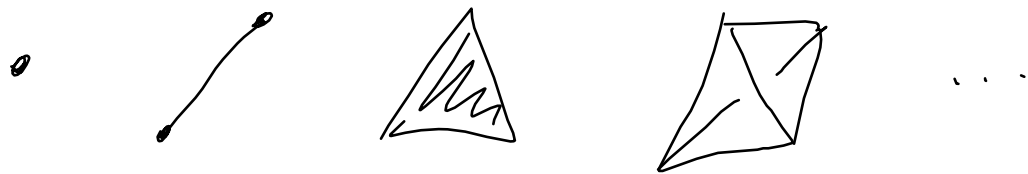


A simplex is an n -dimensional
triangle/tetrahedron etc.



These have many incarnations:

1) topological

$$\Delta^n = \left\{ (x_0, \dots, x_n) \in \mathbb{R}^{n+1} \mid \begin{array}{l} 0 \leq x_i \leq 1 \\ \sum_i x_i = 1 \end{array} \right\}$$

$$\cong \left\{ (x_1, \dots, x_n) \in \mathbb{R}^n \mid \begin{array}{l} 0 \leq x_i \leq 1 \\ \sum x_i \leq 1 \end{array} \right\}$$

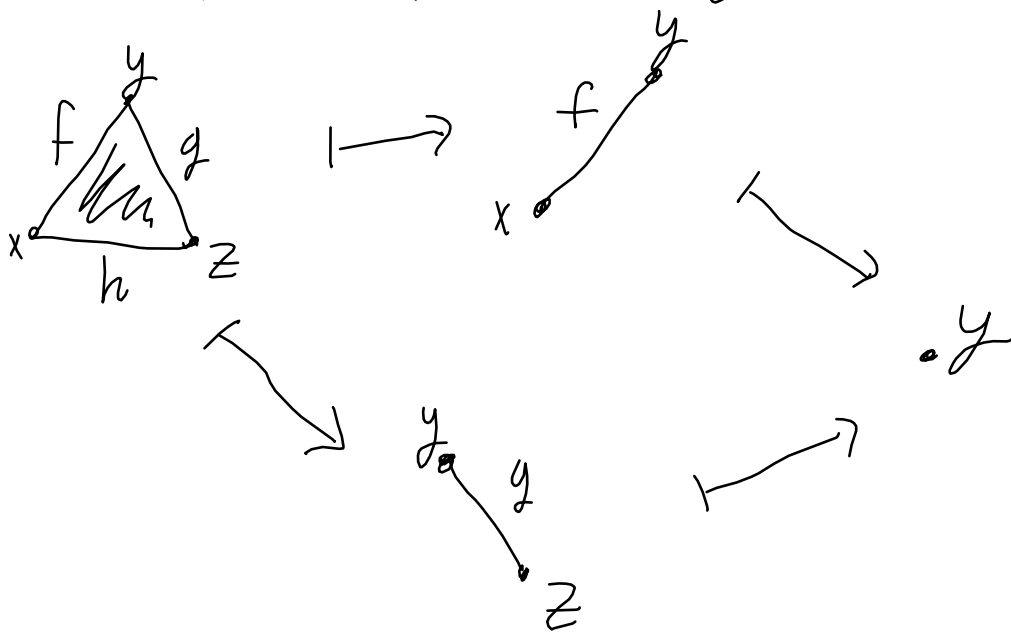
$$\cong D^n$$

but Δ^n has a natural way to
be glued together into cell cxs:

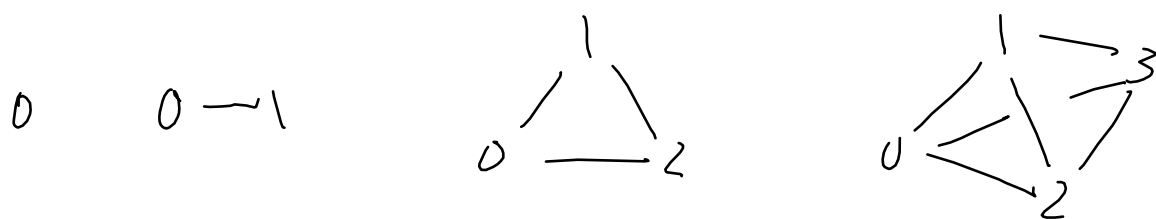
the faces of any n -simplex
are $(n-1)$ -simplices.

A semi-simplicial set is a combinatorial presentation of a way to fit simplices together.

- 1) A set X_n of n -simplices $\forall n \in \mathbb{N}$
- 2) Face maps $d_k: X_{n+1} \rightarrow X_n \quad \forall 0 \leq k \leq n+1$
- 3) $d_k d_l = d_{l-1} d_k$ for $k < l$.



We number the vertices (= 0-simplices)



& d_k is the face that omits the k^{th} vertex. Maybe a little backwards at first

$$d_0 \left(\begin{array}{c} 1 \\ / \quad \backslash \\ 0 \text{---} 2 \end{array} \right) = \begin{array}{c} 1 \\ \backslash \\ 2 \end{array}$$

$$d_1 = \begin{array}{c} 0 \text{---} 2 \end{array}$$

$$d_2 = \begin{array}{c} 1 \\ / \\ 0 \end{array}$$

Let $\Delta_m =$ the category presented by

- objects $[n] \quad \forall n \in \mathbb{N}$
- morphisms $d_k: [n] \rightarrow [n+1] \quad \forall 0 \leq k \leq n+1$
- equations $d_k d_{k-1} = d_{k-1} d_k$.

Then a sss is a functor $\Delta_m^{\text{op}} \rightarrow \text{Set}$.

$$\text{sss} = [\Delta_m^{\text{op}}, \text{Set}] = \text{Set}^{\Delta_m^{\text{op}}}$$

Can think of $[n]$ as the "free-living n -simplex" and d_k as the "inclusion".

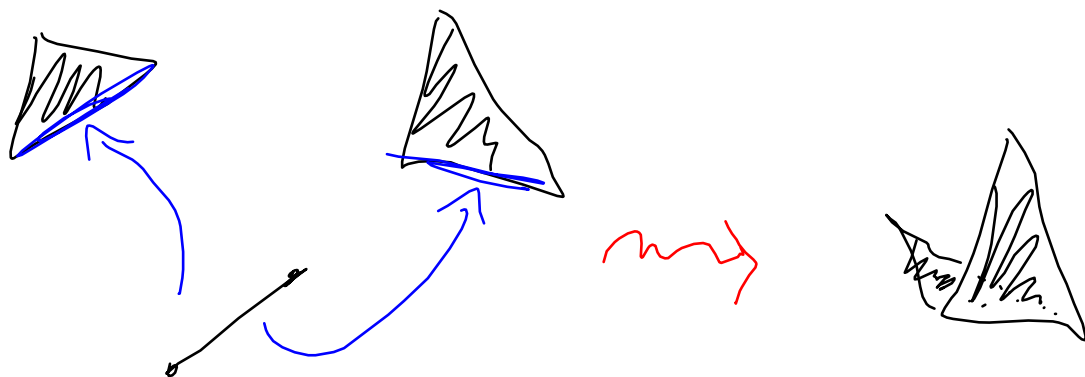
We can "realize" these simplices topologically, recall $\Delta^n \in \text{Top}$.

Have an actual inclusion $\delta_k: \Delta^n \hookrightarrow \Delta^{n+1}$

\rightsquigarrow a functor $\Delta_m \rightarrow \text{Top}$.

A sss X then gives us a way to glue these topological simplices together. The geometric realization

$$|X| = \sum_n X_n \times \Delta^n / \begin{array}{l} x \in X_{n+1}, p \in \Delta^n \\ (d_k x, p) = (x, \delta_k p) \end{array}$$



Not every top space has a cell presentation w/ simplices, but we can ask for a "best approximation."

What are all the possible simplices in a space Y ?

continuous maps $\Delta^n \rightarrow Y$.

$S_n Y = \{ \Delta^n \rightarrow Y \}$ is a SSS,
the singular complex of Y .

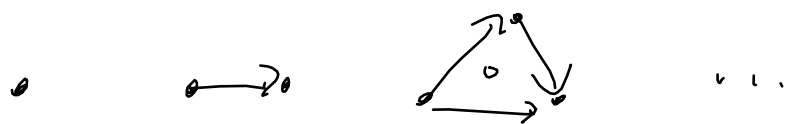
We have an adjunction

$$\text{SSS} \begin{array}{c} \xrightarrow{|-|} \\ \xleftarrow{S} \\ \xrightarrow{S} \end{array} \text{Top}$$

Much algebro-topological info is detected by $S.Y$.

We can do the same thing with
 Top replaced by any other category
 containing "concrete simplices".

eg Cat



$$\mathcal{S}\mathcal{S}\mathcal{S} \begin{array}{c} \xrightarrow{\tau_1} \\ \xleftarrow{\text{Nerve}} \end{array} \text{Cat}$$

eg $\mathcal{S}\mathcal{S}\mathcal{S}$ itself!

$(\Delta^k)_n = \Delta_m(n, k)$ the representable fr.

$$\mathcal{S}\mathcal{S} \begin{array}{c} \xrightarrow{\text{Id}} \\ \xleftarrow{\text{Id}} \end{array} \mathcal{S}\mathcal{S}$$

Density theorem - every presheaf is a
 colimit of representables.

BUT These functors are not as good as we'd like.

eg the SSS $\Delta_{SSS}^1 = \text{hook} = \begin{pmatrix} \vdots \\ 0 \\ \downarrow\downarrow \\ 1 \\ \downarrow\downarrow \\ 2 \end{pmatrix}$

has $|\Delta_{SSS}^1| = \Delta_{\text{top}}^1 = \text{hook}$

of course, $\Delta_{\text{top}}^1 \times \Delta_{\text{top}}^1 = \boxed{\text{scribble}}$

but $\Delta_{SSS}^1 \times \Delta_{SSS}^1 = \begin{matrix} 0 \\ \downarrow\downarrow \\ 1 \\ \downarrow\downarrow \\ 4 \end{matrix}$

$\text{hook} \quad \text{not much like} \quad \boxed{\text{scribble}}$

ie $|-|: \text{SSS} \rightarrow \text{Top}$ doesn't preserve products.



A simplicial set is a sss X w/

• degeneracy maps $s_k: X_n \rightarrow X_{n+1}$, $0 \leq k \leq n$

• $d_k s_l = s_{l-1} d_k$ $k < l$

$d_k s_k = d_{k+1} s_k$

$d_k s_l = s_l d_{k-1}$ $k \geq l+1$

$s_k s_l = s_{l+1} s_k$ $k \leq l$

Not supposed to understand those
deeply.

$\mathcal{A} = \Delta_m$ w/ σ_k 's.

$\text{Set} = [\Delta^{\text{op}}, \text{Set}]$.

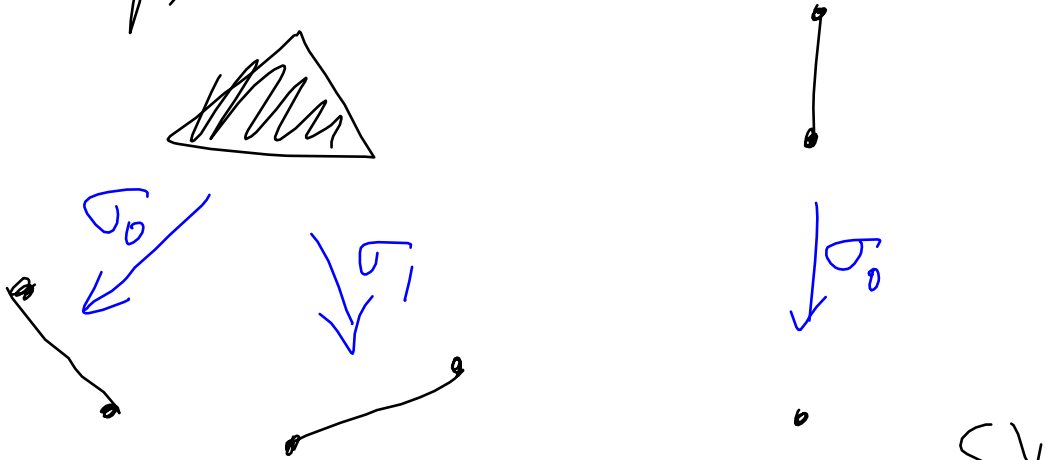
(Cute fact: $\Delta \cong$ finite nonempty
linearly ordered sets)

Idea: $\sigma_k: [n+1] \rightarrow [n]$ "squashes"

an $(n+1)$ -simplex down to an n -simplex

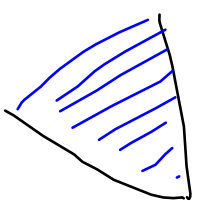
so $S_k: X_n \rightarrow X_{n+1}$ assigns squashed
or degenerate simpl.

eg in Top.



so $S_0(x) = \text{constant paths at } x$

for $f: \Delta_{top}^1 = [0,1] \rightarrow Y$,

$S_0(f) =$  is const along
the blue lines.

in Cat

$$S_0(\text{object } x) = \begin{array}{ccc} & \text{id}_x & \\ & \longrightarrow & \\ x & & x \end{array}$$

$$S_0(x \xrightarrow{f} y) = \begin{array}{ccc} & x & \\ \text{id}_x & \nearrow & f \\ x & \xrightarrow{f} & y \end{array}$$

$$S_1(x \xrightarrow{f} y) = \begin{array}{ccc} & y & \\ f & \nearrow & \\ x & \xrightarrow{f} & y \end{array}$$

Now $|-\!| : \mathcal{S}\text{Set} \rightarrow \text{Top}$, $\tau_i : \mathcal{S}\text{Set} \rightarrow \text{Cat}$

do preserve products,

$(\Delta'_{\text{set}})_n = \Delta([n], [1])$ is more complicated

$$\left(\dots \begin{array}{c} \rightleftharpoons \\ \rightleftharpoons \\ \rightleftharpoons \end{array} 3 \begin{array}{c} \rightleftharpoons \\ \rightleftharpoons \\ \rightleftharpoons \end{array} 2 \right)$$

$$\begin{pmatrix} x \\ |f \\ y \end{pmatrix} \times (z \xrightarrow{g} w) = \begin{array}{ccc} xz & \xrightarrow{s_{x,g}} & xw \\ f|_{s_z} & \searrow fg & |_{s_w} f \\ yz & \xrightarrow{s_{y,g}} & yw \end{array}$$

↪ ↪

s_f, s_g s_f, s_g

ie even if $a \in X_n$, $b \in Y_n$ are degen,
 $(a, b) \in (X \wedge Y)_n$ need not be!

(This is not a proof, but it is true that $|-|$ preserves products... if you define "Top" correctly.)