

Quantum Programming via Linear Homotopy Types

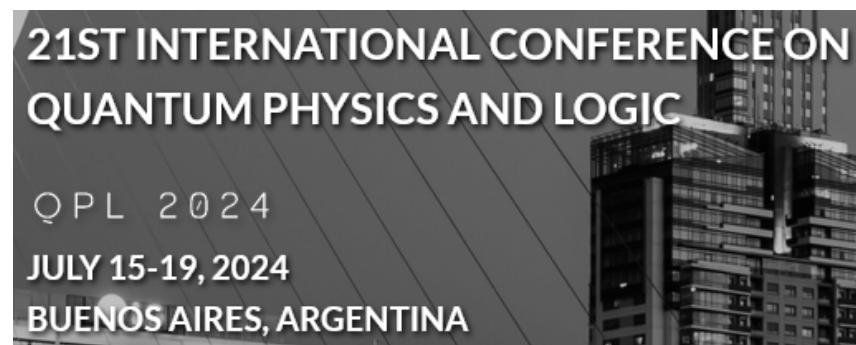
Urs Schreiber



CENTER FOR
QUANTUM &
TOPOLOGICAL
SYSTEMS

جامعة نيويورك أبوظبي
NYU ABU DHABI

talk at:



The Gist

classical logic

lcc category
dependent types

∞ -topos
homotopy types

quantum logic

traditionally:

tensor category
linear types

tensor ∞ -category

homotopify

unify!

tangent ∞ -topos
linear homotopy types

“Quantization via Linear Homotopy Types”
[arXiv:1402.7041]
“Entanglement of Sections”
[arXiv:2309.07245]

Riley: “A bunched Homotopy Type Theory
for Synthetic Stable Homotopy Theory”
[doi:10.14418/wes01.3.139]

“Topological Quantum Gates
in Homotopy Type Theory”
[arXiv:2303.02382]

topological gates
type transport

certification
identity types

expresses
expresses

expresses

quantum effects
motivic yoga

“The Quantum Monadology”
[arXiv:2310.15735]
“Quantum and Reality”
[arXiv:2311.11035]

Riley: “Linear HoTT as a Quantum Certification Language”
[mvr.hosting.nyu.edu/pubs/translation.pdf]

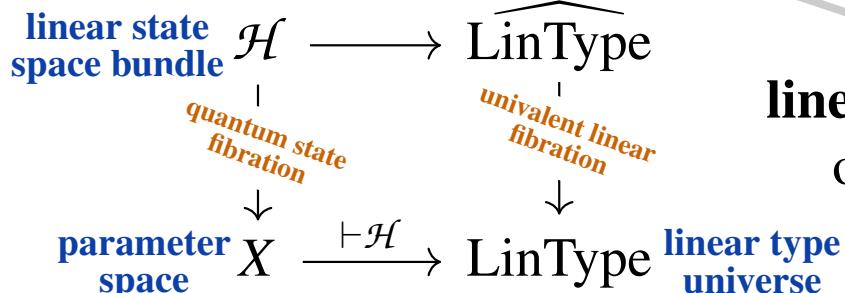
Algebraic Topology

Highlights

local system of state spaces
i.e., flat vector bundle over
classical parameter space

Quantum Physics

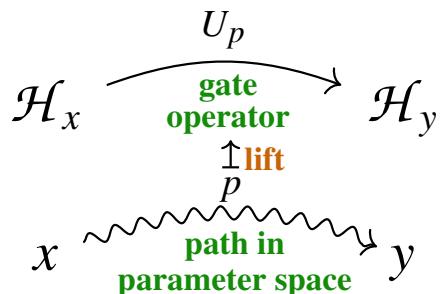
Parameterized Quantum Systems



Linear Homotopy Types

linear types dependent
on homotopy type

Topological Quantum Gates



type transport along
identification in base type

parallel (adiabatic) transport
along path in parameter space

Quantum Measurement

counit of the
descent comonad
("necessity" modality)

$x' : X$
actual
world

$$\vdash \bigoplus_X \mathbb{C} \xrightarrow[\text{state collapse}] {\sum_x q_x |x\rangle \mapsto q_{x'} |x'\rangle} \mathbb{C}$$

$$\square \mathbb{1}_X \xrightarrow[\text{necessity counit}] {\eta^\square} \mathbb{1}_X$$

\downarrow

possible worlds X $\xlongequal{\quad}$ **measurement basis (finite)** X

eliminator of
dependent product
(here: direct sum)

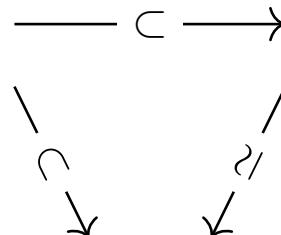
Default Categorical Semantics

Naïve Categorical Semantics:

vector bundles over sets
i.e. over discrete spaces

(coproduct completion of Vect)

Quipper has semantics here



∞ -Categorical Semantics

$H\mathbb{C}$ -module spectra parameterized over ∞ -groupoids

(“tangent ∞ -topos” $T^{\mathbb{C}}\text{Grpd}_{\infty}$)
needs model ctgr presentation

“Quantum Monadology”
§2.1 [arXiv:2310.15735]

Tech glimpse

“Entanglement of Sections”
§3 [arXiv:2309.07245]

Model Categorical Semantics:

simplicial chain complexes parameterized over (simplicial) groupoids
 (“simplicial local systems”)

In all cases: the corresponding Grothendieck construction of tensor categories of linear bundles \mathcal{H} over base spaces X :

$$\underbrace{\int_{X \in \text{ClaType}} \text{QuType}_X}_{\substack{\text{all types are...} \\ \dots \text{linear types over...} \\ \dots \text{classical base types.}}} = \left\{ \begin{array}{ccc} \mathcal{H} & \xrightarrow{\text{fiberwise linear map}} & \mathcal{H}' \\ \downarrow \text{quantum state fibration} & \longrightarrow & \downarrow \text{quantum state fibration} \\ X & \xrightarrow{\text{map of base spaces}} & X' \end{array} \right\}$$

these dependent linear types
are “infinitesimally cohesive”

“Diff. cohomology in cohesive ∞ -topos”
Prop. 4.1.9 [arXiv:1310.7930]

Linear HoTT is foremost HoTT
with infinitesimal cohesive modality \natural

Finster et al: “Synthetic Spectra” [arXiv:2102.04099]

this allows to *construct*
linear/quantum types

$$\text{QuType} \equiv (\mathcal{H} : \text{Type}) \times (\natural \mathcal{H} = *)$$

and a Girard **!-modality** of linear logic
via a Bierman-Benton adjunction

“The Quantum Monadology”
Prop. 2.7 [arXiv:2310.15735]

modality:
“classically”

$$\begin{array}{ccc} \left[\begin{array}{c} 0_X \\ \downarrow \\ X \end{array} \right] & \longleftrightarrow & X \\ \text{Type} & \xrightarrow{\perp \top} & \text{ClaType} \end{array}$$

modality:
“quantumly”

Tech glimpse

$$\begin{array}{ccc} \left[\begin{array}{c} \mathcal{H} \\ \downarrow \\ X \end{array} \right] & \mapsto & X \\ \text{Type} & \xrightarrow{\perp} & \mathcal{H} \\ \left[\begin{array}{c} \mathcal{H} \\ \downarrow \\ * \end{array} \right] & \longleftarrow & \mathcal{H} \\ \left[\begin{array}{c} \mathcal{H} \\ \downarrow \\ X \end{array} \right] & \mapsto & \bigoplus_{x:X} \mathcal{H}_x \\ \text{Type} & \xrightarrow{\top} & \text{QuType} \end{array}$$

“quantized” $Q \equiv \Sigma_+^\infty$

“classicized” $C \equiv \Omega_+^\infty$

$$\begin{array}{ccccc} !\text{QuType} & \xleftarrow{\quad \dashv \quad} & \text{Type} & \xleftarrow{\quad \vdash \quad} & \text{ClaType} \\ \downarrow & \dashv \quad \vdash & & & \uparrow \\ \xleftarrow{\perp} & & \xleftarrow{\perp} & & \xleftarrow{\perp} \\ - \natural(\mathbb{1} \rightarrow (-)) \rightarrow & & & & \end{array}$$

Syntax

Semantics

$$\text{h-FORM } \frac{\Gamma \vdash A : \text{Type}}{\Gamma \vdash \underline{h}A : \text{Type}}$$

$$\frac{}{(5) \underline{h}\Gamma^{\text{rel}}A \longrightarrow \underline{h}(\eta_{\Gamma}^{\natural})^*A - \underline{h}q_A \rightarrow \underline{h}A}{\begin{array}{c} A \\ \downarrow p_A \\ \underline{h}\Gamma \end{array}} \\ \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \\ \Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \underline{h}\eta_{\Gamma}^{\natural} \rightarrow \underline{h}\underline{h}\Gamma$$

$$\text{h-INTRO } \frac{\Gamma \vdash a : A}{\Gamma \vdash a^{\natural} : \underline{h}A}$$

$$\frac{}{\Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \vdash a \rightarrow A}{\begin{array}{c} \vdash a \rightarrow A \\ \parallel \\ \underline{h}\Gamma \end{array}} \\ \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \\ \Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \eta_{\Gamma}^{\natural} \rightarrow \underline{h}\underline{h}\Gamma$$

$$\text{h-ELIM } \frac{\Gamma \vdash b : \underline{h}A}{\Gamma \vdash b^{\natural} : A}$$

$$\frac{}{\Gamma \longrightarrow \vdash b \longrightarrow \underline{h}\Gamma^{\text{rel}}A \rightarrow \underline{h}((\eta_{\Gamma}^{\natural})^*A)}{\begin{array}{c} \vdash b^{\natural} : A \\ \parallel \\ \underline{h}\Gamma \end{array}} \\ \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \\ \Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \eta_{\Gamma}^{\natural} \rightarrow \underline{h}\underline{h}\Gamma$$

$$\frac{}{\Gamma \longrightarrow \vdash b \longrightarrow \underline{h}\Gamma^{\text{rel}}A \rightarrow \underline{h}((\eta_{\Gamma}^{\natural})^*A) \rightarrow \underline{h}A - \varepsilon_A^{\natural} \rightarrow A}{\begin{array}{c} \vdash b^{\natural} : A \\ \parallel \\ \underline{h}\Gamma \end{array}} \\ \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \\ \Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \eta_{\Gamma}^{\natural} \rightarrow \underline{h}\underline{h}\Gamma \xrightarrow{\varepsilon_{\underline{h}\Gamma}^{\natural}} \underline{h}\Gamma$$

$$\frac{(24) \frac{\Gamma \vdash a : A}{\Gamma \vdash \underline{a} : A} \quad (30) \frac{\Gamma \vdash a : A}{\Gamma \vdash \underline{a}^{\natural} : \underline{h}A}}{\Gamma \vdash \underline{a}^{\natural} : \underline{h}A}$$

$$\frac{}{\Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \varepsilon_{\Gamma}^{\natural} \rightarrow \Gamma - \vdash a \rightarrow A - \eta_A^{\natural} \rightarrow \underline{h}A \xrightarrow{(\eta_A^{\natural})^{-1}} \underline{h}\Gamma}{\begin{array}{c} \eta_A^{\natural} \circ (\vdash a) \\ \downarrow \quad \quad \quad \downarrow \\ \Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \eta_A^{\natural} \rightarrow \underline{h}\underline{h}\Gamma \end{array}} \\ \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \\ \Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \eta_{\Gamma}^{\natural} \rightarrow \Gamma - \eta_{\Gamma}^{\natural} \rightarrow \underline{h}\underline{h}\Gamma$$

some inference rules for the *classically modality*

$$\frac{(19) \frac{\Gamma \vdash A : \text{Type}}{\Gamma, a:A \vdash a : A} \quad (34) \frac{\Gamma, a:A \vdash a^{\natural} : A}{\Gamma \vdash (a \mapsto a^{\natural}) : A \rightarrow \underline{h}A}}{\Gamma \vdash (a \mapsto a^{\natural}) : A \rightarrow \underline{h}A}$$

$$\frac{\begin{array}{c} A \longrightarrow \eta_A^{\natural} \longrightarrow \underline{h}A \\ \downarrow p_A \quad \quad \quad \downarrow \\ \underline{h}\Gamma^{\text{rel}}A \end{array}}{\begin{array}{c} \underline{h}\Gamma^{\text{rel}}A \xrightarrow{\alpha_A, q_A^{\natural}} \underline{h}\Gamma \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma \end{array}} \quad \quad \quad \frac{\begin{array}{c} \underline{h}\Gamma^{\text{rel}}A \xrightarrow{\text{pb}} \underline{h}\Gamma \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma \end{array}}{\begin{array}{c} \underline{h}\Gamma - \eta_{\Gamma}^{\natural} \rightarrow \underline{h}\Gamma \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma \end{array}}$$

some inference rules for linear types:

Syntax

Semantics

$$\frac{\Gamma \vdash A : \text{Type} \quad \Gamma \vdash \underline{x} : \underline{h}A}{\Gamma \vdash A_{\underline{x}} := (a : A) \times \text{Id}(a^{\natural}, \underline{x})}$$

linear fiber

$$\frac{\begin{array}{c} A_{\underline{x}} \xrightarrow{\text{pb}} \underline{h}\Gamma \xrightarrow{\vdash \underline{x}} \underline{h}A \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma^{\text{rel}}A \xrightarrow{\text{pb}} \Delta_{\underline{x}}^{\text{rel}}A \end{array}}{\begin{array}{c} \Delta_{\underline{x}}^{\text{rel}}A \xrightarrow{(\vdash \underline{x}, \text{id})} \underline{h}A \times \underline{h}\Gamma \xrightarrow{\text{id} \times \vdash \underline{x}} \underline{h}A \times \underline{h}A \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma - \Delta_{\underline{x}} \rightarrow \underline{h}\Gamma \times \underline{h}\Gamma \end{array}}$$

$$\frac{\Gamma \vdash A : \text{Type} \quad \Gamma \vdash \underline{x} : \underline{h}A}{\Gamma \vdash \underline{h}(A_{\underline{x}}) \simeq *}$$

linear fibers are indeed linear

$$\frac{\begin{array}{c} \underline{h}A_{\underline{x}} \xrightarrow{\sim} \underline{h}\Gamma \xrightarrow{\vdash \underline{x}} \underline{h}A \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma \xrightarrow{\text{pb}} \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\underline{h}\Gamma \end{array}}{\begin{array}{c} \underline{h}\Gamma - \eta_{\Gamma}^{\natural} \rightarrow \underline{h}(A_{\underline{x}}) \xrightarrow{\sim} \underline{h}(\underline{h}A \times \underline{h}\Gamma) \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}A - \underline{h}(\eta_A^{\natural}, \eta_A^{\natural} \circ p_A) \rightarrow \underline{h}(\underline{h}A \times \underline{h}\Gamma) \end{array}}$$

$$\frac{\Gamma \vdash A : \text{Type}}{\Gamma \vdash A \simeq \sum_{x : \underline{h}A} A_x}$$

types are sums of their linear fibers

$$\frac{\begin{array}{c} A \xrightarrow{\text{pb}} \eta_A^{\natural} \longrightarrow \underline{h}A \\ \downarrow p_A \quad \quad \quad \downarrow \\ \underline{h}\Gamma^{\text{rel}}A \xrightarrow{\text{pb}} \Delta_A \end{array}}{\begin{array}{c} \Delta_A \xrightarrow{(\text{id}_A, \eta_A^{\natural})} \underline{h}\Gamma^{\text{rel}}A \xrightarrow{\text{pb}} \underline{h}\Gamma \\ \downarrow \quad \quad \quad \downarrow \\ A \times_{\Gamma} \underline{h}A \xrightarrow{\text{id}_{\underline{h}A}} \underline{h}\Gamma \times_{\Gamma} \underline{h}A \\ \downarrow \quad \quad \quad \downarrow \\ \underline{h}\Gamma^{\text{rel}}A \times_{\Gamma} \underline{h}\Gamma^{\text{rel}}A \xrightarrow{\text{pb}} \underline{h}\Gamma \end{array}} \quad \quad \quad \frac{\Gamma \longrightarrow \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma}{\Gamma \vdash \eta_{\Gamma}^{\natural} \longrightarrow \underline{h}\Gamma}$$

key: **quantum measurement** effect handling
captured by dependent linear type formation



recall: given a context extension

then dependent sum and dependent product form the base change adjoint triple, where...

$$\begin{array}{ccc}
 \text{classical type } X & \xrightarrow{\quad\quad\quad} & *
 \\[10pt]
 \text{---} & \text{dependent sum} & \text{---} \\
 & \xrightarrow{\Sigma_X} & \\
 \text{---} & \perp & \text{---} \\
 \text{X-dependent...} & \xleftarrow{\star_X \times (-)} & \text{...classical types} \\
 \text{ClaType}_X & \xleftarrow{\quad\quad\quad} & \text{ClaType} \\
 & \xrightarrow{\Pi_X} & \\
 & \text{dependent products} &
 \end{array}$$

...the induced (co-)modalities exhibit S5 modal logic

now for dependent *linear* types over finite bases this is the direct sum...



quantum necessity monad

models “dynamic lifting” of quantum measurement;
Kleisli equivalence proves deferred msrmnt principle

$$\begin{array}{ccc}
 \text{---} & \text{direct sum} & \text{---} \\
 & \xrightarrow{\oplus_X} & \\
 \text{X finite:} & \perp & \text{---} \\
 \text{---} & \xleftarrow{1_X \otimes (-)} & \text{...linear types} \\
 \text{---} & \xrightarrow{\oplus_X} & \\
 & \text{direct sum} &
 \end{array}$$

... which yields pair of Frobenius monads

“Quantum Monadology”
§2 [arXiv:2310.15735]

quantum indefiniteness monad
 \simeq Coecke’s measurement monad
 as now used in xz-calculus

this way, quantum-logic & -effects are incarnated in homotopy type theory by a **system of modalities/monads**

- | | |
|---|--|
| axiomatized:
<small>Riley: “A bunched HoTT”
[doi:10.14418/wes01.3.139]</small> | \vdash classically
\otimes entangled
\triangleright quantumly
$!$ 2 nd quantized
\square necessarily
\circlearrowleft indefinitely |
| induced: | |
| <small>“Quantum Monadology”
§2 [arXiv:2310.15735]</small> | |
| | |
| | |
| | |
| | |

via do-notation for monadic effects this permits a domain specific **quantum language** to be embedded into LHoTT

e.g. snippet from embedded error-correction pseudocode:

```
correct_error : LgclQBit → OSyndromeLgclQBit
correct_error ≡
  [for |b1, b2, b3⟩
   do [for |ψ⟩ in measure_syndrome(|b1, b2, b3⟩)
        do [if measured (s1, s2)
             then BitFlip(s1, s2)|ψ⟩]]]
```

“Quantum Monadology”
§3 [arXiv:2310.15735]

quantum measurement effects

1st punchline

“Quantum Monadology”
Lit 1.19 [arXiv:2310.15735] —

corresponding do-notation

```
prog : D1 → E(D3)
prog ≡
  [for d2 in this d1
   do [for d3 in that d2
        do return d3]]]
```

Kleisli morphism for monadic effect E

$$\begin{array}{c}
 D_1 \xrightarrow{\text{id}(D_1 \xrightarrow{\text{this}} E(D_2))} \\
 \hline
 \hookrightarrow E(D_2) \xrightarrow{\text{bind}^E(D_2 \xrightarrow{\text{that}} E(D_3))} \\
 \hline
 \rightarrow E(D_3) \xrightarrow{\text{bind}^E(\text{return}_{D_3}^E)} E(D_3)
 \end{array}$$

beyond offering certification for traditional quantum languages (e.g. Quipper)
 the homotopy-typing reflects
topological quantum gates

2nd punchline

Riley: “Linear HoTT as a Quantum Certification Language”
[\[mvr.hosting.nyu.edu/pubs/translation.pdf\]](http://mvr.hosting.nyu.edu/pubs/translation.pdf)

Theorem: *Type transport in the following dependent type family expresses, in the above simplicial categorical semantics, topological braid quantum gates of $\mathfrak{su}(2)$ -anyons (Majorana, Fibonacci, ...)*

$$\left. \begin{array}{lll} \text{punctures} & \text{degree} & \text{shifted level} \\ N : \mathbb{N}_+, \quad n : \mathbb{N}, \quad \kappa : \mathbb{N}_{\geq 2} \\ w_{(-)} : N \rightarrow \{0, \dots, \kappa - 2\} \end{array} \right\} \vdash \left(\vec{z} \mapsto \left[(t : \mathbf{BC}^\times) \rightarrow \left(\text{fib}_{(t, \vec{z})} (\text{pr}_N^{N+n}, \tau_{(\kappa, w_\bullet)}) \rightarrow \mathbf{B}^n (\mathcal{C}_t \mathbb{C}_{\text{udl}}) \right) \right]_0 \right) : \mathbf{BPBr}(N) \rightarrow \text{Type}$$

weights

$$(206) \quad \text{pr}_N^{N+n} : \mathbf{BPBr}(N+n) \longrightarrow \mathbf{BPBr}(N)$$

$$\begin{array}{ccc} \text{pt} \\ b_{Ii} \end{array} \mapsto \begin{array}{ccc} \text{pt} \\ e \end{array}$$

$$\begin{array}{ccc} \text{pt} \\ b_{IJ} \end{array} \mapsto \begin{array}{ccc} \text{pt} \\ b_{IJ} \end{array}$$

$$\tau_{(\kappa, w_\bullet)} : \mathbf{BPBr}(N+n) \longrightarrow \mathbf{BC}^\times$$

(53) (223)

$$\begin{array}{ccc} \text{pt} \\ b_{Ii} \end{array} \mapsto \begin{array}{ccc} \text{pt} \\ \exp(2\pi i \frac{w_I}{\kappa}) \end{array}$$

$$\begin{array}{ccc} \text{pt} \\ b_{ij} \end{array} \mapsto \begin{array}{ccc} \text{pt} \\ \exp(2\pi i \frac{2}{\kappa}) \end{array}$$

$$\begin{array}{ccc} \text{pt} \\ \text{pt} \end{array} \mapsto \begin{array}{ccc} \text{pt} \\ \text{pt} \end{array}$$

(2)

in combination \Rightarrow major potential use-case:
certifying topological quantum compilation
 (arguably inevitable for industry-scale TQC)

the reported results are theoretical (\exists on paper)

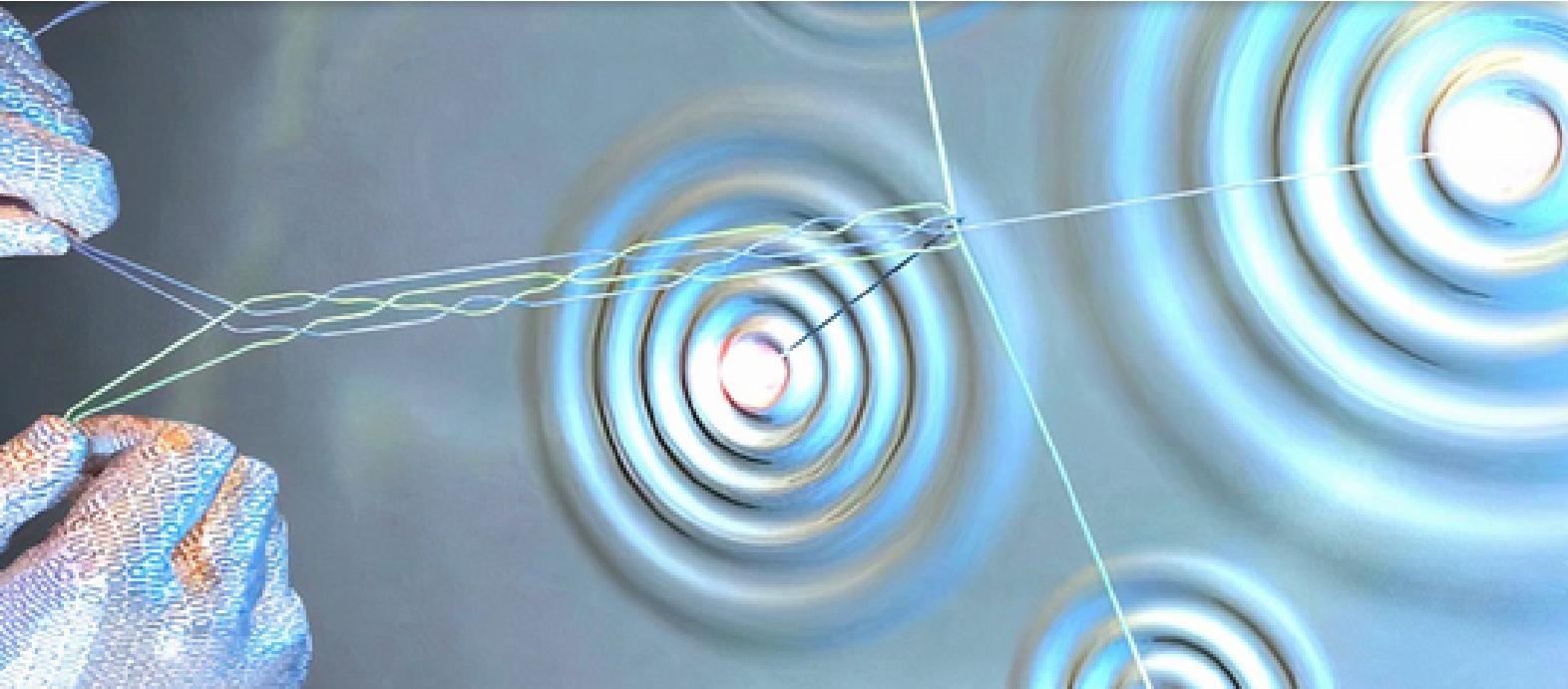
If you are a young energetic type theorist
interested in coding implementations \Rightarrow
contact us: nyuad.cqts.info@nyu.edu !

Thanks!

جامعة نيويورك أبوظبي

NYU ABU DHABI

Academics Admissions Research Campus Life Public Programs About



Center for Quantum and Topological Systems

Researchers Conferences

Home / Research / Faculty Labs and Projects / Center for Quantum and Topological Systems

Center for Quantum and Topological Systems