMMA The arrow

$$\rho: A(U,U) \rightarrow A(G)$$

that sends $[H_{\pi},\phi]~(\phi:H_{\pi} \not \to H_{\pi})$ to the representative function

$$\sigma \rightarrow \operatorname{tr}(\pi(\sigma)\phi) \quad (\sigma \in G)$$

is a linear bijection.

[Note: This can be sharpened in that A(U,U) carries a canonical Hopf algebra structure which is preserved by ρ , i.e., ρ is an isomorphism of Hopf algebras.]